Maximal avoidable sets of words

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We use the following notations.

 Σ : an alphabet (a finite set of letters),

 Σ^* : the set of words over Σ ,

 Σ^{ω} : the set of infinite words (sequences),

 $\Sigma^{\#} := \Sigma^{*} \cup \Sigma^{\omega}$.

 Σ^{+} := Σ^{*} - {1}, where 1 is the empty word.

For $x = a_1 a_2 \cdots$, $y = b_1 b_2 \cdots \in \Sigma^{\#}$ define the distance of x and y by $d(x, y) = \frac{1}{\min \{n \mid a_n \neq b_n\}}.$

As is well-known ([4]), ($\Sigma^{\#}$, d) is a compact totally disconnected metric space.

Let $x, y \in \Sigma^*$ and $X \in \Sigma^\#$. We say y <u>avoids</u> x, if y does not contain x as subword, and y <u>avoids</u> X, if y avoids every x in X. X is called <u>avoidable</u>, if there is an infinite word y avoiding X, otherwise X is called <u>unavoidable</u>. Avoidability of sets of words called patterns were studied in [1].

Example 1. Let $X=\{v^2\mid v\in \Sigma^\dagger\}$. Then y avoids X if and only if y is square-free. It is a famous fact that X is avoidable if $|\Sigma|\geq 3$

([5]).

An avoidable set X is $\underline{\text{maximal}}$, if any set properly containing X is $\underline{\text{unavoidable}}$.

Theorem 1. For any avoidable set X, there is a maximal avoidable set containing X.

For a given $X \subset \Sigma^*$, the set

 $Min(X) = \{x \in X \mid any \ x' \in X \text{ is not a proper subword of } X \}$ is called the <u>base</u> of X. Easily we see y avoids X if and only if y avoids Min(X). X is <u>finitely based</u> if Min(X) is a finite set. The base of a maximal avoidable set is called a <u>critical set</u> of words.

Corollary. For any avoidable set X which is factor-free, that is, any word in X is not a subword of another word in X, there is a critical set containing X.

Example 2. Let Σ = {a, b}. Then, {a², ab, ba} and {a², b²} are critical sets.

An infinite word x is <u>recurrent</u>, if for any subword v of x, there is an integer k(v) > 0 such that any subword of length k of x contains v as subword. In this situation v is said to be recurrent in x. If $x = v^{\omega}$ for some $v \in \Sigma^*$, x is called <u>periodic</u>; the shortest such v is the <u>period</u> of x. A periodic infinite word is recurrent, but the converse is not true.

The <u>shift transformation</u> τ is a mapping from Σ^ω to itself defined

 $\tau(x) = a_2 a_3 \cdots$ for $x = a_1 a_2 \cdots$.

Obviously, t is a surjective continuous mapping.

A <u>subshift</u> S is a non-empty closed subset of Σ^{ω} invariant under τ . S is <u>minimal</u>, if it does not contain a subshift properly. For a given set X C Σ^* of words, S(X) is the set of infinite words avoiding X. For a given subshift S C Σ^{ω} , X(S) is the set of words which do not appear as subwords of elements of S.

Theorem 2. For an avoidable set X, S(X) is a subshift. If X is maximal, then S(X) is minimal. Conversely, if S is a subshift, then X(S) is an avoidable set. If S is minimal, then X(S) is maximal. This gives a 1-1 correspondence between maximal avoidable sets and minimal subshifts.

Lemma 1. An avoidable set X is maximal if and only if any word out of X is recurrent in any infinite word avoiding X.

Theorem 3 (Morse-Hedlund [4]). S C Σ^ω is a minimal subshift if and only if

$$S = \{ \tau^{n}(x) \mid n = 0, 1, 2, \dots \}$$

for some recurrent infinite word x. Moreover,

- (1) S is perfect, if x is non-periodic. In this case every element in S is non-periodic.
- (2) S is finite, if x is periodic. In this case every element in S is periodic.

Corollary (c.f. [3, Theorem 4.2]). Let X be an avoidable set such that for any $v \in \Sigma^+$, v^n does not avoid X for x >> 0, then S(X) contains a perfect subset

Theorem 4. Let X be a maximal avoidable set. Then, S(X) is finite if and only if X is finitely based.

For an avoidable set X, the <u>radical</u> rad(X) of X is the intersection of all the maximal avoidable sets containing X. X is called <u>reduced</u>, if X = rad(X).

Lemma 2. A word v is in rad(X), if and only if any recurrent infinite word avoiding X avoids v.

Corollary. Any word out of rad(X) is extensible to a recurrent infinite word avoiding X.

Theorem 5. If X is a reduced avoidable set, then every isolated point of S(X) is periodic.

Corollary. If X is a reduced avoidable set such that for any $v\in \Sigma^+$, v^n does not avoid X for n >> 0, then S(X) is perfect.

A set X of words is quasi maximal, if rad(X) is maximal.

Theorem 6. Let X be an avoidable set. Then following statements are equivalent.

- (1) X is quasi-maximal.
- (2) S(X) contains a unique minimal subshift.
- (3) For any n > 0, there is a word v of length n such that $X \cup \{v\}$ is unavoidable.
- (4) For any word w such that X U $\{w\}$ is unavoidable and for any n > 0, there is a word v of length n such that X U $\{wv\}$ is unavoidable.

Example 3. Let $X = \{a^2, bab\} \subset \{a, b\}^*$. Then, b^{ω} and ab^{ω} are only infinite words avoiding X, and $X \cup \{b^n\}$ is unavoidable for any n > 0.

Thus X is quasi-maximal.

An unavoidable set X is said to be $\underline{minimal}$, if X - {v} is avoidable for any $v \in X$. As is easily seen ([2]), a minimal unavoidable set is finite.

Conjecture I (Ehrenfeucht, see [2]). For any unavoidable set X, there is a word $x \in X$ and a letter $a \in \Sigma$ such that $(X - \{x\}) \cup \{xa\}$ is unavoidable.

Conjecture II. For any minimal unavoidable set X, there is a word x in X such that $X - \{x\}$ is a quasi-maximal avoidable set.

Theorem 7. Conjecture I and Conjecture II are equivalent.

References

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