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Zeta function of a linear code
and its Riemann hypothesis property

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1 Introduction

Duursma defined zeta function of code first in 1999. After that, the definition of it was expanded even general linear code. Furthermore, a Riemann hypothesis analogue for self-dual linear code was formulated. In this paper, we introduce Duursma's theory.

2 Preliminaries

Let $C$ be a linear code of length $n$ and minimum distance $d$ over the finite field of $q$ elements. Let $A_i$ be the number of words of weight $i$ in $C$. The weight distribution may be represented by a polynomial

$$W_C(x, y) = \sum_{i}^{n} A_i x^{n-i} y^i$$

called the weight enumerator.

**Definition 2.1** The zeta polynomial $P(T)$ of $C$ is the unique polynomial of degree at most $n - d$ such that generating function

$$\frac{P(T)}{(1 - T)(1 - qT)} (y(1 - T) + x T)^n$$

has expansion

$$\ldots + \frac{W_C(x, y) - x^n}{q-1} T^{n-d} + \ldots.$$ 

The quotient $Z(T) = P(T)/((1 - T)(1 - qT))$ is called the zeta function of the linear code.
Definition 2.2 Let $C$ be a linear code over the field $F_q$ of $q$ elements has as main parameters its length $n$, dimension $k$, and minimum distance $d$. Then dual code of $C$ is defined by

$$C^\perp = \{ u \in F_q \mid u \cdot v = 0 \ \forall v \in C \},$$

where for all $u = (u_1, \ldots, u_n)$ and $v = (v_1, \ldots, v_n)$ in $F_q$, inner product $u \cdot v$ is defined by

$$u \cdot v = u_1v_1 + \cdots + u_nv_n.$$ 

Dimension and minimum distance of $C^\perp$ is denoted by $k^\perp$ and $d^\perp$ respectively.

Definition 2.3 If $C$ is equal to its dual code $C^\perp$, then the code is called self-dual code.

Theorem 2.1 For zeta polynomial $P(T)$, the following holds.

(i) $\deg P(T) = n + 2 - d - d^\perp$

(ii) Let zeta polynomial and zeta function of $C^\perp$ be $P^\perp(T)$ and $Z^\perp(T)$ respectively. Then

$$P^\perp(T) = P\left(\frac{1}{qT}\right)q^gT^{g+g^\perp},$$

$$Z^\perp(T) = Z\left(\frac{1}{qT}\right)q^{g-1}T^{g+g^\perp-2},$$

where $g = n + 1 - k - d$, $g^\perp = n + 1 - k^\perp - d^\perp$.

In particular, if $C$ is self-dual code, since $P(T) = P^\perp(T)$, the following hold.

(i)' $\deg P(T) = 2g$

(ii)'

$$P(T) = P\left(\frac{1}{qT}\right)q^gT^{2g}$$

$$Z(T) = Z\left(\frac{1}{qT}\right)q^{g-1}T^{2g-2}$$
Proof. [2, p59].

By the way, like these equations, there are some equations for weight enumerator.

**Theorem 2.2** For weight enumerator of $C$, the following hold.

(i) $\overline{W}_C(x, y) := W_C(x + y, y) \Rightarrow \overline{W}_{C^\perp}(x, y) = \frac{1}{|C|} \overline{W}_C(qy, x)$

(ii) $\overline{W}_C(z) := \overline{W}_C(1, z) \Rightarrow \overline{W}_{C^\perp}(z) = \frac{(qz)^n}{|C|} \overline{W}_C\left(\frac{1}{qz}\right)$

(iii) $W_{C}^{\dim}(x, y) := \sum_{R \subseteq N} \dim C(R)x^{n-|R|}y^{|R|}$

$\Rightarrow W_{C^\perp}^{\dim}(x, y) = (x + y)^{n-1}\{(n - k)y - kx\} + W_{C}^{\dim}(y, x)$

### 3 A Riemann hypothesis analogue for self-dual codes

**Definition 3.1** [3, p119 Def4.1] Let $C$ be self-dual code, $P(T)$ be its zeta polynomial. $C$ is called that $C$ has the Riemann hypothesis property, when for all zeros $\alpha$ of $P(T)$, $|\alpha| = \frac{1}{\sqrt{q}}$.

**Definition 3.2** Let $C$ be a self-dual code. $C$ is called extremal when equality holds in the following upper bounds.

- (Type I) $d \leq 2[n/8] + 2$
- (Type II) $d \leq 4[n/24] + 4$
- (Type III) $d \leq 3[n/12] + 3$
- (Type IV) $d \leq 2[n/6] + 2$
Four type is a classification of a non-trivial divisible self-dual code defined over $F_q$. A code is said to be divisible when all weights are divisible by an integer $c$ greater than one. Type I, II, III and IV means $(q, c) = (2, 2), (2, 4), (3, 3)$ and $(4, 2)$ respectively.

**Problem** [3, p119 open problem4.2] Do all extremal weight enumerators have the Riemann hypothesis property?

**Example 3.1** [8,4,4] extended hamming code $C_8$ is a self-dual binary extremal doubly even code. It’s weight enumerators is

$$W_{C_8}(x, y) = x^8 + 14x^4y^4 + y^8.$$ 

Hence, it’s zeta polynomial is

$$P(T) = \frac{1}{5}(1 + 2T + 2T^2).$$

Since $\alpha = \frac{1 \pm i}{2}$, so $|\alpha| = \frac{1}{\sqrt{2}} = \frac{1}{\sqrt{q}}$. So $C_8$ has the Riemann hypothesis property.

**Example 3.2** [72,36,16] code

If such a code exists, then the zeros all have same absolute value $\frac{1}{\sqrt{2}}$.

**Example 3.3** $C_8 \oplus C_8 \oplus C_8$ is the set of words $(a|b|c)$ where $a$, $b$, $c$ are arbitrary words of $C_8$. This code is type II and not extremal. This code is not satisfy the Riemann hypothesis property.

**Theorem 3.1** Extremal self-dual code of type IV has the Riemann hypothesis property.

In [4], Duursma obtained this theorem. But, it seems that it can’t be proved as for three other types yet. Like this, the necessary and sufficient condition for zeta function of code to satisfy Riemann hypothesis property doesn’t get clear yet.
References


