Relative *t*-designs in Johnson association schemes for P-polynomial structure

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1 Introduction

Relative t-designs were defined in Q-polynomial and P-polynomial association schemes respectively by Delsarte [6] in 1977 and Bannai–Bannai–Suda–Tanaka [1] in 2015. Actually, relative t-designs in P-polynomial association schemes were first introduced by Delsarte-Seidel [7] in 1998 for binary Hamming association scheme H(n, 2). They called such designs regular t-wise balanced designs which are equivalent to relative t-designs in H(n, 2) for P-polynomial structure with respect to the fixed point $(0, 0, \ldots, 0)$. In 2015, Bannai-Bannai-Suda-Tanaka [1] proposed the definition for general P-polynomial association schemes and proved that the concepts of relative t-designs in Hamming association schemes H(n,q) $(q \ge 2)$ for Ppolynomial structure and Q-polynomial structure are equivalent. Using this good property, Bannai–Bannai–Zhu [2] proved a necessary and sufficient condition for relative t-designs on two shells of H(n,2). In addition, they proved that if (Y,w) is a relative t-design in H(n,2)on p shells then the subset of (Y, w) on each shell must be a usual (weighted) combinatorial (t+1-p)-design. It is an interesting question to ask how the situation is in the case of Johnson association scheme J(v, k). Each nontrivial shell X_r of J(v, k) is known to be a commutative association scheme which is the product of two smaller Johnson association schemes. Bannai-Zhu [4] studied relative t-designs on one shell X_r in J(v, k) for Q-polynomial structure and proved that they are \mathcal{T} -designs in X_r for $\mathcal{T} = \{(t_1, t_2) \mid 0 \leq t_1 + t_2 \leq t\}$. In particular, if the weight is constant, then relative t-designs on one shell in J(v, k) are mixed t-designs in $J(k, k_1) \otimes J(v-k, k_2)$ which were introduced and studied by Martin [8]. It is well known that Johnson association schemes are both P- and Q-polynomial association schemes. Therefore it would be interesting to ask the similar question for P-polynomial structure whether we can expect to regard relative t-designs on one shell of J(v, k) as weighted \mathcal{T} -designs in X_r for some set \mathcal{T} .

In this report, we investigate relative t-designs in J(v, k) for P-polynomial structure and give the answer to this question. The main result is that if (Y, w) is a relative t-design supported by one shell X_r in J(v, k) for P-polynomial structure, then (Y, w) is a weighted \mathcal{T} -design in X_r with $\mathcal{T} = \{(t_1, t_2) \mid 0 \leq t_1, t_2 \leq t\}$. We also discuss the existence problem of tight relative t-designs. We make an algorithm to construct tight relative 2-designs in one shell and obtain many examples. In addition, we can construct some tight relative 3-designs on one shell X_{2u} in J(8u, 4u) for integer $u \geq 1$.

2 Preliminaries

In this section, we will give the definition of t-designs in Q-polynomial association schemes introduced by Delsarte [5], designs in product of Q-polynomial association schemes by Martin [9] and relative t-designs in P-polynomial association schemes by Bannai–Bannai–Suda–Tanaka [1]. (Please refer to [3] for more information on P-polynomial or Q-polynomial association schemes.)

Throughout this report (Y, w) is assumed to be a weighted subset of X, namely, Y is a non-empty finite subset of X and $w: Y \longrightarrow \mathbb{R}_{>0}$.

2.1 Definition of *t*-designs

Definition 2.1 ([5, Theorem 3.10]). Let $\mathfrak{X} = (X, \{R_r\}_{r=0}^k)$ be a Q-polynomial association scheme with respect to the ordering E_0, E_1, \ldots, E_k . A weighted subset (Y, w) of X is called a weighted t-design in \mathfrak{X} if $E_j\chi_{(Y,w)} = 0$ for all $1 \leq j \leq t$, where $\chi_{(Y,w)}$ is the weighted characteristic vector of (Y, w) defined by

$$\chi_{(Y,w)} = \begin{cases} w(y), & \text{if } y \in Y, \\ 0, & \text{if } y \notin Y. \end{cases}$$

Definition 2.2. Let $X = {V \choose k}$ be the set of all k-subsets of V with |V| = v. A weighted subset (Y, w) of X is called a weighted t- (v, k, λ_t) design if for any $z \in {V \choose t}$ the following value

$$\sum_{y \in Y, z \subset y} w(y) = \lambda_t$$

is a constant depending only on t but not on the choice of z.

Remark 1. Delsarte [5, Theorem 4.7] proved that a *t*-design in Johnson association scheme J(v, k) for Q-polynomial structure, which is a weighted *t*-design with constant weight w = 1, is equivalent to a combinatorial t- (v, k, λ_t) design. We should remark that this result is also true for (non-constant) weighted *t*-designs in J(v, k) and weighted *t*- (v, k, λ_t) designs.

2.2 Designs in product of Q-polynomial association schemes

In this sunsection, we recall the concept of designs in product of Q-polynomial association schemes introduced by Martin [9]. For any positive integer k_i , let C_i be the totally ordered chain on $\{0, 1, \ldots, k\}$ and set $C = C_1 \times C_2$. Consider the poset (C, \trianglelefteq) defined by

$$\mathcal{C} = \{ \underline{\ell} = (\ell_1, \ell_2) \mid 0 \le \ell_i \le k_i, i = 1, 2 \}$$

with partial order $\underline{\ell'} \leq \underline{\ell}$ if $\ell'_1 \leq \ell_1$ and $\ell'_2 \leq \ell_2$. A subset \mathcal{T} of \mathcal{C} is called a *downset* in (\mathcal{C}, \leq) if $\underline{\ell} \in \mathcal{T}$ and $\underline{\ell'} \leq \underline{\ell}$ imply $\underline{\ell'} \in \mathcal{T}$. For any set $\mathcal{E} \subset \mathcal{C}$, denote

$$\mathcal{E} + \mathcal{E} := \{ (\ell_1 + \ell'_1, \ell_2 + \ell'_2) \mid (\ell_1, \ell_2), (\ell'_1, \ell'_2) \in \mathcal{E} \}.$$

Let $\mathfrak{X}^{(i)} = (X^{(i)}, \{R_r^{(i)}\}_{r=0}^{k_i})$ be a Q-polynomial association scheme and $E_0^{(i)}, E_1^{(i)}, \ldots, E_{k_i}^{(i)}$ its Q-polynomial ordering of primitive idempotents for i = 1, 2. For two points $x = (x_1, x_2)$ and $y = (y_1, y_2)$ from $X^{(1)} \times X^{(2)}$, define a relation $R_{(h_1, h_2)}$ on $X^{(1)} \times X^{(2)}$ by

$$(x,y) \in R_{(h_1,h_2)}$$
 if $(x_1,y_1) \in R_{h_1}^{(1)}, (x_2,y_2) \in R_{h_2}^{(2)}.$

Define $\mathfrak{X} = \mathfrak{X}^{(1)} \otimes \mathfrak{X}^{(2)} = (X^{(1)} \times X^{(2)}, \mathcal{R})$ as the product of these two association schemes, where the relations set \mathcal{R} is

$$\mathcal{R} = \{ R_{(h_1, h_2)} \mid 0 \le h_1 \le k_1, 0 \le h_2 \le k_2 \}.$$

In particular, if $\mathfrak{X} = J(v_1, k_1) \otimes J(v_2, k_2)$, then, for $x = (x_1, x_2), y = (y_1, y_2) \in \binom{V_1}{k_1} \times \binom{V_2}{k_2}$, $(x, y) \in R_{(h_1, h_2)}$ means $|x_1 \cap y_1| = k_1 - h_1$ and $|x_2 \cap y_2| = k_2 - h_2$. Moreover, for a set $Y \subset \binom{V_1}{k_1} \times \binom{V_2}{k_2}$, we define the *distance set* of Y by

$$A(Y) = \{(h_1, h_2) \mid (x, y) \in R_{(h_1, h_2)}, x, y \in Y, x \neq y\}.$$

Note that $(0,0) \notin A(Y)$.

Definition 2.3 ([9, Theorem 2.3]). Let \mathcal{T} be a downset of \mathcal{C} . The weighted subset (Y, w) of $X = X^{(1)} \times X^{(2)}$ is called a weighted \mathcal{T} -design in $\mathfrak{X} = \mathfrak{X}^{(1)} \otimes \mathfrak{X}^{(2)}$ if

$$(E_{t_1}^{(1)} \otimes E_{t_2}^{(2)}) \chi_{(Y,w)} = 0, \text{ for all } (t_1, t_2) \in \mathcal{T} \setminus \{(0,0)\}$$

In particular, if $\mathfrak{X} = J(v_1, k_1) \otimes J(v_2, k_2)$ with points set $X = \binom{V_1}{k_1} \times \binom{V_2}{k_2}$, then we have another equivalent definition of \mathcal{T} -designs given by Martin [9, Lemma 2.2]. Define a partial order \preceq on the set $\tilde{X} = \binom{V_1}{\leq k_1} \times \binom{V_2}{\leq k_2}$ by $z \preceq y$ if $z_1 \subseteq y_1$ and $z_2 \subseteq y_2$ for $z = (z_1, z_2)$ and $y = (y_1, y_2)$ in \tilde{X} .

Definition 2.4 ([8]). Let (Y, w) be a weighted subset of $\binom{V_1}{k_1} \times \binom{V_2}{k_2}$. The pair (Y, w) is called a weighted (t_1, t_2) - $(v_1, k_1, v_2, k_2, \lambda_{(t_1, t_2)})$ design if for any $(z_1, z_2) \in \binom{V_1}{t_1} \times \binom{V_2}{t_2}$, the following value

$$\sum_{\substack{(y_1, y_2) \in Y \\ z_1, z_2) \preceq (y_1, y_2)}} w(y_1, y_2) = \lambda_{(t_1, t_2)}$$

is a constant depending only on the pair (t_1, t_2) but not on the choice of (z_1, z_2) .

By Definition 2.4, we obtain the following lemma (see also [8]).

Lemma 2.5. Let (Y, w) be a weighted (t_1, t_2) - $(v_1, k_1, v_2, k_2, \lambda_{(t_1, t_2)})$ design. Then it is also a weighted (s_1, s_2) - $(v_1, k_1, v_2, k_2, \lambda_{(s_1, s_2)})$ design satisfying

$$\binom{v_1 - s_1}{t_1 - s_1} \binom{v_2 - s_2}{t_2 - s_2} \lambda_{(t_1, t_2)} = \binom{k_1 - s_1}{t_1 - s_1} \binom{k_2 - s_2}{t_2 - s_2} \lambda_{(s_1, s_2)},$$
(2.1)

whenever $0 \leq s_1 \leq t_1$ and $0 \leq s_2 \leq t_2$. In particular, $\lambda_{(0,0)} = |Y|$.

Lemma 2.5 implies that a weighted (t, t)- $(v_1, k_1, v_2, k_2, \lambda_{(t,t)})$ design is exactly a weighted \mathcal{T} -design in $J(v_1, k_1) \otimes J(v_2, k_2)$ with $\mathcal{T} = \{(t_1, t_2) \mid 0 \leq t_1, t_2 \leq t\}.$

2.3 Relative *t*-designs in P-polynomial association schemes

Now let us give the definition of relative t-designs in P-polynomial association schemes introduced by Bannai–Bannai–Suda–Tanaka [1]. Let $\mathfrak{X} = (X, \{R_r\}_{r=0}^k)$ be a P-polynomial association scheme and u_0 a fixed point in X. For $0 \leq r \leq k$, we call the subset $X_r := \{x \in X \mid (x, u_0) \in R_r\}$ the r-th shell of \mathfrak{X} . Let $\mathbb{R}^{|X|}$ be the space consisting of column vectors indexed by the points in X. Given a point $z \in X_j$, define a vector $f_z \in \mathbb{R}^{|X|}$ whose x-th entry is

$$f_z(x) = \begin{cases} 1 & \text{if } x \in X_i, (x, z) \in R_{i-j}, i \ge j, \\ 0 & \text{otherwise.} \end{cases}$$

Let $\operatorname{Hom}_j(X) := \operatorname{span}\{f_z \mid z \in X_j\}$ for $0 \le j \le k$, then

 $\mathbb{R}^{|X|} = \operatorname{Hom}_0(X) + \operatorname{Hom}_1(X) + \dots + \operatorname{Hom}_k(X).$

We say Y is a subset of X on p shells $X_{r_1} \cup \cdots \cup X_{r_p}$ if $\{r \mid Y \cap X_r \neq \emptyset\} = \{r_1, r_2, \ldots, r_p\}$. Denote $Y_{r_{\nu}} = Y \cap X_{r_{\nu}}$ for $1 \leq \nu \leq p$.

Definition 2.6 ([1, Definition 1.1]). Let $\mathfrak{X} = (X, \{R_r\}_{r=0}^k)$ be a P-polynomial association scheme and u_0 a fixed point in X. A weighted subset (Y, w) is called a relative t-design on p shells $X_{r_1} \cup \cdots \cup X_{r_p}$ in \mathfrak{X} with respect to u_0 if the following

$$\sum_{\nu=1}^{p} \frac{W(Y_{r_{\nu}})}{|X_{r_{\nu}}|} \sum_{x \in X_{r_{\nu}}} f(x) = \sum_{y \in Y} w(y)f(y)$$
(2.2)

holds for any $f \in \operatorname{Hom}_0(X) + \operatorname{Hom}_1(X) + \dots + \operatorname{Hom}_t(X)$, where $W(Y_{r_\nu}) = \sum_{y \in Y_{r_\nu}} w(y)$.

3 Relative *t*-designs in one shell of J(v, k)

In this section, we discuss relative t-designs on one shell of Johnson association scheme J(v,k) for P-polynomial structure. We first describe the structure of each nontrivial shell of J(v,k). Let $X = \binom{V}{k}$ be the set of all k-subsets of $V = \{1, \ldots, v\}$ with $v \ge 2k$ and $J(v,k) = (X, \{R_r\}_{r=0}^k)$ Johnson association scheme. For any fixed point $u_0 \in X$, denote $X_r := \{x \in X \mid |x \cap u_0| = k - r\}$. Without loss of generality, we may assume $u_0 = \{1, 2, \ldots, k\}$. Let $V_1 = u_0, V_2 = V \setminus u_0$ and denote $v_i = |V_i|$ for i = 1, 2.

Proposition 3.1. Each nontrivial shell X_r of J(v, k) is identified with the product of two smaller Johnson association schemes $J(k, k_1) \otimes J(v - k, k_2)$. More explicitly,

- (1) $(X_r, \{R_h\}_{h=0}^k) \cong J(k, r) \otimes J(v k, r) \text{ if } 1 \le r \le \frac{k}{2}.$
- (2) $(X_r, \{R_h\}_{h=0}^k) \cong J(k, k-r) \otimes J(v-k, r)$ if $\frac{k}{2} < r \le \frac{v-k}{2}$.
- (3) $(X_r, \{R_h\}_{h=0}^k) \cong J(k, k-r) \otimes J(v-k, v-k-r)$ if $\frac{v-k}{2} < r \le k$.

Proof. For cases (1), (2) and (3), define the bijection $\phi: X_r \longrightarrow {\binom{V_1}{k_1}} \times {\binom{V_2}{k_2}}$ respectively by

$$\phi(x) = (V_1 \setminus x, V_2 \cap x), \quad (V_1 \cap x, V_2 \cap x), \quad (V_1 \setminus x, V_2 \setminus x)$$

We can check that, for $x, y \in X_r$, $(x, y) \in R_{h_1+h_2}$ if and only if $(\phi(x), \phi(y)) \in R_{(h_1,h_2)}$. Therefore $(X_r, \{R_h\}_{h=0}^k) \cong J(k, k_1) \otimes J(v-k, k_2)$. **Remark 2.** Without confusion, we identify X_r with the product association scheme $J(k, k_1) \otimes J(v - k, k_2)$ whenever we mention designs in X_r . In addition, if (Y, w) is a weighted subset of one shell X_r of J(v, k), then the above map ϕ preserves the weight as well, namely, $w(y) = w(\phi(y))$ for any $y \in Y$.

Now we are ready to give our main theorem. By definition 2.6, if p = 1, i.e., $Y \subset X_r$, then we have the specific condition for relative *t*-designs supported by one shell X_r as follows.

$$\frac{W(Y_r)}{|X_r|} \sum_{x \in X_r} f_z(x) = \sum_{y \in Y} w(y) f_z(y), \quad \forall z \in X_0 \cup X_1 \cup \dots \cup X_t.$$
(3.1)

Theorem 3.2. If (Y, w) is a relative t-design in J(v, k) supported by one shell X_r , then (Y, w) is a weighted \mathcal{T} -design in X_r with $\mathcal{T} = \{(t_1, t_2) \mid 0 \leq t_1, t_2 \leq t\}.$

Using Lemma 2.5, it is enough to prove that (Y, w) is a (t, t)- $(v_1, k_1, v_2, k_2, \lambda_{(t,t)})$ design for some constant $\lambda_{(t,t)}$ if $(X_r, \{R_h\}_{h=0}^k) \cong J(v_1, k_1) \otimes J(v_2, k_2)$. More precisely, we need to prove that $\sum_{\substack{y \in \phi(Y) \\ (z_1, z_2) \leq y}} w(y)$ is constant for any $(z_1, z_2) \in \binom{V_1}{t} \times \binom{V_2}{t}$.

Remark 3. In Theorem 3.2, (Y, w) is a weighted \mathcal{T} -design in X_r means $(\phi(Y), w)$ is a weighted \mathcal{T} -design in $J(k, k_1) \otimes J(v - k, k_2)$, where ϕ is the bijection defined in the proof of Proposition 3.1.

4 Lower bound for relative *t*-designs on one shell

In this section, we give the lower bound for relative t-designs on one shell X_r of J(v, k). By Theorem 3.2, it is equivalent to obtain the lower bound for weighted \mathcal{T} -designs in product association scheme $J(k, k_1) \otimes J(v - k, k_2)$ with $\mathcal{T} = \{(t_1, t_2) \mid 0 \leq t_1, t_2 \leq t\}$. The following lower bound of designs in product association schemes was proved by Martin [8].

Lemma 4.1 ([8, Theorem 3.2]). Let \mathcal{T} be a downset in $(\mathcal{C}, \trianglelefteq)$ and \mathcal{E} a set satisfying $(\mathcal{E} + \mathcal{E}) \cap \mathcal{C} \subseteq \mathcal{T}$. If Y is a \mathcal{T} -design in $J(v_1, k_1) \otimes J(v_2, k_2)$, then

$$|Y| \ge \sum_{(j_1,j_2)\in\mathcal{E}} \left[\binom{v_1}{j_1} - \binom{v_1}{j_1-1} \right] \cdot \left[\binom{v_2}{j_2} - \binom{v_2}{j_2-1} \right].$$
(4.1)

Moreover, if equality holds, then for any $(h_1, h_2) \in A(Y)$ we have

$$\sum_{(j_1,j_2)\in\mathcal{E}} Q_{j_1}^{(1)}(h_1) Q_{j_2}^{(2)}(h_2) = 0,$$

where $Q_{j_i}^{(i)}(x)$ is the following Hahn polynomial corresponding to $J(v_i, k_i)$

$$Q_{j}^{(i)}(x) = \left(\begin{pmatrix} v_{i} \\ j \end{pmatrix} - \begin{pmatrix} v_{i} \\ j-1 \end{pmatrix} \right) {}_{3}F_{2} \begin{pmatrix} -x, -j, -v_{i} - 1 + j \\ -k_{i}, -v_{i} + k_{i} \end{cases}; 1 \right).$$

From the proof of Theorem 3.2 in [8], one can check that inequality (4.1) also holds for weighted \mathcal{T} -designs in $J(v_1, k_1) \otimes J(v_2, k_2)$.

Corollary 4.2. If (Y, w) is a relative 2*e*-design in J(v, k) on one shell X_r with *P*-polynomial structure, then

$$|Y| \ge \binom{k}{e} \binom{v-k}{e}.$$
(4.2)

It follows from Theorem 3.2 that (Y, w) is a weighted \mathcal{T} -design for $\mathcal{T} = \{(t_1, t_2) \mid 0 \leq t_1, t_2 \leq 2e\}$. Take $\mathcal{E} = \{(j_1, j_2) \mid 0 \leq j_1, j_2 \leq e\}$. Using Lemma 4.1, we obtain

$$|Y| \ge \sum_{j_1=0}^{e} \sum_{j_2=0}^{e} \left[\binom{k}{j_1} - \binom{k}{j_1-1} \right] \cdot \left[\binom{v-k}{j_2} - \binom{v-k}{j_2-1} \right] = \binom{k}{e} \binom{v-k}{e}$$

Proposition 4.3. If (Y, w) is a relative (2e + 1)-design in J(v, k) on one shell X_r with *P*-polynomial structure, then

$$|Y| \ge 4 \binom{k-1}{e} \binom{v-k-1}{e}.$$
(4.3)

A relative t-design on one shell of J(v, k) is called *tight* if equality holds in (4.2) or (4.3).

At the end of this subsection, we introduce the concept of projections for designs which will be used later. Given a (t_1, t_2) - $(v_1, k_1, v_2, k_2, \lambda_{(t_1, t_2)})$ design Y, define the left and right projection of Y as follows.

$$\begin{split} Y^{(L)} &= \{y^{(L)} \mid (y^{(L)}, y^{(R)}) \in Y\}.\\ Y^{(R)} &= \{y^{(R)} \mid (y^{(L)}, y^{(R)}) \in Y\}. \end{split}$$

Then $Y^{(L)}$ are $t_1 - (v_1, k_1, \lambda_{(t_1, 0)})$ designs and $Y^{(R)}$ are $t_2 - (v_2, k_2, \lambda_{(0, t_2)})$ designs.

5 Tight relative 2-designs

In this section, we give an algorithm to construct tight relative 2-designs (with constant weight) on one shell of J(v, k). We also provide two explicit examples.

Assume the weight function is constant, i.e., w = 1. Let (Y, 1) be a tight relative 2-design on one shell X_r of J(v, k). Using Lemma 4.1, we obtain the distance set A(Y) from the zeros of following equation.

$$F(h_1, h_2) = \sum_{\underline{j} \in \mathcal{E}} Q_{\underline{j}}(\underline{h}) = 1 + Q_1^{(1)}(h_1) + Q_1^{(2)}(h_2) + Q_1^{(1)}(h_1)Q_1^{(2)}(h_2),$$

where $\mathcal{E} = \{(0,0), (0,1), (1,0), (1,1)\}$. More explicitly, we have

$$F(h_1, h_2) = \frac{v_1 v_2 \left(v_1 h_1 - h_1 - v_1 k_1 + k_1^2\right) \left(h_2 v_2 - h_2 - k_2 v_2 + k_2^2\right)}{k_1 k_2 \left(v_1 - k_1\right) \left(v_2 - k_2\right)}.$$
(5.1)

Then we know that

$$A(Y) \subseteq \left\{ \left(\frac{k_1(v_1 - k_1)}{v_1 - 1}, h_2\right), \left(h_1, \frac{k_2(v_2 - k_2)}{v_2 - 1}\right) \mid 0 \le h_1 \le k_1, 0 \le h_2 \le k_2, (h_1, h_2) \ne (0, 0) \right\}$$

The construction of explicit tight relative 2-designs on one shell X_r in J(v, k) (i.e. tight (2, 2)- $(v_1, k_1, v_2, k_2, \lambda_{(2,2)})$ designs) Y is equivalent to the following problem.

Problem. For a point set $X = X^{(L)} \times X^{(R)} = \binom{V_1}{k_1} \times \binom{V_2}{k_2}$, find a pair of 2- $(v_1, k_1, \lambda_{(2,0)})$ design $Y^{(L)} = \{y_1^{(L)}, \ldots, y_{\lambda_{(0,0)}}^{(L)}\} \subset X^{(L)}$ and 2- $(v_2, k_2, \lambda_{(0,2)})$ design $Y^{(R)} = \{y_1^{(R)}, \ldots, y_{\lambda_{(0,0)}}^{(R)}\} \subset X^{(R)}$ with the same cardinality, so that $Y = \{(y_i^{(L)}, y_i^{(R)}) \mid 1 \leq i \leq \lambda_{(0,0)}\} \subset Y^{(L)} \times Y^{(R)}$ satisfies the condition that, for each $\{p_1, q_1\} \subset V_1$ and $\{p_2, q_2\} \subset V_2$, there exists exactly $\lambda_{(2,2)}$ element(s) $y \in Y$ such that $(\{p_1, q_1\}, \{p_2, q_2\}) \preceq y$. The partial order \preceq means $\{p_1, q_1\} \subset y^{(L)}$ and $\{p_2, q_2\} \subset y^{(R)}$. Moreover, the parameters $\lambda_{(0,0)}, \lambda_{(2,0)}, \lambda_{(0,2)}$ and $\lambda_{(2,2)}$ satisfy the relation given by Eq. (2.1).

5.1 Basic idea for the construction

Now we explain the basic idea of an algorithm to construct (2, 2)- $(v_1, k_1, v_2, k_2, \lambda_{(2,2)})$ designs for a given 2- $(v_1, k_1, \lambda_{(2,0)})$ design as the left projection $Y^{(L)}$. The algorithm is applicable for cases when the size of 2- $(v_2, k_2, \lambda_{(2,2)})$ design is equal to $\lambda_{(2,0)}$. Define $\tilde{Y}^{(L)} = (y_1^{(L)}, y_2^{(L)}, \dots, y_{\lambda_{(0,0)}}^{(L)})$ and $\tilde{Y}^{(R)} = (y_1^{(R)}, y_2^{(R)}, \dots, y_{\lambda_{(0,0)}}^{(R)})$. We introduce the notation of $\rho(p,q;z)$ for a given $z = (z_1, z_2, \dots, z_{\lambda_{(0,0)}})$ and $p, q \in V_1$ $(p \neq q)$, defined as

$$\rho(p,q;z) = (z_{i_1}, z_{i_2}, \dots, z_{i_{\lambda_{(2,0)}}}) \text{ for } \{i_1, i_2, \dots, i_{\lambda_{(2,0)}}\} = \{i \mid \{p,q\} \subset y_i^{(L)} \in \tilde{Y}^{(L)}\}.$$

We observe that $\rho\left(p,q;\tilde{Y}^{(R)}\right)$ is a 2- $(v_2,k_2,\lambda_{(2,2)})$ design. This will be the key fact in the algorithm. We use this fact as a condition to fix the right (ordered) projection $\tilde{Y}^{(R)}$.

The essential idea of the algorithm is the following. We choose one 2- $(v_1, k_1, \lambda_{(2,0)})$ design as the left (ordered) projection $Y^{(L)}$. The right (ordered) projection $\tilde{Y}^{(R)}$ is unknown. Denote $c = (\emptyset, \ldots, \emptyset)$ of length $\lambda_{(0,0)}$. Since $\rho(p, q; \tilde{Y}^{(R)})$ for $\{p, q\} \subset V_1$ is a 2- $(v_2, k_2, \lambda_{(2,2)})$ design, we assign one 2- $(v_2, k_2, \lambda_{(2,2)})$ design to $\rho(p, q; c)$. We do this for all choices of $\{p, q\} \subset V_1$. If we find a c that is consistent with all conditions from $\rho(p, q; c)$, then we take it as $\tilde{Y}^{(R)}$.

5.2 Examples for tight relative 2-designs

Example 1. One trivial example for tight relative 2-design (i.e., (2,2)- $(v_1, k_1, v_2, k_2, \lambda_{(2,2)})$ design) Y is the product of a symmetric 2- (v_1, k_1, λ_1) design $Y^{(L)}$ and a symmetric 2- (v_2, k_2, λ_2) design $Y^{(R)}$ with the same cardinality and $\lambda_1 \cdot \lambda_2 = \lambda_{(2,2)}$. Namely,

$$Y = \{(y_i, y_j) \mid y_i \in Y^{(L)}, y_j \in Y^{(R)}, 1 \le i \le v_1, 1 \le j \le v_2\}.$$

Example 2. Tight relative 2-designs on X_3 in J(14,7), i.e., tight (2,2)-(7,3,7,3,1) designs.

Using the algorithm, we obtain many tight (2, 2)-(7, 3, 7, 3, 1) designs. Since both the left and right projections of a (2, 2)-(7, 3, 7, 3, 1) design are 2-(7, 3, 7) designs, take the second 2-(7, 3, 7) design on Spence's homepage [10] as the left projection. We give two examples which is constructed from that 2-(7, 3, 7) design. Let

$$Y_1 = \{(y_{i_1}, y_{j_1}), (y_{i_2}, y_{j_2}), (y_{10}, y_{j_1}) \mid i_\ell \in I_\ell, j_\ell \in J_\ell, \ell = 1, 2\}, Y_2 = \{(y_{i_1}, y_{j_1}), (y_{i_2}, y_{j_2}), (y_{10}, y_{j_3})) \mid i_\ell \in I_\ell, j_m \in J_m, \ell = 1, 2, m = 1, 1, 3.\}$$

where $I_1 = \{1, 2, 3, 4, 7, 9\}$, $I_2 = \{5, 6, 8, 11\}$, $J_1 = \{1, 2, 3, 4, 7, 9\}$, $J_2 = \{10\}$, $J_3 = \{1, 12, 13, 14, 15, 9\}$. Here y_i denotes the *i*-th block of the list given below.

 $\{1, 2, 3\}, \{1, 4, 5\}, \{1, 6, 7\}, \{2, 4, 6\}, \{2, 4, 7\}, \{2, 5, 6\}, \{2, 5, 7\}, \{3, 4, 6\}, \{3, 4, 7\}, \{3, 5, 6\}, \\ \{3, 5, 7\}, \{1, 4, 6\}, \{1, 5, 7\}, \{2, 4, 5\}, \{2, 6, 7\}.$

Then both Y_1 and Y_2 are tight (2, 2)-(7, 3, 7, 3, 1) designs with the following distance sets.

$$A(Y_1) = \{(0,2), (2,0), (2,2), (1,2), (3,2)\},\$$

$$A(Y_2) = \{(0,2), (2,0), (2,2), (2,1), (2,3), (1,2), (3,2)\}$$

Example 3. A tight relative 2-design on X_4 in J(26, 13), i.e., tight (2, 2)-(13, 4, 13, 4, 1) design.

Both the left and right projections of a (2, 2)-(13, 4, 13, 4, 1) design are 2-(13, 4, 13) designs. It is known that there is a unique symmetric 2-(13, 4, 1)-design \mathcal{D}_{13} up to automorphism. Take twelve copies of \mathcal{D}_{13} and one more design by permuting the points of \mathcal{D}_{13} labelled by 12 and 13, then we have a 2-(13, 4, 13) design as the left projection $Y^{(L)}$. Using the algorithm, we construct a tight (2, 2)-(13, 4, 13, 4, 1) design Y in the following.

$$Y = \{ (y_{i_1}, y_{j_1}), (y_{i_2}, y_{j_2}), (y_{i_3}, y_{18}) \mid i_{\ell} \in I_{\ell}, j_m \in J_m, \ell = 1, 2, 3, m = 1, 2 \}$$

where $I_1 = \{1, 4, 7, 8, 10, 22, 30\}$, $I_2 = \{13, 16, 20, 23, 26, 29\}$, $I_3 = \{14, 15, 19, 24, 27, 28\}$, $J_1 = \{1, 4, 7, 8, 10, 13, 16, 20, 22, 23, 26, 29, 30\}$, $J_2 = \{2, 3, 5, 6, 9, 11, 12, 17, 21, 25, 31, 32\}$. Here y_i denotes the *i*-th block of the list given below

$$\begin{split} &\{1,2,3,4\}, \{1,2,3,13\}, \{1,4,7,12\}, \{1,5,6,7\}, \{1,5,9,11\}, \{1,6,8,10\}, \{1,8,9,10\}, \\ &\{1,11,12,13\}, \{2,4,9,10\}, \{2,5,8,11\}, \{2,5,8,12\}, \{2,6,7,11\}, \{2,6,9,12\}, \{2,6,9,13\}, \\ &\{2,7,10,12\}, \{2,7,10,13\}, \{3,4,8,11\}, \{3,5,7,10\}, \{3,5,9,12\}, \{3,5,9,13\}, \{3,6,9,12\}, \\ &\{3,6,10,11\}, \{3,7,8,12\}, \{3,7,8,13\}, \{4,5,6,13\}, \{4,5,10,12\}, \{4,5,10,13\}, \{4,6,8,12\}, \\ &\{4,6,8,13\}, \{4,7,9,11\}, \{7,8,9,13\}, \{10,11,12,13\}. \end{split}$$

The distance set of Y equals

$$A(Y) = \{(0,3), (1,3), (2,3), (3,0), (3,1), (3,2), (3,3), (3,4), (4,3)\}.$$

5.3 Tight relative 3-designs on one shell of J(v, k)

In this subsection, we will discuss the existence problem of tight relative 3-designs (Y, w) in J(v, k) on one shell with constant weight, i.e., w = 1.

Let $V_1 = u_0, V_2 = V \setminus u_0$. We introduce the notation

$$Y_{(W_1,W_2)}^{(U_1,U_2)} = \{ (y_1 - W_1, y_2 - W_2) \mid (y_1, y_2) \in \phi(Y), W_i \subset y_i \subset V_i - U_i \},\$$

where W_i and U_i are subsets of V_i such that $W_i \cap U_i = \emptyset$ for i = 1, 2. Using the proof of Proposition 4.3, if there exists a tight relative 3-design Y on one shell X_r of J(v, k), then we have four tight $\{(2,2)\}$ -designs $Y_{(\alpha,\beta)}^{(\emptyset,\emptyset)}$, $Y_{(\alpha,\emptyset)}^{(\emptyset,\beta)}$, $Y_{(\emptyset,\beta)}^{(\alpha,\emptyset)}$ and $Y_{(\emptyset,\emptyset)}^{(\alpha,\beta)}$ respectively in product association schemes $J(v'_1, k_1 - 1) \otimes J(v'_2, k_2 - 1)$, $J(v'_1, k_1 - 1) \otimes J(v'_2, k_2)$, $J(v'_1, k_1) \otimes J(v'_2, k_2 - 1)$ and $J(v'_1, k_1) \otimes J(v'_2, k_2)$, where $v'_1 = k - 1$, $v'_2 = v - k - 1$ and

$$(k_1, k_2) = \begin{cases} (r, r), & \text{if } 1 \le r \le \frac{k}{2}, \\ (k - r, r), & \text{if } \frac{k}{2} < r \le \frac{v - k}{2}, \\ (k - r, v - k - r), & \text{if } \frac{v - k}{2} < r \le k - 1. \end{cases}$$

According to the computer search, there is only one family of possible parameters, i.e., (v, k, r) = (8u, 4u, 2u) when $v \leq 3000$.

We can construct a tight relative 3-design on one shell X_{2u} in J(8u, 4u). Choose any tight relative 2-design on one shell X_{2u-1} in J(8u-2, 4u-1) as $Y_{(\alpha,\beta)}^{(\emptyset,\emptyset)}$. Replacing the left (resp. right) projection of $Y_{(\alpha,\beta)}^{(\emptyset,\emptyset)}$ by its complementary design, then this new design is $Y_{(\alpha,\emptyset)}^{(\emptyset,\beta)}$ (resp. $Y_{(\emptyset,\beta)}^{(\alpha,\emptyset)}$). Replace both the left and right projection of $Y_{(\alpha,\beta)}^{(\emptyset,\emptyset)}$ by their complementary designs and denote this new design as $Y_{(\emptyset,\beta)}^{(\alpha,\beta)}$. Define the sets Y_1, Y_2, Y_3, Y_4 as follows.

$$\begin{split} Y_1 &= \left\{ (Y' \cup \{\alpha\}, Y'' \cup \{\beta\}) \mid (Y', Y'') \in Y_{(\alpha,\beta)}^{(\emptyset,\emptyset)} \right\}, \\ Y_2 &= \left\{ (Y' \cup \{\alpha\}, Y'') \mid (Y', Y'') \in Y_{(\alpha,\emptyset)}^{(\emptyset,\beta)} \right\}, \\ Y_3 &= \left\{ (Y', Y'' \cup \{\beta\}) \mid (Y', Y'') \in Y_{(\emptyset,\beta)}^{(\alpha,\emptyset)} \right\}, \\ Y_4 &= Y_{(\emptyset,\emptyset)}^{(\alpha,\beta)}. \end{split}$$

Then $Y = Y_1 \cup Y_2 \cup Y_3 \cup Y_4$ is a tight relative 3-design on the shell X_{2u} in J(8u, 4u).

5.4 Further problem

According to the computer search, if $r \leq \frac{k}{2}$, then we get the following table for the possible parameters of tight relative 2-designs on one shell X_r in J(v,k), (i.e., tight (2,2)- $(k, r, v - k, r, \lambda_{(2,2)})$ designs).

| v | 14 | 22 | 26 | 29 | 30 | 32 | 32 | 37 | 37 | 38 | |
|-------------------|----|----|----|----|----|----|----|----|----|----|--|
| k | 7 | 11 | 13 | 13 | 15 | 11 | 16 | 15 | 16 | 19 | |
| r | 3 | 5 | 4 | 6 | 7 | 5 | 6 | 7 | 6 | 9 | |
| $\lambda_{(2,2)}$ | 1 | 4 | 1 | 5 | 9 | 2 | 4 | 6 | 3 | 16 | |

There are two interesting two families in the table above, namely, $(k, r, \lambda_{(2,2)}) = (u^2 + u + 1, u + 1, 1)$ and $(k, r, \lambda_{(2,2)}) = (4u - 1, 2u - 1, (u - 1)^2)$ for integer $u \ge 2$. We obtained some non-trivial examples for the first family when u = 2, 3, 4, 5. **Problem one** is the general construction for tight relative 2-designs with the possible parameters from the above two families.

According to the computer search for $v \leq 3000$, there is only one family of parameters (v, k, r) = (8u, 4u, 2u) for tight relative 3-designs and we can construct such designs if the tight relative 2-designs involved exist. **Problem two** is whether all the tight relative 3-designs on one shell of J(v, k) have the parameter (v, k, r) = (8u, 4u, 2u).

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