

On Physarum solver on Riemannian surface

宮路 智行, 大西 勇 Tomoyuki Miyaji** and Isamu Ohnishi*

Department of Mathematical and Life Sciences,
Graduate School of Science, Hiroshima University,
1-3-1 Kagamiyama, Higashi-Hiroshima, 739-8526 JAPAN.

** denpaatama@hiroshima-u.ac.jp

* isamu_o@math.sci.hiroshima-u.ac.jp

概要

In this paper we report a mathematical theory inspired by *P.polycephalum*. It is an amoeba-like organism that exhibits the shortest path-finding behavior in a labyrinth. Physarum solver is a model equation of this interesting behavior of the creature. In this paper, we prove mathematically rigorously it on graph included in two dimensional manifold with Riemannian metric, that is, Riemannian surface. In other words, we give a rigorous proof of the fact that the equilibrium point corresponding to the shortest path is globally asymptotically stable for Physarum solver on Riemannian surface. Moreover, we extend our theorem to a surface composed of the finite number of Riemannian surfaces pasted with one another continuously.

1 Introduction

The plasmodium of true slime mold *P.polycephalum* is a large amoeba-like organism and very interesting. It possesses a lot of nucleuses, but composed of a single cell. Its body contains a tube network by means of which nutrients and signals circulate through the body in effective manner. When food sources were presented to a starved plasmodium that was spread over the entire agar surface, it concentrated at every food source, respectively. Almost the entire plasmodium accumulated at the food sources and covered each of them in order to absorb nutrients [9]. Only a few tube remained connecting the quasi-separated components of the plasmodium through the short path. Nakagaki *et al.* showed that this simple organism had the ability

*The first author is supported by a grant of "Initiative program of facinating educations of graduate school" in Ministry of Education, Culture, Sports, Science and Technology in Japan, and moreover, he is also an encouraging young researcher with resaerch grant of JSPS, and the second author is supported by Grant No.17540118 of JSPS.

to find the minimum-length solution of a labyrinth [10, 11]. The connecting tube traces the shortest path even in a complicated labyrinth. This adaptation process of the tube network is based on an underlying physiological mechanism, that is, a tube becomes thicker as a flux in the tube is larger.

Tero *et al.* made a mathematical model in consideration of the qualitative mechanisms clarified by experiments [13]. They considered the tube network of *P.polycephalum* on a labyrinth to be a plane graph, set some variables on vertices and edges of the graph, and described the process of growth and degeneracy of the tube. The model consists of two parts, equations for flux in tubes and nonlinear ODEs for adaptation of tubes. In a special case of nonlinear terms (in which the adaptation term has just a linear form), the model is called *Physarum solver*. According to numerical simulation results, the minimum-length solution of a labyrinth can be obtained as an asymptotic steady state of *Physarum solver* [13, 15].

We have already obtained several results in interesting cases to report it in [5]. In this paper, we present an extension of the theorem in a further more general case on two dimensional surface, in which the length of any two points on it can be computable. We note that the first announcement of this paper's result has been already reported in [6], but this is a little bit more general theorem proved here than it.

As telling from the mathematically technical point of view, it is an interesting point that we prove the globally asymptotical stability without using Lyapunov function, which is defined globally in the whole space. Instead of the function, we use a kind of "weak" Lyapunov function defined in an open bounded set to operate a graph under consideration. From the viewpoint of dynamical system theory, such an operation of graph is corresponding to reduction of the system to an exponential attracting invariant set, which is the "lower" boundary of the open bounded set. We make the reduction of the system by using of "exponential tracking" property of *Physarum solver*. (See, for instance, [3] for the detail of the reduction.) This property is also proved by our "weak" Lyapunov function. Moreover, the inertial form of the system of equations has expressed explicitly and has a similar form to the original one, but its dimension is lower than of the original one, which can let us make a progress of it to the next step. We thus make this procedure a lot of times to get the result. Here it is necessary for the length of each edge of the graph to be calculated. Therefore, we put the assumption about the two dimensional surface under consideration.

On the other hand, in laboratory, *P.polycephalum* sometimes makes a mistake, for example, when there is a kind of "double-edge" like parallel

circuit in a graph. Then, it sometimes chooses a longer path. The model equation can explain this case, that is, if the adaptation term has a super-linear form, then numerical computations show that it sometimes chooses a longer path according to initial data ([13]). We also have already studied such a case in [7] to get some mathematical results. In that paper, we explain this fact from the viewpoint of dynamical system mathematically. In the consequence, we do not only show why *P.polycephalum* makes a mistake, but also show how it makes a mistake. It is a kind of analysis of a simple but interesting transition process, which can be significant also in real world application. Actually, as *P.polycephalum* is in various situations, we can consider that it changes important parameters according to the situation. This is the reason why it sometimes takes the right answer, but sometimes makes a mistake. It is very interesting that this model can explain such a kind of mysterious behavior of the organism of life just by the subtle difference of the adaptation term.

As telling about an application of Physarum solver, in [15], they have attempted to apply Physarum solver to navigation system of road map. In fact, Physarum solver has found the shortest path connecting between Seattle and Houston admirably. In such a way, as it has a lot of possibilities to apply, we here consider that it is very important to ensure the shortest path decision property of Physarum solver mathematically rigorously.

Finally, we remark that the definitions, the notations and the results stated below can be extended to a surface composed of the finite number of Riemannian surfaces pasted with one another continuously, and a graph on such a surface. This is because the length of each edge are computable on the graph, even if the edge which crosses a pasted line between two surfaces exists. Therefore, we can understand a biochemical experiments in Nakagaki et.al. [14] recently by use of *P.polycephalum* where “the shortest path” is found on a non-homogenized place by *P.polycephalum*, in this mathematical framework. This is because Physarum solver with weight on tube can be regarded as Physarum solver on a kind of non-homogenized place by translating the weight to the field where the system of equations are defined.

2 Preliminaries

2.1 Graphs on Riemannian Surface

Physarum solver is defined on a finite graph in Riemannian surface. First, we briefly introduce some notations. We will assume that the reader is

familiar with basic terms and results from graph theory. See, for details, [2, 4, 8].

In Riemannian surface M^2 , a graph $G = (V, E)$ is a pair of sets, $V, E \subset M^2$, where V is a nonempty and E is a set of 2-element subsets of V . Throughout this paper we assume that V is finite. The elements of V are called vertices of G , the elements of E are the edges of G . Let $|G|$ denote the number of vertices. $|G|$ is called the order of the graph G . The degree of a vertex v , denoted $d(v)$, is the number of edges incident with v .

Let G_1, \dots, G_k be subgraphs of the graph G . The union $G_1 \cup \dots \cup G_k$ is the graph $H \subseteq G$ with $V(H) = V(G_1) \cup \dots \cup V(G_k)$ and $E(H) = E(G_1) \cup \dots \cup E(G_k)$.

For $U \subset V$ we denote by $G - U$ the subgraph of G obtained by deleting from G the vertices in U and all edges incident with them. If $F \subset E$, then $G - F$ is the subgraph of G obtained by removing from G the edges in the set F .

A path is a nonempty graph $P = (V, E)$ such that

$$V = \{x_0, x_1, \dots, x_k\} \quad E = \{x_0x_1, x_1x_2, \dots, x_{k-1}x_k\},$$

where $x_i x_j$ is the edge joining x_i and x_j . We denote $P = x_0 x_1 \dots x_k$ and call it the path from x_0 to x_k , or x_0 - x_k path.

If $P = x_0 \dots x_{k-1}$ is a path and $k \geq 3$, then a graph $C = P + x_{k-1}x_0$ is called a cycle.

A graph is in M^2 , if it is drawn in M^2 in such a way that no edges intersect on M^2 , except at a common end-vertex. Then, we call the graph M^2 graph. A planar graph is a simple example. For any graph G , a set $M^2 \setminus G$ is an open subset of M^2 . Its region is called a face of G . Since G is bounded, one of the faces can be unbounded. The unbounded face of G is called outer face, if it exists, and the remainders are called inner faces. If M^2 is compact, then we take an adequate face in all the faces $\subset M^2$, and we call this be outer and call the others inner. (Later, we choose the face including s and t on its boundary as the outer.)

Note that, in what follows, we consider only M^2 graph as a graph. Therefore we will state a graph simply, but it means M^2 graph.

2.2 Physarum solver

Now, let us briefly introduce *Physarum solver*. The formulation and physiological backgrounds are detailed in [13](See also [15]).

Let $G = (V, E)$ be a graph. We consider a set $N = (G, s, t, L)$, where $s, t \in V$ are two distinguished vertices, and L is a function from E to \mathbb{R}_+ .

Assume that G is a connected graph and $|G| \geq 2$. Let $V = \{v_0, v_1, \dots, v_n\}$, where $n \geq 1$. Let e_{ij} denote the edge joining v_i and v_j if it exists. If G has multiple edges joining v_i and v_j , we denote them by $v_{ij}^1, v_{ij}^2, \dots$. Let $L(e_{ij}) = L_{ij} > 0$ be a length (or weight) of the edge e_{ij} . We can calculate the length, because M^2 is Riemannian. Assume that $L(e_{ij}) = L(e_{ji})$. G is considered as a flow network flowing out from v_0 and sinking into v_n . To distinguish those two vertices, let us call $s = v_0$ a source and $t = v_n$ a sink (or a target). Assume that there exist exactly one source and one target. For a path $P = v_{\beta_0} \cdots v_{\beta_k}$ in G , define the length of the path P to be

$$L(P) = \sum_{i=0}^{k-1} L(e_{\beta_i \beta_{i+1}}).$$

Remark that, throughout this paper, the length of a path does not mean the number of edges which compose the path.

Let τ denote the time variable. For each i , the variable $p_i(\tau)$ is a pressure at the vertex v_i . For each edge e_{ij} , $D_{ij}(\tau)$ and $Q_{ij}(\tau)$ are its conductivity (corresponding to a thickness of tube) and flux, respectively. In addition, as described above, e_{ij} has its length L_{ij} . For each edge, D_{ij} should be non-negative and $D_{ij} = D_{ji}$. Let $\mathbf{D}(\tau) = (D_{ij}(\tau))_{i,j}$ be the set of all $D_{ij}(\tau)$'s. Note that p_i, D_{ij} and Q_{ij} are variables depending on time τ , and L_{ij} is a positive constant. We want to obtain the s - t path such that its length is smaller than that of any other s - t paths.

First, we give two rules for flux. The flux Q_{ij} is given by

$$Q_{ij} = \frac{D_{ij}}{L_{ij}}(p_i - p_j) = g_{ij}(p_i - p_j), \quad (1)$$

where $g_{ij} = D_{ij}/L_{ij}$ is the conductance of the edge e_{ij} . (1) is an analogy of *Ohm's law* for electric circuits. It is clear that $Q_{ij} = -Q_{ji}$. Moreover, we assume the *Kirchhoff's law* at each node:

$$\sum_{j \neq i} Q_{ij} = \begin{cases} I_0 & \text{if } i = 0, \\ 0 & \text{if } 0 < i < n, \\ -I_0 & \text{if } i = n, \end{cases} \quad (2)$$

where I_0 is the flux from the source vertex. In this model, it is assumed that I_0 is a positive constant.

Next, we give an adaptation rule of conductivity:

$$\dot{D}_{ij} = |Q_{ij}| - D_{ij}, \quad (3)$$

where we use \dot{x} to represent the derivative $dx/d\tau$.

By setting $p_n = 0$ as a basic pressure level, all p_i 's are determined by (2). Then Q_{ij} 's are determined by (1), and the evolution of D_{ij} 's is described by (3). We call the system (1),(2),(3) *Physarum solver* for $N = (G, s, t, L)$.

We sometimes consider an orientation of graph G by means of the direction of flux. If $Q_{ij} > 0$, then we suppose that e_{ij} is oriented from v_i to v_j . Naturally, this orientation can change as time passes.

3 Mathematical analysis

In this section, we analyze Physarum solver(1),(2),(3) mathematically. In the first two subsections, we discuss the system without time evolution. Then, in §3.3, we give a proof of the main theorem.

3.1 Kirchhoff's Law

Here we discuss a system of linear equations derived from Kirchhoff's law (2). Solving this system, we obtain the values of pressure at each vertex.

Let $G = (V, E)$ be a graph with $|G| = n + 1$ for $n \geq 1$. Let $V = \{v_0, \dots, v_n\}$, $s = v_0$, and $t = v_n$. In this subsection, we interpret that $g_{ij} > 0$ if $e_{ij} \in G$, otherwise $g_{ij} = 0$. It allows us to simplify the summation notation because we can discuss as if G is a complete graph.

Substitute (1) for (2), then the equation to be solved is

$$\sum_{j \neq i} g_{ij}(p_i - p_j) = \begin{cases} I_0 & \text{if } i = 0, \\ 0 & \text{if } 1 \leq i \leq n - 1, \end{cases} \quad (4)$$

where $p_n = 0$ and $g_{ij} > 0$. (4) is written in matrix form

$$\mathbf{A}\mathbf{p} = \mathbf{b}, \quad (5)$$

where

$$\mathbf{p} = {}^t(p_0, p_1, \dots, p_{n-1}) \quad , \quad \mathbf{b} = {}^t(I_0, 0, \dots, 0), \quad (6)$$

and $\mathbf{A} = (\mathbf{A}_{ij})$ is a square matrix of order n given by

$$\mathbf{A}_{ij} = \begin{cases} \sum_{l \neq i} g_{il} & \text{if } i = j, \\ -g_{ij} & \text{otherwise,} \end{cases} \quad (7)$$

where $i, j = 0, \dots, n - 1$. We first study the solution of (5).

Proposition 1. *The coefficient matrix \mathbf{A} is a symmetric nonsingular M -matrix.*

Proof. A symmetric matrix whose off-diagonal entries are all non-positive is a symmetric nonsingular M-matrix if and only if it is positive definite(cf.[1]). It is obvious that the off-diagonal entries of \mathbf{A} are non-positive from the definition. Hence \mathbf{A} is a matrix of the class $Z^{n \times n}$ in [1]. It is also obvious that the coefficient matrix \mathbf{A} is symmetric from its definition. Now we consider $\langle \mathbf{p}, \mathbf{A}\mathbf{p} \rangle$ for $\mathbf{p} \neq 0$, where $\langle \cdot, \cdot \rangle$ is the standard inner product on \mathbb{R}^n . We can obtain $\langle \mathbf{p}, \mathbf{A}\mathbf{p} \rangle > 0$. In fact,

$$\langle \mathbf{p}, \mathbf{A}\mathbf{p} \rangle = \sum_{i=0}^{n-1} p_i \sum_{j \neq i} g_{ij}(p_i - p_j) = \sum_{i=0}^{n-1} \sum_{j \neq i} g_{ij} p_i^2 - \sum_{i=0}^{n-1} \sum_{j \neq i} g_{ij} p_i p_j.$$

As $p_n = 0$, we get

$$\langle \mathbf{p}, \mathbf{A}\mathbf{p} \rangle = \sum_{i=0}^{n-1} \sum_{j \neq i, j < n} g_{ij} p_i^2 - \sum_{i=0}^{n-1} \sum_{j \neq i, j < n} g_{ij} p_i p_j + \sum_{i=0}^{n-1} g_{in} p_i^2.$$

In addition, as $g_{ij} = g_{ji}$, the following two hold:

$$\begin{aligned} \sum_{i=0}^{n-1} \sum_{j \neq i, j < n} g_{ij} p_i^2 &= \sum_{0 \leq i < j < n} g_{ij} (p_i^2 + p_j^2) \\ \sum_{i=0}^{n-1} \sum_{j \neq i, j < n} g_{ij} p_i p_j &= \sum_{0 \leq i < j < n} 2g_{ij} p_i p_j \end{aligned}$$

Thus we obtain

$$\langle \mathbf{p}, \mathbf{A}\mathbf{p} \rangle = \sum_{0 \leq i < j < n} g_{ij} (p_i - p_j)^2 + \sum_{i=0}^{n-1} g_{in} p_i^2.$$

Since $g_{ij} \geq 0$, we obtain $\langle \mathbf{p}, \mathbf{A}\mathbf{p} \rangle > 0$ for $\mathbf{p} \neq 0$. It implies that \mathbf{A} is positive definite. Therefore \mathbf{A} is a symmetric nonsingular M-matrix. \square \square

Proposition 2. *The system (5) has a unique non-negative solution \mathbf{p} .*

Proof. Since \mathbf{A} is an M-matrix, $\mathbf{A}\mathbf{p} = \mathbf{b} > 0$ implies $\mathbf{p} > 0$. In addition, as \mathbf{A} is nonsingular, \mathbf{p} is unique. \square \square

The next proposition guarantees that the vertex $s = v_0$ actually works as a source and $t = v_n$ works as a sink.

Proposition 3. *For any $v_j \in G$, $p_0 \geq p_j$.*

Proof. Let $\beta = \{\beta_i\}_{i=1}^k$ and suppose that $p_0 < p_{\beta_i}$ for $i = 1, 2, \dots, k$. Then we obtain

$$\sum_{i=1}^k \sum_{j \neq \beta_i} Q_{\beta_i j} = \sum_{i=1}^k \sum_{j \notin \beta} g_{\beta_i j} (p_{\beta_i} - p_j) > 0.$$

However, this sum must be zero from (2). This is contradiction. \square \square

It follows that all Q_{ij} 's are bounded.

Proposition 4. For any $e_{ij} \in G$, $|Q_{ij}| \leq I_0$.

Proof. Now we may assume without loss of generality that the pressures satisfy

$$p_0 \geq p_1 \geq \cdots \geq p_{n-1} \geq p_n = 0. \quad (8)$$

Thus Q_{ij} is nonnegative whenever $i < j$.

According to Kirchhoff's law around the vertex $s = v_0$, we have $\sum_{j=1}^n Q_{0j} = I_0$. Since $Q_{0j} \geq 0$ for $j > 0$, we obtain $0 \leq Q_{0j} \leq I_0$.

Next, according to Kirchhoff's law around the vertex v_1 , we have $\sum_{j \neq 1} Q_{1j} = 0$. It can be rewritten as

$$Q_{01} = \sum_{j=2}^n Q_{1j}. \quad (9)$$

Since $Q_{01} \leq I_0$ and $Q_{1j} \geq 0$ for $j > 1$, we obtain $0 \leq Q_{1j} \leq I_0$ for $j > 1$.

Generally, according to Kirchhoff's law around the vertex v_i with $1 \leq i \leq n-1$, we have

$$\sum_{j=0}^{i-1} Q_{ji} = \sum_{j=i+1}^n Q_{ij}. \quad (10)$$

We show that $0 \leq Q_{ij} \leq I_0$ for $i < j$. As $Q_{ij} \geq 0$ for $j > i$, it suffices to prove $\sum_{j=i+1}^n Q_{ij} \leq I_0$. For the term $Q_{i-1,i}$ in the left hand side of (10), we get

$$Q_{i-1,i} \leq \sum_{j=i}^n Q_{i-1,j} = \sum_{j=0}^{i-2} Q_{j,i-1}.$$

Hence we obtain

$$\sum_{j=i+1}^n Q_{ij} \leq \sum_{j=0}^{i-2} Q_{ji} + \sum_{j=0}^{i-2} Q_{j,i-1} = \sum_{j=0}^{i-2} (Q_{ji} + Q_{j,i-1}).$$

For the terms $Q_{i-2,i}$, $Q_{i-2,i-1}$ in the right hand side, we get

$$Q_{i-2,i} + Q_{i-2,i-1} \leq \sum_{j=i-1}^n Q_{i-2,j} = \sum_{j=0}^{i-3} Q_{j,i-2}.$$

Therefore we get

$$\sum_{j=i+1}^n Q_{ij} \leq \sum_{j=0}^{i-3} (Q_{ji} + Q_{j,i-1} + Q_{j,i-2}).$$

Iteratively, it follows that

$$\sum_{j=i+1}^n Q_{ij} \leq \sum_{j=0}^i Q_{0j}.$$

Now, it is obvious that

$$\sum_{j=0}^i Q_{0j} \leq \sum_{j=0}^n Q_{0j} = I_0.$$

Thus we obtain

$$\sum_{j=i+1}^n Q_{ij} \leq \sum_{j=1}^n Q_{0j} = I_0. \quad (11)$$

Since $Q_{ij} \geq 0$ for $i < j$, we obtain $0 \leq Q_{ij} \leq I_0$ for $i < j$. As $Q_{ij} = -Q_{ji}$, it follows that $|Q_{ij}| \leq I_0$ for any $e_{ij} \in G$. \square \square

Proposition 5. Let $\beta = \{\beta_i\}_{i=0}^k \subset \{0, 1, \dots, n\}$ and $\beta_0 = 0, \beta_l = n$, and $l < n$. Suppose that $D_{\beta_i \beta_{i+1}} > 0$ for $0 \leq i \leq l-1$. When $D_{rs} \rightarrow 0$ for every r, s such that $(r, s) \neq (\beta_i, \beta_{i+1})$ for any $0 \leq i \leq l-1$, all p_i 's remain finite.

Proof. Since p_0 is the largest pressure value, we show that p_0 is upper bounded. According to the previous proposition, we obtain

$$\begin{aligned} p_0 &= p_0 - p_n = \sum_{i=0}^{l-1} (p_{\beta_i} - p_{\beta_{i+1}}) \\ &= \sum_{i=0}^{l-1} \frac{Q_{\beta_i \beta_{i+1}}}{g_{\beta_i \beta_{i+1}}} \leq \sum_{i=0}^{l-1} \frac{I_0}{g_{\beta_i \beta_{i+1}}}. \end{aligned}$$

Thus p_0 is upper bounded. \square \square

Now we consider the orientation of G by the direction of flow as described in the previous section. If $p_i > p_j$ and $e_{ij} \in G$ for $v_i, v_j \in G$, let e_{ij} be oriented from v_i to v_j . Let \vec{E} be the set of orientable edges in such a way and $\vec{G} = (V, \vec{E})$, then \vec{G} is a directed graph with one source $s = v_0$ and one target $t = v_n$. Note that $E = \vec{E}$ does not always hold and \vec{E} can change as time passes. The next is, however, a universal property which holds through the adaptation process.

Proposition 6. \vec{G} is acyclic, that is, \vec{G} has no directed cycle.

Proof. Assuming that there exists a directed cycle $\vec{C} = v_i v_j \cdots v_k v_i$, then we obtain $p_i > p_j > \cdots > p_k > p_i$. This is contradiction. \square \square

3.2 Equilibria and s - t paths

Here we discuss the relation between equilibria of an adaptation equation and s - t paths in a graph G . In an equilibrium state, i.e. $\dot{D}_{ij} = 0$, it holds that

$$D_{ij} (|p_i - p_l| - L_{ij}) = 0 \quad (12)$$

$$\sum_j \frac{D_{ij}}{L_{ij}}(p_i - p_j) = \begin{cases} I_0 & \text{if } i = 0, \\ 0 & \text{if } i \neq 0, n, \\ -I_0 & \text{if } i = n. \end{cases} \quad (13)$$

According to (13), there always exists at least one s - t path such that $D_{ij} > 0$ for all e_{ij} in the path.

We assume that the length of s - t path is different each other. For given any s - t path $P = v_{\beta_0} \cdots v_{\beta_k}$, let $\bar{\mathbf{D}}_P = (D_{ij})_{i,j}$ be the conductivities with

$$\begin{cases} D_{ij} > 0 \text{ and } |p_i - p_j| = L_{ij} & \text{if } e_{ij} \in P, \\ D_{ij} = 0 & \text{otherwise,} \end{cases} \quad (14)$$

where $v_{\beta_0} = s$ and $v_{\beta_k} = t$. Then $\bar{\mathbf{D}}_P$ satisfies (12). Let us call $\bar{\mathbf{D}}_P$ the equilibrium point corresponding to the path P .

It is easy to characterize $\bar{\mathbf{D}}_P$. The next proposition immediately follows from (12)-(14).

Proposition 7. *The following two hold.*

1. *The pressures at vertex $v_{\beta_i} \in P$ is given by*

$$p_i = \begin{cases} \sum_{j=i}^{k-1} L_{\beta_j \beta_{j+1}} & \text{if } i = 0, \dots, k-1, \\ 0 & \text{if } i = k. \end{cases} \quad (15)$$

Epecially, p_0 is equal to the length of P .

2. *The flux and the conductivity at edge $e_{\beta_i \beta_{i+1}} \in P$ are*

$$Q_{\beta_i \beta_{i+1}} = D_{\beta_i \beta_{i+1}} = I_0.$$

Therefore, the equilibrium point $\bar{\mathbf{D}}_P$ is given by

$$D_{ij} = \begin{cases} I_0 & \text{if } e_{ij} \in P, \\ 0 & \text{otherwise.} \end{cases} \quad (16)$$

Our goal is to prove that the equilibrium point $\bar{\mathbf{D}}_{P^*}$ corresponding to the *shortest* s - t path P^* is globally asymptotically stable for Physarum solver.

Remark. If G has h edges such that $L(P_1) = \cdots = L(P_h)$ ($h > 2$), the number of equilibria corresponding to P_1, \dots, P_h is uncountably infinite. Let $G' = P_1 \cup P_2 \cup \cdots \cup P_h$, $V(G') = \{v_{\gamma_0}, \dots, v_{\gamma_k}\}$, and $v_{\gamma_0} = s, v_{\gamma_k} = t$. Since $\dot{D}_{ij} = 0$ implies $Q_{ij} = D_{ij}$, the set of equilibria corresponding to the paths consists of all the point such that

$$\begin{cases} \sum_{j=1}^k D_{\gamma_0 \gamma_j} = \sum_{j=0}^{k-1} D_{\gamma_j \gamma_k} = I_0, \\ \sum_{j>i} D_{\gamma_i \gamma_j} = \sum_{j<i} D_{\gamma_i \gamma_j} & \text{if } i \neq 0, k. \end{cases}$$

3.3 Main Theorem

Here we prove the main theorem of this paper:

Theorem 1. *Assume that $N = (G, s, t, L)$ satisfies the following properties:*

1. G is a connected planar graph with $|G| \geq 2$.
2. The source s and the target t are on the same face of G .
3. G has exactly one shortest s - t path $P_* = v_{\alpha_0} \cdots v_{\alpha_k}$, where $v_{\alpha_0} = s, v_{\alpha_k} = t$.

Then the equilibrium point $\bar{\mathbf{D}}_{P_}$ corresponding to the path P_* is globally asymptotically stable for Physarum solver (1)-(3).*

First, the results in the previous subsections allow us to restrict the phase space of Physarum solver.

Lemma 1. *The hypercube*

$$H = \{\mathbf{D} | 0 \leq D_{ij} \leq I_0 \text{ for all } e_{ij} \in G\} \quad (17)$$

is attracting and invariant for Physarum solver.

Proof. According to the proposition 4, we get

$$-D_{ij} \leq \dot{D}_{ij} \leq I_0 - D_{ij}. \quad (18)$$

Hence any positive orbit is attracted to H . □ □

Therefore, we can always restrict the phase space into H . Hereafter we suppose that $\mathbf{D}(0) \in H$. The lower boundary of H is invariant for Physarum solver.

Lemma 2. *For any edge e_{ij} , the set $\{\mathbf{D} | D_{ij} = 0\}$ is invariant for Physarum solver. More generally, for a s - t path P the set*

$$\{\mathbf{D} | D_{ij} = 0 \text{ for } e_{ij} \notin P\} \quad (19)$$

is an invariant subset.

Proof. According to (3), if $D_{ij}(0) = 0$, then $\dot{D}_{ij}(\tau) = 0$ holds for all $\tau > 0$. □ □

Lemma 3. *Let G be a connected graph, then the following three hold.*

1. *If G has a vertex $v (\neq s, t)$ with $d(v) = 1$, then the conductivity of the incident edge tends to zero as $\tau \rightarrow \infty$.*

2. If G is s - t path, i.e. $G = P_*$, then $\bar{\mathbf{D}}_{P_*}$ is globally asymptotically stable.
3. If G is a tree, then $\bar{\mathbf{D}}_{P_*}$ is globally asymptotically stable

Proof. 1. At the edge e , the flux is always zero and we have $\dot{D} = -D$.
 2. At all the edge e , the flux is always equal to I_0 and we have $\dot{D} = I_0 - D$.
 3. It is clear from 1 and 2.

Note that if G is a path or a tree, the assumptions of our main theorem hold. □ □

Lemma 4. *Assume that G is not a path and satisfies the assumption (iii). Then there is no inner equilibrium point, that is, the interior of H contains no equilibrium point.*

Proof. If G is a tree, it is trivial. Assume that G contains at least two s - t paths. Suppose that the system has an inner equilibrium point $\bar{\mathbf{D}}_I$. It implies that for every $e_{ij} \in G$

$$|p_i - p_j| = L_{ij} \quad \text{and} \quad D_{ij} > 0.$$

For the shortest s - t path, we obtain

$$L(P_*) = \sum_{i=0}^k |p_{\alpha_i} - p_{\alpha_{i+1}}| \geq p_0.$$

On the other, for any s - t directed path $P = v_{\beta_0} \cdots v_{\beta_l}$, we obtain

$$p_0 = \sum_{j=0}^l (p_{\beta_j} - p_{\beta_{j+1}}) = L(P).$$

Thus we have $L(P_*) \geq L(P)$ and it is contradiction. □ □

Now, we introduce the main tool.

Definition 1. *For an edge $e_{ij} \in G$, we define a function $F_{ij}(\mathbf{D})$ as*

$$F_{ij}(\mathbf{D}) = L_{ij} \log D_{ij}. \quad (20)$$

For a path $P = v_{\beta_0} \cdots v_{\beta_k}$ in G , we define a function $F_P(\mathbf{D})$ as

$$F_P(\mathbf{D}) = \sum_{i=0}^{k-1} F_{\beta_i \beta_{i+1}}. \quad (21)$$

Lemma 5. *The derivative of F_{ij} is calculated as*

$$\dot{F}_{ij} = |p_i - p_j| - L_{ij}. \quad (22)$$

Therefore we obtain

$$\dot{F}_P = \sum_{i=0}^{k-1} |p_{\beta_i} - p_{\beta_{i+1}}| - L(P). \quad (23)$$

Proof. The proof is straight forward.

$$\begin{aligned} \dot{F}_{ij} &= \frac{L_{ij}}{D_{ij}} \dot{D}_{ij} = \frac{L_{ij}}{D_{ij}} (|Q_{ij}| - D_{ij}) \\ &= \frac{L_{ij}}{D_{ij}} \left(\frac{D_{ij}}{L_{ij}} |p_i - p_j| - D_{ij} \right) = |p_i - p_j| - L_{ij}. \end{aligned}$$

And (23) can be obtained immediately. \square \square

Lemma 6. *Assume that (G, s, t, L) satisfies the assumption (iii). Let $P = v_{\beta_0} \cdots v_{\beta_l} \subset G$ be a s - t path such that $P \neq P_*$. Assume that $Q_{\beta_i \beta_{i+1}}(\tau) > 0$ for all i and $\tau \geq 0$. Then any orbit from the interior of H is exponentially attracted onto the lower boundary of H . Especially, there exists at least one edge $e_{ij} \in P \setminus P_*$ such that $D_{ij} \rightarrow 0$ as $\tau \rightarrow \infty$. Therefore any orbit is exponentially attracted onto an invariant subset*

$$\{\mathbf{D} \in H \mid D_{ij} = D_{ji} = 0\}. \quad (24)$$

Proof. According to Lemma 5 and the triangle inequality, we get

$$\begin{aligned} \dot{F}_{P_*} - \dot{F}_P &= \sum_{i=0}^{k-1} |p_{\alpha_i} - p_{\alpha_{i+1}}| - L(P_*) - \sum_{i=0}^{l-1} (p_{\beta_i} - p_{\beta_{i+1}}) + L(P) \\ &\geq |p_0 - p_n| - (p_0 - p_n) + L(P) - L(P_*) \\ &= L(P) - L(P_*) > 0. \end{aligned}$$

Thus $F_{P_*} - F_P$ is continuous and strictly increasing with respect τ , except on the lower boundary of H . Hence $F_{P_*} - F_P$ is a Lyapunov function. Now, if we assume that the ω -limit set of $\mathbf{D}(0) > 0$ is in the interior of H , then it consists of equilibrium point. As there is no inner equilibrium point from Lemma 4, it is contradiction. Therefore the interior of H does not contain ω -limit set, and any orbit is attracted to the lower boundary of H . Especially, it follows that this attraction is exponential. The function $F_P - F_{P_*}$ is calculated as

$$F_P - F_{P_*} = \log \frac{\prod_{i=0}^{l-1} D_{\beta_i \beta_{i+1}}^{L_{\beta_i \beta_{i+1}}}}{\prod_{i=0}^{k-1} D_{\alpha_i \alpha_{i+1}}^{L_{\alpha_i \alpha_{i+1}}}}.$$

Let

$$W = \frac{\prod_{i=0}^{l-1} D_{\beta_i \beta_{i+1}}^{L_{\beta_i \beta_{i+1}}}}{\prod_{i=0}^{k-1} D_{\alpha_i \alpha_{i+1}}^{L_{\alpha_i \alpha_{i+1}}}}.$$

The derivative of W with respect to time is calculated as

$$\dot{W} = (\dot{F}_P - \dot{F}_{P_*})W \leq (L(P_*) - L(P))W. \quad (25)$$

Since W is positive in the interior of H and $L(P_*) < L(P)$, W exponentially tends to zero as $\tau \rightarrow +\infty$. In addition, $\dot{W} = 0$ holds if and only if $D_{ij} = 0$ for at least one edge $e_{ij} \in P \setminus P_*$. Therefore, D_{ij} exponentially tends to zero as $\tau \rightarrow \infty$. \square \square

According to Lemma 6, we can restrict the system into (24) to know the ω -limit set of Physarum solver for N . The reduced system corresponds to the Physarum solver on the graph $G - e_{ij}$. The proof of our main theorem is as follows:

Proof. Now we consider $N = (G, s, t, L)$ with the assumptions (i)-(iii). We can assume that s and t are on the boundary of outer-face Γ . If we consider \vec{G} by the direction of flow, \vec{G} is a directed acyclic graph. Since G is planar, Γ consists of two s - t directed paths. It implies that s - t path of Lemma 6 do exist.

Let $P \neq P_*$ be an s - t path on γ . According to Lemma 6, for at least one edge $e_{ij} \in P$, $D_{ij} \rightarrow 0$ as $\tau \rightarrow \infty$. Let $G_1 = G - e_{ij}$, then we can consider the Physarum solver for $N_1 = (G_1, s, t, L)$ as an attracting invariant manifold of that for N . Repeat such a reduction. At i th step, suppose that we obtain G_i . Lemma 6 implies that for a certain edge e on the boundary of outer-face of G_i the conductivity tends to zero as $\tau \rightarrow \infty$. Then let $G_{i+1} = G_i - e$ and we reduce the system to Physarum solver for $N_{i+1} = (G_{i+1}, s, t, L)$. If G_{i+1} has a vertex $v \neq s, t$ with $d(v) = 1$, then the conductivity of the edge incident with v tends to zero as $\tau \rightarrow \infty$. So replace G_{i+1} by $G_{i+1} - v$. As $G_{i+1} \subset G_i$, G_{i+1} is also planar and $P \subset G_{i+1}$. Since G is a finite graph, there exists a finite number f such that $G_f = P$. As the graph is an s - t path, then all edges survive from Lemma 3. Therefore \bar{D}_P is globally asymptotically stable. \square \square

4 Discussion

We have proved mathematically rigorously that Physarum solver must find the shortest path between two specified vertices on the same face of a given planar graph. In other words, we have mathematically rigorously

justified the ability of Physarum solver, by which we can find the minimum-length solution of a maze built on a petri dish. Thus, in this paper we have answered affirmatively completely for the problem presented in [13].

As telling from the mathematically technical point of view, it is a very interesting point that we prove the globally asymptotical stability without using usual Lyapunov function, which is defined globally in the whole domain. Instead of this, we use a kind of “local” Lyapunov function to operate a graph under consideration on which Physarum solver is defined.

As telling about an application of Physarum solver, in [15], they have attempted to apply Physarum solver to navigation system of road map. In fact, Physarum solver has found the shortest path connecting between Seattle and Houston admirably. Moreover, they say that Physarum solver has solved it in a little bit shorter time than using Dijkstra algorithm directly, which is usually used to get the shortest path on a graph. It is, therefore, meaningful very much from the viewpoint of application to verify that Physarum solver solves the shortest path decision problem mathematically rigorously.

参考文献

- [1] A. Berman, and R. J. Plemmons, *Nonnegative Matrices in the Mathematical Science, 2nd Edition*, SIAM, Philadelphia, PA, (1994).
- [2] R. Diestel, *Graph Theory*, Springer-Verlag, New York, (2000).
- [3] A. Eden, C. Foias, B. Nicolaenko and R. Temam, *Exponential Attractors for Dissipative Evolution Equations*, Wiley, 1994.
- [4] J. L. Gross, and J. Yellen, *Handbook of Graph Theory*, CRC PRESS, (2003).
- [5] T. Miyaji, and I. Ohnishi, Mathematical analysis to an adaptive network of the Plasmodium system, *Hokkaido Mathemaical Journal*, **36**, No. 2, pp. 445-465 (2007).
- [6] T. Miyaji, and I. Ohnishi, Can Physarum solve the shortest path problem mathematically rigorously ?, Research report in RIMS (2007).
- [7] T. Miyaji, I. Ohnishi, A. Tero, and T. Nakagaki, Failure to the shortest path decision of an adaptive transport network with double edges of Plasmodium system, To appear in *IJDSDE* (2008).

- [8] B. Mohar, and C. Thomassen, *Graphs on Surfaces*, The Johns Hopkins University Press, (2001).
- [9] T. Nakagaki, H. Yamada, and H. Hara, Smart network solution by an amoeboid organism, *Biophys. Chem.* **107** (2004), pp.1-5.
- [10] T. Nakagaki, H. Yamada, and Á. Tóth, Maze-solving by an amoeboid organism, *Nature* **407** (2000), pp.470.
- [11] T. Nakagaki, H. Yamada, and Á. Tóth, Path finding by tube morphogenesis in an amoeboid organism, *Biophys. Chem.* **92** (2001), pp.47-52.
- [12] A. Tero, R. Kobayashi, and T. Nakagaki, A coupled-oscillator model with a conservation law for the rhythmic amoeboid movements of plasmodial slime molds, *Physica D* **205** (2005), pp.125-135.
- [13] A. Tero, R. Kobayashi, and T. Nakagaki, A Mathematical model for adaptive transport network in path finding by true slime mold, *J.Theor.Biol.* **244** (2007), pp.553-564.
- [14] A. Tero, R. Kobayashi, and T. Nakagaki, Research Report in RIMS **1499** (2006), pp.159-166.
- [15] A. Tero, R. Kobayashi, and T. Nakagaki, Physarum solver: - a biologically inspired method of road-network navigation -, *Physica A.* **363** (2006), pp.115-119.