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Representation of successor-type proof-theoretically regular ordinals via limits

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Abstract
In this paper, we extend a result in [Ta04], that is, we show that every successor-type proof-theoretically regular ordinal has its own representation as a limit of a sequence consisting of certain canonical elements.

1 Introduction
In our previous paper [Ta04], we defined a set $\text{Reg}(\mathcal{T}(M))$ based on $\mathcal{T}(M)$, which was a primitive recursive well-ordered set defined by M. Rathjen to establish the proof theoretic ordinal of KRM. We call elements of $\text{Reg}(\mathcal{T}(M))$ "proof-theoretically regular ordinals based on $\mathcal{T}(M)$ (ptros)". In [Ta04], we also characterized some sort of ptros as proof-theoretical analogues of (hyper) inaccessible cardinals up to the least Mahlo cardinal. Since the characterization is based on $\text{Reg}(\mathcal{T}(M))$ as an analogue of the set of regular cardinals up to the least Mahlo cardinal, it is significant to characterize ptros and find the relationship between $\text{Reg}(\mathcal{T}(M))$ and the set of regular cardinals up to the least Mahlo cardinal. For these purpose, we are in the process of establishing a "canonical" fundamental sequence of each limit-type element of $\mathcal{T}(M)$. A coherent way to establish an appropriate fundamental sequence of each limit-type element of $\mathcal{T}(M)$ can be expected to be a coherent way to re-construct each element of $\mathcal{T}(M)$ as a more familiar concept, and hence, it turns out to provide a desirable characterization of ptros as proof-theoretical analogues of regular cardinals.

In this paper, we extend a result in [Ta04] (cf. Theorem 2.11 in this paper). The result gives a fundamental sequence of the least "successor-type" ptro $\psi^{\Omega_{i}}_{M}(\Omega_{1})$, by which $\psi^{\Omega_{i}}_{M}(\Omega_{1})$ can be characterized as the least fixed point of the function enumerating strongly critical ordinals. We here give a similar sequence $\{\gamma_{n}\}_{n \in \omega}$ of every successor-type ptro $\gamma$. Compared with the previous result in [Ta04], the proof of the property that $\gamma = \lim_{n \in \omega} \gamma_{n}$ needs some special attentions. Therefore, for (a certain type of) a given ordinal $\delta$ less than $\gamma$, we construct a labeled tree informing us the number $n \in \omega$ with $\delta < \gamma_{n}$.

In Section 2, we explain several definitions and results in [Ta04]. In Section 3, we show the extended version of the result above.

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2 Preliminaries

In this paper, $M$ denotes the least Mahlo cardinal, and $\varphi$ the veblin function. For more details, one can refer to [Bu92], [Ra98], [Ra99] or [Ta04].

**Definition 2.1** (Rathjen98,99). For given ordinals $\alpha$ and $\beta$, we define a set $C^M(\alpha, \beta)$ called a Skolem's hull as well as functions $\chi^\alpha$ and $\psi^\alpha_M$ called collapsing functions, as follows:

(M1) $\beta \cup \{0, M\} \subset C^M(\alpha, \beta)$;
(M2) $\gamma = \gamma_1 + \gamma_2 \& \gamma_1, \gamma_2 \in C^M(\alpha, \beta) \Rightarrow \gamma \in C^M(\alpha, \beta)$;
(M3) $\gamma = \varphi \gamma_1 \gamma_2 \& \gamma_1, \gamma_2 \in C^M(\alpha, \beta) \Rightarrow \gamma \in C^M(\alpha, \beta)$;
(M4) $\gamma = \Omega \gamma_1 \& \gamma_1 \in C^M(\alpha, \beta) \Rightarrow \gamma \in C^I(\alpha, \beta)$;
(M5) $\gamma = \chi^\xi(\delta) \& \xi, \delta \in C^M(\alpha, \beta) \& \delta \neq \gamma \in C^M(\xi, \gamma) \& \delta < M \Rightarrow \gamma \in C^M(\alpha, \beta)$

(M6) $\gamma = \psi^\xi_M(\kappa) \& \xi, \kappa \in C^M(\alpha, \beta) \& \xi \neq \alpha \& \xi \in C^M(\xi, \gamma) \Rightarrow \gamma \in C^M(\alpha, \beta)$;

$\chi^\alpha(\delta) \simeq$ the $\delta^{th}$ regular cardinal $\pi < M$ with $C^M(\alpha, \pi) \cap M = \pi$;

$\psi^\alpha_M(\kappa) \simeq \min\{\rho < \kappa : C^M(\alpha, \rho) \cap \kappa = \rho \land \kappa \in C^M(\alpha, \rho)\}$.

**Definition 2.2**

(i) $\gamma =_{nf} \alpha + \beta : \Leftrightarrow \gamma = \alpha + \beta \& \gamma > \alpha \geq \beta \& \beta$ is an additive principal number.
(ii) $\gamma =_{nf} \varphi \alpha \beta : \Leftrightarrow \gamma = \varphi \alpha \beta \& \alpha, \beta < \gamma$.
(iii) $\gamma =_{nf} \Omega \alpha : \Leftrightarrow \gamma = \Omega \alpha \& \alpha < \gamma$.
(iv) $\gamma =_{nf} \psi^\alpha_M(\kappa) : \Leftrightarrow \gamma = \psi^\alpha_M(\kappa) \& \alpha \in C^I(\alpha, \gamma)$.
(v) $\gamma =_{nf} \chi^\alpha(\beta) : \Leftrightarrow \gamma = \chi^\alpha(\beta) \& \beta < \gamma \& \alpha \in C^M(\alpha, \gamma)$.

**Definition 2.3** (Rathjen95,98). We define a set $T(M)$ called an elementary ordinal representation system for KPM and the degree $d(\alpha) < \omega$ of each element $\alpha$ of $T(M)$, as follows:

(i) $0, M \in T(M) \& d(0) = d(M) = 0$;
(ii) ($\gamma =_{nf} \alpha + \beta \& \alpha, \beta \in T(M)$)

$\Rightarrow (\gamma \in T(M) \& d(\gamma) = \max\{d(\alpha), d(\beta)\} + 1)$;
(iii) ($\gamma =_{nf} \varphi \alpha \beta \& \alpha, \beta \in T(M) \& (\gamma < M \text{ or } \alpha = 0)$)

$\Rightarrow (\gamma \in T(M) \& d(\gamma) = \max\{d(\alpha), d(\beta)\} + 1)$;
(iv) ($\gamma =_{nf} \Omega \alpha < M \& \alpha > 0 \& \alpha \in T(M)$)

$\Rightarrow (\gamma \in T(M) \& d(\gamma) = d(\alpha) + 1)$;
(v) ($\gamma =_{nf} \chi^\xi(\alpha) \& \xi, \alpha \in T(M)$)

$\Rightarrow (\gamma \in T(M) \& d(\gamma) = d(\alpha) + 1)$;
(vi) ($\gamma =_{nf} \psi^\alpha_M(\kappa) \& \kappa, \alpha \in T(M)$)

$\Rightarrow (\gamma \in T(M) \& d(\gamma) = \max\{d(\kappa), d(\alpha)\} + 1)$.

**Theorem 2.4** (Rathjen91, Buchholz92). (1) Each element of $T(M)$ has a unique representation with $0, M, +, \varphi, \Omega, \chi, \psi_M$.

(2) $|KPM| \leq \psi^{2M+1}_M(\Omega_1) = T(M) \cap \Omega_1$, where $|KPM|$ denotes the proof theoretic ordinal of KPM.
Definition 2.5 An ordinal \( \gamma \) is called a \textit{proof-theoretically regular ordinal based on} \( \mathcal{T}(M) \) if \( \gamma \) is (expressed by) an element of \( \mathcal{T}(M) \) having the form of \( \psi_{M}^{\kappa}(\Omega_{1}) \) with \( \kappa \in \text{Reg} \), where \( \text{Reg} \) denotes the set of regular cardinals.

Definition 2.6 A \textit{ptro} \( \gamma \) is called a \textit{successor-type ptro} if \( \gamma \) has an element \( \theta \in \mathcal{T}(M) \) satisfying that \( \gamma \) is the least ptro larger than \( \theta \).

Definition 2.7 An ordinal \( \gamma \) is called a \textit{proof-theoretically inaccessible ordinal based on} \( \mathcal{T}(M) \) if \( \gamma \) is an element of \( \text{Reg}(\mathcal{T}(M)) \) as well as the supremum of \( \text{Reg}(\mathcal{T}(M)) \cap \gamma \), where \( \text{Reg}(\mathcal{T}(M)) \) denotes the set of ptros based on \( \mathcal{T}(M) \).

Theorem 2.8 (Takaki 04). All ptros are classified into the following two types:
\( \text{(i)} \) Successor-type ptros, which are of the form \( \psi_{M}^{\Omega_{\alpha+1}}(\Omega_{1}) \) or \( \psi_{M}^{\Omega_{1}}(\Omega_{1}) \);
\( \text{(ii)} \) Proof-theoretically inaccessible ordinals, which are of the form \( \psi_{M}^{\chi^{\alpha}(\beta)}(\Omega_{1}) \) or \( \psi_{M}^{\beta}(\Omega_{1}) \).

Definition 2.9 For each \( n \in \omega \), we define \( \Psi_{n} \) by:
\[
\Psi_{n} = \begin{cases} 
0 & \text{if } n = 0; \\
\psi_{M}^{\Psi_{n-1}^{\beta}(\alpha)}(\Omega_{1}) & \text{if } n > 0.
\end{cases}
\]

Lemma 2.10 For each \( n \in \omega \), \( \Psi_{n} \in \mathcal{T}(M) \) and \( \Psi_{n} < \Psi_{n+1} \).

The purpose of this paper is to extend the following theorem.

Theorem 2.11 (cf. Theorem 4 in [Ta04]). \( \psi_{M}^{\Omega_{1}}(\Omega_{1}) = \lim_{n \in \omega} \Psi_{n} \).

3 Representation of successor-type ptros

Definition 3.1 Let \( \alpha \) and \( \beta \) be elements of \( \mathcal{T}(M) \). Then, for each \( n \in \omega \), we define an ordinal \( \Psi_{n}^{\beta}(\alpha) \), as follows:
\[
\Psi_{n}^{\beta}(\alpha) = \begin{cases} 
\beta & \text{if } n = 0; \\
\psi_{M}^{\Psi_{n-1}^{\beta}(\alpha)}(\Omega_{\alpha+1}) & \text{otherwise}.
\end{cases}
\]

In particular, \( \Psi_{n}(\alpha) := \Psi_{n}^{0}(\alpha) \).

\( \Psi_{n}^{\beta}(\alpha) \) also satisfies properties of \( \Psi_{n} \).

Lemma 3.2 For each \( \alpha, \beta \in \mathcal{T}(M) \), if
\[
\beta < \psi_{M}^{\beta}(\Omega_{\alpha+1}) \quad \text{and} \quad \forall \xi \left( \alpha < \xi \Rightarrow \beta \in \mathcal{C}^{M}(\beta, \xi) \right)
\]
then, for each \( n \in \omega \),
\[
\Psi_{n}^{\beta}(\alpha) \in \mathcal{T}(M) \quad \text{and} \quad \Psi_{n}^{\beta}(\alpha) < \Psi_{n+1}^{\beta}(\alpha). \quad (1)
\]

In particular, for each \( \alpha \in \mathcal{T}(M) \) and \( n < \omega \),
\[
\Psi_{n}(\alpha) \in \mathcal{T}(M) \quad \text{and} \quad \Psi_{n}(\alpha) < \Psi_{n+1}(\alpha).
\]
Proof. This lemma is shown by checking the properties in (1) as well as
\[ \forall \xi \ ( \alpha < \xi \Rightarrow \Psi_n^\beta(\alpha) \in C^M(\Psi_n^\beta(\alpha), \xi) ) , \]
by using induction on \( n \).

Now we give a representation of each successor-type ptro via \( \Psi_n(\alpha) \) and the concept of limit.

**Theorem 3.3** For each \( \alpha \) with \( \psi^\alpha M_1(\Omega_1) \in \mathcal{T}(M) \),
\[ \psi^\alpha M_1(\Omega_1) = \lim_{n \in \omega} \psi^n M_1(\alpha)(\Omega_1) . \quad (2) \]

Proof. Since in [Ta04] we dealt with the case where \( \alpha = 0 \), it suffices to show (2) in the case where \( \alpha > 0 \).

[1] One can show that \( \psi^\alpha M_1(\Omega_1) \geq \lim_{n \in \omega} \psi^n M_1(\alpha)(\Omega_1) \), by the following two claims.

**Claim 1** (cf. Lemmas 9.3 and 11 in [Ta04]). For each \( \alpha \) and \( \beta \), \( \psi^\beta M_1(\alpha+1) \) is defined and \( \Omega_\alpha < \psi^\beta M_1(\alpha+1) < \Omega_\alpha+1 \).

**Claim 2** (cf. Lemma 10 in [Ta04]). For each \( \alpha_1 \), \( \alpha_2 \) and \( \pi \in \text{Reg} \), if \( \psi^\alpha_1 M_1(\pi) \) and \( \psi^\alpha_2 M_1(\pi) \) are defined and if \( \alpha_1 \leq \alpha_2 \), then \( \psi^\alpha_1 M_1(\pi) \leq \psi^\alpha_2 M_1(\pi) \).

[2] In order to show that \( \psi^\alpha M_1(\Omega_1) \leq \lim_{n \in \omega} \psi^n M_1(\alpha)(\Omega_1) \), we show that, for each \( \gamma < \psi^\alpha M_1(\Omega_1) \), there exists an \( n \in \omega \) with \( \gamma \leq \psi^n M_1(\alpha)(\Omega_1) \), by using induction on \( d(\gamma) \).

Since it is easy to check the property above in any case except the case where \( \gamma = \psi^\xi M_1(\pi) \)

[1] More precisely, we should assume that \( \gamma =_{\text{nf}} \psi^\xi M_1(\pi) \). However, we use only the symbol "=" unless we need special attention.
(ii) We define each node of $T_2(\xi)$ and its label, by using recursion on the
distance from the root of $T_2(\xi)$, as follows.

(ii.0) If $s \in T_2(\xi)$ is the root, then $l_s$ is $\xi$.

Let $s$ be a node of $T_2(\xi)$. Then, we define the successors (successor nodes)
of $s$ as well as their labels, according to the following conditions of $l_s$.

(ii.i) If $l_s = 0$, then $s$ is a leaf, that is, $s$ has no successor node.

(ii.ii) If $l_s = \delta + \eta$ or $l_s = \phi \delta \eta$, then $s$ has successors $s_1$ and $s_2$, and

$l_{s_1} := \delta, l_{s_2} := \eta$.

(ii.iii) If $l_s = \Omega_\beta$ and $l_s = \chi^\delta(\eta)$, then $s$ is a leaf.

(ii.iv) Let $l_s = \psi^\delta_M(\tau)$. In this case, $\tau \leq \Omega_{\alpha+1}$ since $l_s \leq \xi$.

(ii.iv.i) If $\tau < \Omega_{\alpha+1}$, then $s$ is a leaf.

(ii.iv.ii) If $\tau = \Omega_{\alpha+1}$, then $s$ has a successor $s_1$ and $l_{s_1} := \delta$.

Claim 3 $T_2(\xi)$ is well-defined to be a finite tree.

(Proof of Claim 3: In order to show that $T_2(\xi)$ is well-defined, we show that,
for each node $s$ of $T_2(\xi)$, $l_s$ satisfies the properties (i.i)～(i.iii) above, by using
induction on the distance from the root to $s$.

If $s$ is the root, it is trivial since $l_s = \xi$.

We let $l_s = \psi^\delta_M(\Omega_{\alpha+1})$ and show that $\delta$ satisfies (i.i)～(i.iii), as follows. By
induction hypothesis, $l_s$ is a subterm of $\xi$, $l_s \leq \xi$ and $l_s \in C^M(\xi, \gamma)$. Then, $\delta$ is
also a subterm of $\xi$. On the other hand, $l_s > \Omega_1 > \gamma$. So, we have $\delta \in C^M(\xi, \gamma)$
and $\delta < \xi$ from Definition 2.1.(M5) and $l_s \in C^M(\xi, \gamma)$.

Any other case is similar to the case above.

Moreover, for each node $s \in T_2(\xi)$ and each successor $s'$ of $s$, it holds that
distance $d(s) > d(s')$. So, $T_2(\xi)$ is finite. \(\Box\)

Definition 3.5 (1) A node $s$ of $T_2(\xi)$ (= $T_2(\xi, \alpha)$) is said to be critical when
$l_s = \psi^\delta_M(\Omega_{\alpha+1})$ for some $\delta$. CN denotes the set of critical nodes (of $T_2(\xi)$).

(2) For each path $p$ of each subtree of $T_2(\xi)$, the number of critical nodes in $p$
is called the weight of $p$. Moreover, for each subtree $T$ of $T_2(\xi)$, the maximum
number of weights of all paths of $T$ is called the weight of $T$, and denoted by
$\text{wt}(T)$. Furthermore, for each node $s$ of $T_2(\xi)$, the weight of the subtree of $T_2(\xi)$
with root $s$ is called the weight of $s$, and denoted by $\text{wt}(s)$.

(3) For each subtree $T$ of $T_2(\xi)$, the maximum length of all paths of $T$ is called
the height of $T$. Moreover, for each node $s$ of $T_2(\xi)$, the height of the subtree of $T_2(\xi)$
with root $s$ is called the depth of $s$, and denoted by $\text{dp}(s)$.

Claim 4 For each node $s$ of $T_2(\xi)$, it holds that $l_s < \Psi_{\text{wt}(s)+1}(\alpha)$.

(Proof of Claim 4: We show the claim by induction on the depth of $s$.

(i) If $s$ is a leaf, then $l_s \leq \Omega_\alpha$. So, since $\Omega_\alpha < \Psi_n(\alpha)$ for each $n > 0$, we have
$l_s < \Psi_1(\alpha)$.\)
(ii) Assume that $s$ is not any leaf. Then, $l_s = \text{nf} \delta + \eta$, $l_s = \text{nf} \varphi \delta \eta$, or $l_s = \text{nf} \psi_M^\delta(\Omega_{\alpha+1})$.

Let $l_s = \text{nf} \psi_M^\delta(\Omega_{\alpha+1})$. Then, $l_s \in \text{CN}$ and $s$ has one successor $s_1$ with $l_{s_1} = \delta$. Since $\mathrm{wt}(s_1) = \mathrm{wt}(s) - 1$ and $d_p(s_1) < d_p(s)$, the induction hypothesis implies that $l_{s_1} < \Psi_{\mathrm{wt}(s)}(\alpha)$. On the other hand, since $l_s \in T(M)$ and $\Psi_{\mathrm{wt}(s)+1}(\alpha) \in T(M)$, we have $l_s < \Psi_{\mathrm{wt}(s)+1}(\alpha)$ (cf. Lemma 16 in [Ta04]).

Any other case is similar to or easier than the case above. \hfill \Box

By Claim 4, we have $\xi < \Psi_{\mathrm{wt}(T_2(\xi))+1}(\alpha)$, and hence, by Claim 2,

$$\gamma \leq \psi_M^{\Psi_{\mathrm{wt}(T_2(\xi))+1}(\alpha)}(\Omega_1).$$

So, the proof of Theorem 3.3 is completed. \hfill \Box

We can also expect that each $\psi_M^{\Psi_n(\alpha)}(\Omega_1)$ has itself as its regular expression, that is, $\psi_M^{\Psi_n(\alpha)}(\Omega_1) \in T(M)$. Unfortunately, we have not yet completed the proof of the property. However, it is not hard to show this property for each $\alpha$ less than a certain ordinal. For example, one can easily show the following proposition.

**Proposition 3.6** For each $\alpha \in T(M)$ and $n \in \omega$, if $\alpha \in C^M(\Psi_n(\alpha), \psi_M^{\Psi_n(\alpha)}(\Omega_1))$, then

$$\psi_M^{\Psi_n(\alpha)}(\Omega_1) \in T(M) \text{ and } \psi_M^{\Psi_n(\alpha)}(\Omega_1) < \psi_M^{\Psi_{n+1}(\alpha)}(\Omega_1).$$

By Theorem 3.3 and Proposition 3.6, each successor-type ptro $\psi_M^{\Omega_{\alpha+1}}(\Omega_1)$ has a fundamental sequence $\{\psi_M^{\Psi_n(\alpha)}(\Omega_1)\}_{n \in \omega}$ if $\alpha \in C^M(\Psi_n(\alpha), \psi_M^{\Psi_n(\alpha)}(\Omega_1))$.

**Reference**


[Ta04] O. Takaki, Primitive recursive analogues of regular cardinals based on ordinal representation systems for KPi and KPM, to appear in AML.