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<th>Title</th>
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</thead>
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The derivational complexity of string-rewriting systems

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1 Derivational complexity

Let $\Sigma$ be a (finite) alphabet and let $\Sigma^* = \cup_{n \geq 0} \Sigma^n$ be the free monoid generated by $\Sigma$. A (string)-rewriting system $R$ is a nonempty subset of $\Sigma^* \times \Sigma^*$. An element $r = (u, v)$ in $R$ is called a rule of $R$ and written $u \rightarrow v$. Suppose that a word $x \in \Sigma^*$ contains $u$ as a subword, that is, $x = x_1ux_2$ with $x_1, x_2 \in \Sigma^*$, then we can apply the rule $r$ to $x$ and $x$ is rewritten to the word $y = x_1vx_2$. In this situation we write as $x \rightarrow_r y$. If there is some rule $r \in R$ such that $x \rightarrow_r y$, we write $x \rightarrow_R y$, and we call the relation $\rightarrow_R$ the one-step derivation on $\Sigma^*$ by $R$.

A rewriting system $R$ is terminating on $x \in \Sigma^*$ if there is no infinite sequence of derivation:

$$x \rightarrow_R x_1 \rightarrow_R \cdots \rightarrow_R x_n \rightarrow_R \cdots$$

starting with $x$. $R$ is terminating (or noetherian), if it is terminating on every $x \in \Sigma^*$.

The maximal length of a derivation sequence starting with $x$ is denoted by $\delta_R(x)$. For $x$ on which $R$ is not terminating, we set $\delta_R(x) = \infty$. The function $d_R : \mathbb{N} \rightarrow \mathbb{N} \cup \{\infty\}$ defined by

$$d_R(n) = \max\{\delta_R(x) | x \in \Sigma^n\}$$

for $n \in \mathbb{N}$ is the derivational complexity of $R$.

We are interested in what functions can be derivational complexities of terminating finite rewriting systems.

Let $\mathbb{R}_+ = \{x \in \mathbb{R} | x \geq 0\}$. For two functions $f, g : \mathbb{N} \rightarrow \mathbb{R}_+ \cup \{\infty\}$, if there is a constant $C > 0$ such that $f(n) \leq C \cdot g(n)$ for any sufficiently large $n \in \mathbb{N}$, we write as $f \leq O(g)$. If moreover $g \leq O(f)$, $f$ and $g$ are called equivalent, and written as $f = O(g)$.

A function $f : \mathbb{N} \rightarrow \mathbb{R}_+ \cup \{\infty\}$ is super-additive if

$$f(m + n) \geq f(m) + f(n)$$

holds for any $m, n \in \mathbb{N}$. A super-additive function is non-decreasing. It is easy to see that the derivational complexity of a rewriting system is super-additive.
For an integer \(k \geq 1\), a rewriting system \(R\) has polynomial (derivational) complexity of degree \(k\), if \(d_R(n) = O(n^k)\). Any (nonempty) rewriting system \(R\) has at least linear complexity, that is, \(d_R(n) \geq O(n)\).

**Example 1.1.** Let \(k \geq 2\) and let \(\Sigma_k = \{a_1, a_2, \ldots, a_k\}\). For \(2 \leq \ell \leq k\) let

\[
C_\ell = \{a_1a_\ell \to a_\ell a_{\ell-1}, a_2a_\ell \to a_\ell a_{\ell-1}, \ldots, a_{\ell-1}a_\ell \to a_\ell a_1\}.
\]

Define a system \(P_k\) on \(\Sigma_k\) inductively as follows.

\[
P_2 = C_2 = \{a_1a_2 \to a_2a_1\},
\]

and

\[
P_k = P_{k-1} \cup C_k
\]

for \(k \geq 3\). Then, \(P_k\) has polynomial complexity of degree \(k\).

A rewriting system \(R\) has *exponential complexity*, if there are constants \(C \geq D > 1\) such that

\[
D^n \leq d_R(n) \leq C^n
\]

for sufficiently large \(n \in \mathbb{N}\). The one-rule system \(\{ab \to b^2a\}\) has an exponential derivational complexity.

Due to [4], a derivational complexity exists in each level of the Grzegorczyk hierarchy of primitive recursive functions. Even the Ackermann's function is attained ([5]). Actually, a derivational complexity can exceed any recursive function (see Section 2). Many studies have been done about the derivational complexity of term rewriting systems under specific termination techniques (see [7] and the references cited there). Here we shall discuss the derivational complexity of *string* rewriting systems under a general situation.

## 2 Q-systems and Turing machines

In this article we only consider deterministic Turing machines. Let

\[
M = M(\Sigma, Q, q_0, F, \delta)
\]

be a \(k\)-tape Turing machine, where \(\Sigma\) is a tape alphabet, \(Q\) is a set of states, \(q_0\) is an initial state, \(F\) is a set of final states and \(\delta\) is a transition function. We assume that the tapes are one-way infinite and each head never moves to the left of the initial position.

Let \(\Sigma_b = \Sigma \cup \{b\}\), where \(b\) denotes the blank symbol. The transition function \(\delta\) is a mapping from \((Q \setminus F) \times \Sigma_b^k\) to \(Q \times (\Sigma_b \cup \{L, R\})^k\), where \(L\) and \(R\) are the symbols for the right and left moves of the heads respectively. If for each \(i\) with \(1 \leq i \leq k\), \(x_iy_i\) is a word written on the \(i\)-th tape and the machine is looking at the leftmost letter of \(y_i\) in state \(q\), then the \(k\)-ple

\[
c = (x_1qy_1, x_2qy_2, \ldots, x_kqy_k)
\] (2.1)
is a configuration of $M$. The size $|c|$ of a configuration $c$ in \((2.1)\) is defined by

$$|c| = |x_1y_1x_2y_2\cdots x_ky_k|.$$  

For $x \in \Sigma^*$, let $\tau_M(x)$ be the number of steps taken until $M$ halts when it runs with input $x$ written in the first tape of $M$. The time function $t_M : \mathbb{N} \to \mathbb{N} \cup \{\infty\}$ of $M$ is defined by

$$t_M(n) = \max \{ \tau_M(x) \mid x \in \Sigma^n \}.$$  

For a configuration $c$, let $\tau'_M(c)$ be the number of steps taken until $M$ halts when it starts with $c$. In particular, $\tau_M(x) = \tau'_M(q_0x, q_0, \ldots, q_0)$ for $x \in \Sigma^*$. Define the total time function function $t'_M : \mathbb{N} \to \mathbb{N} \cup \{\infty\}$ of $M$ by

$$t'_M(n) = \max \{ \tau'_M(c) \mid c : \text{configuration of size } n \}.$$  

Clearly,

$$t'_M(n) \geq t_M(n)$$  

for any $n \in \mathbb{N}$.  

A Q-system is a finite rewriting system $R$ over an alphabet

$$\Sigma = Q \cup \Sigma_1 \cup \Sigma_2 \cup \{\$\} \quad \text{(disjoint union)}$$

consisting of rules only of the form

$$vqu \rightarrow v'q'u', \quad \text{or} \quad vqu\$ \rightarrow v'q'u'\$$

where $q, q' \in Q, u, u' \in \Sigma_1^*$ and $v, v' \in \Sigma_2^*$.  

A word $x \in \Sigma^*$ is admissible (resp. weakly admissible), if it is of the form $vqu$ with $q \in Q, v \in \Sigma_2^*$ and $u \in \Sigma_1^*$ (resp. $u \in \Sigma_1^* \cup \Sigma_1^*\$).  

For a Q-system $R$ and for $n \in \mathbb{N}$, define

$$ad_R(n) = \max \{ \delta_R(x) \mid x \text{ is admissible and } |x| = n + 2 \}$$

**Lemma 2.1.** For a Q-system $R$, we have

$$ad_R(n) \leq d_R(n + 2)$$  

for any $n \in \mathbb{N}$. If $ad_R$ is super-additive, then

$$d_R(n + 1) \leq ad_R(n)$$  

for any $n \in \mathbb{N}$. If $ad_R$ is equivalent to a non-zero super-additive function, then

$$d_R(n + 1) \leq O(ad_R(n)).$$
There is a natural way to simulate one-tape Turing machines by string-
rewriting systems ([3]).

Let $M = M(\Sigma, Q, q_0, F, \delta)$ be a one-tape Turing machine. Here, $\delta$ is a
mapping from $(Q \setminus F) \times \Sigma_b$ to $Q \times (\Sigma_b \cup \{L, R\})$. We define a Q-system $R_M$
associated with $M$ as follows. $R_M$ is a rewriting system on the alphabet

$$\Omega = Q \cup \Sigma_b \cup \overline{\Sigma}_b \cup \{$$ \}

where $\overline{\Sigma}_b = \{\overline{a} | a \in \Sigma_b\}$ is a copy of $\Sigma_b$,
and consists of the rules:

$$qa \rightarrow \overline{a}q', \quad \delta(q, a) = (q', R),$$

$$a'qa \rightarrow q'a'a, \quad \delta(q, a) = (q', L),$$

$$qa \rightarrow q'a', \quad \delta(q, a) = (q', a'),$$

$$q\$ \rightarrow \overline{b}q\$, \quad \delta(q, b) = (q', R),$$

$$\overline{a}q\$ \rightarrow qa\$, \quad \delta(q, b) = (q', L),$$

$$q\$ \rightarrow q'a\$, \quad \delta(q, b) = (q', a),$$

for $a, a' \in \Sigma_b$, $q \in Q \setminus F$ and $q' \in Q$.

For a word $x \in \Sigma_b^*$, $\overline{x}$ denotes the word obtained from $x$ by replacing every
letter $a$ in $x$ by $\overline{a}$. Since one step of the Turing machine $M$ just corresponds to
one rewriting by $R_M$ we have

**Lemma 2.2.** It holds that

$$\delta_{R_M}(q_0x\$) = \tau_M(x), \quad \delta_{R_M}((\overline{x})qy\$) = \tau_M'(xqy)$$

for $x, y \in \Sigma_b^*$ and $q \in Q$.

**Corollary 2.3.** We have

$$d_{R_M}(n + 2) \geq ad_{R_M}(n) = t'_M(n) \geq t_M(n)$$

for $n \geq 0$.

If $R$ is finite and terminating, then we can compute $d_R$ by tracing all the
derivation sequences (see Section 4), and it is a recursive function. Actually it
can exceed any recursive function.

**Corollary 2.4.** For any recursive function $f$, there exists a finite terminating
rewriting system $R$ such that

$$d_R(n) \geq f(n)$$

for any positive $n \in \mathbb{N}$.

### 3 Time functions and derivational complexity

As we have seen in the last section, derivational complexity is related to the
time functions of Turing machines.
Lemma 3.1. (cf. [2], [6]) For any k-tape Turing machine $M$ with time function $f(n) \geq O(n)$, there exists a one-tape Turing machine $M'$ such that $t_{M'}(n) = O(t_{M'}'(n)) = O(f(n)^2)$.

Suppose that $f$ is the time function of a k-tape Turing machine $M$ such that $f \geq O(n)$ and $f^2$ is equivalent to a super-additive function $g$. Let $M'$ be the one-tape Turing machine Lemma 3.1. We have

$$t_{M'}'(n) = O(f(n)^2) = O(g(n)).$$

Let $R$ be the $Q$-system associated with $M'$, then by Lemma 2.1 and Corollary 2.3, we see

$$d_{R}(n + 2) \geq t_{M'}'(n) = ad_{R}(n) \geq O(d_{R}(n + 1)).$$

It follows that

$$O(f(n - 2)^2) \leq d_{R}(n) \leq O(f(n - 1)^2).$$

Thus, we have

Theorem 3.2. Let $f(n)$ be a time function of a Turing machine such that $f \geq O(n)$ and $f(n)^2$ is equivalent to a super-additive function. Then there exists a finite rewriting system $R$ such that

$$O(f(n - 2)^2) \leq d_{R}(n) \leq O(f(n - 1)^2).$$

We say that a function $f : \mathbb{N} \rightarrow \mathbb{N}$ is computable in time $O(g(n))$, if there exists a (deterministic) algorithm computing $f(n)$ within time $O(g(n))$, more precisely, if there exists a multi-tape Turing machine which computes binary $f(n)$ for given binary $n$ with time function $t_{M}(n) \leq O(g(n))$.

Lemma 3.3. If $f : \mathbb{N} \rightarrow \mathbb{N}$ is a function such that $f(n) \geq O(n^2)$ and the binary $f(n)$ is computable in time $O(\sqrt{f(n)})$ for binary $n \in \mathbb{N}$, then $\lfloor \sqrt{f(n)} \rfloor$ is equivalent to a time function of a Turing machine.

Combining this lemma with Theorem 3.1 we have

Theorem 3.4. Suppose that a function $f(n) \geq O(n^2)$ is computable in time $O(\sqrt{f(n)})$ in binary and equivalent to a super-additive function. Then, there exists a finite rewriting system $R$ such that

$$O(f(n - 2)) \leq d_{R}(n) \leq O(f(n - 1)).$$

4 Computing the derivational complexity

Let $R$ be a rewriting system on $\Sigma$. Consider a derivation sequence of length 2:

$$x = x'ux'' \rightarrow_{R} x'vx''' = y = y'u'y'' \rightarrow_{R} y'y'' = z,$$

where $u \rightarrow v, u' \rightarrow v' \in R$. This sequence is left canonical, if

$$|x'| < |y'u'|.$$
Lemma 4.1. For a derivation sequence of length $n$ from $x \in \Sigma^*$ to $y \in \Sigma^*$, there is a left canonical sequence from $x$ to $y$ of the same length $n$.

For a derivational sequence

$$p : x_0 \rightarrow_R x_1 \rightarrow_R x_2 \rightarrow_R \cdots \rightarrow_R x_n,$$

we define a number $L(p)$ by induction on $n$ as follows. When $n = 1$ and

$$p : x_0 = x_0'ux_0'' \rightarrow_r x_0'vx_0''$$

with $r = (u \rightarrow v) \in R$, define

$$L(p) = |x_0'v| = |x_0| - |x_0''|.$$

Suppose that $n \geq 2$ and

$$x_{n-2} = x_{n-2}'u'x_{n-2}'' \rightarrow_{r'} x_{n-2}'v'x_{n-2}'' = x_{n-1} = x_{n-1}'ux_{n-1}'' \rightarrow_r x_{n-1}'vx_{n-1}'' = x_n$$

with $r = (u \rightarrow v), r' = (u' \rightarrow v') \in R$. Then, define

$$L(p) = L(p') + |x_{n-1}'| - |x_{n-2}'| + |u| + K - 1,$$

where $p'$ is the subsequence

$$x_0 \rightarrow_R x_1 \rightarrow_R \cdots \rightarrow_R x_{n-1}$$

of $p$ and

$$K = \max \{ |u|, |v| \mid u \rightarrow v \in R \}.$$

Lemma 4.2. For any derivation sequence $p$ of length $n \geq 1$ starting with $x \in \Sigma^*$ we have

$$L(p) \leq (2K - 1)(n - 1) + |x|.$$

Lemma 4.3. A left canonical derivation sequence $p$ can be found by tracing at most $L(p)$ letters in the words appearing in $p$.

Theorem 4.4. Let $R$ be a finite rewriting system on $\Sigma$ with derivational complexity $f$. Then, given $n \in \mathbb{N}$, $f(n)$ can be computed deterministically in time $C^{f(n)}$ for some constant $C > 1$.

5 Complexities of the forms $n^\alpha$ and $\alpha^n$

In this section we give the results that there are finite rewriting systems with derivational complexities equivalent to $n^\alpha$ (and $\alpha^n$), if the computational complexity of the real number $\alpha$ is relatively low, but there are no such systems if the complexity of $\alpha$ is high. The author has been inspired by the discussions in [8].

A real number $\alpha > 0$ is computable in time $f(n)$, if a binary rational approximation $a/b$ ($a, b \in \mathbb{N}$) of $\alpha$ such that $b \leq O(2^n)$ and

$$|\alpha - \frac{a}{b}| < \frac{1}{2^n}$$

can be computed in time $f(n)$ (refer to [9] for computable real numbers). We denotes this rational $a/b$ by $\alpha[n]$. 
Lemma 5.1. Let $\alpha > 0$ be a real number computable in time $O(f(n))$. Then for an integer $\nu$, the function $g_{\alpha,\nu}(n) = 2^\lfloor \alpha\lceil \log_2 n \rceil - \nu \rfloor \cdot n \rfloor$ is equivalent to $2^{\alpha n}$ and can be computed in time $O(f(\lceil \log_2 n \rceil - \nu) + n)$.

Theorem 5.2. Let $\alpha \geq 2$ be a real number computable in time $O(C^{2^n})$ for some constant $C > 1$. Then, there is a finite rewriting system $R$ with derivational complexity equivalent to $\alpha^n$.

Next, we consider the exponential function $\alpha^n$. Because it is not super-additive, we need the following

Lemma 5.3. Let $\alpha > 1$ be a real number, then the function $f_\alpha$ defined by

$$f_\alpha(n) = \begin{cases} \alpha^n & \text{if } n \geq 1/\log \alpha \\ (e \log \alpha) \cdot n & \text{if } 0 \leq n < 1/\log \alpha \end{cases}$$

is super-additive.

The computational complexities of $\alpha$ and $\log_2 \alpha$ are closely related.

Lemma 5.4. Let $\alpha (> 1)$ be a real number computable in time $O(f(n))$. Then, $\log_2 \alpha$ is computable in time $O(f(n + 2) + 4^n n^2)$, and $2^\alpha$ is computable in time $O(f(n + \lceil \alpha \rceil + 2) + 8^n n^2)$.

If we use a faster algorithm to compute the product of two integers, for example, Schönhag-Strassen's algorithm (see [1]), we can improve Lemma 5.4, but this is enough for our purpose.

Theorem 5.5. If a real number $\alpha > 1$ is computable in time $O(C^{2^n})$ for some constant $C > 1$, then there is a finite rewriting system $R$ with derivational complexity equivalent to $\alpha^n$.

By our results we see that, for example, the functions $n^\alpha (\alpha \geq 2)$, $\alpha^n (\alpha > 1)$ and $2^{\alpha n} (\alpha > 0)$ for a rational (or more generally an algebraic) number $\alpha$ are equivalent to the derivational complexities of finite rewriting systems. For a transcendental number $\alpha$ with low complexity such as $\pi$ and $e$, they are also equivalent to the derivational complexities.

Using Theorem 4.4, we can give the other direction as follows.

Theorem 5.6. Let $\alpha > 1$ be a real number.

(1) If there is a finite rewriting system with derivational complexity equivalent to $n^\alpha$, then $\alpha$ is computable in time $C^{2^n}$ for some constant $C > 1$.

(2) If there is a finite rewriting system with derivational complexity equivalent to $\alpha^n$, then $\alpha$ is computable in time $C^{2^n}$ for some constant $C > 1$. 
References


