A Scheduling Scheme for Continuous Media Data Broadcasting With a Single Channel

Tomoki Yoshihisa, Masahiko Tsukamoto, and Shojiro Nishio

Abstract—Recently, various schemes for broadcasting continuous media data such as audio and video have been studied. Some of them have focused on reducing the waiting time of clients under the condition that clients can play data without interruption from beginning to end. These schemes usually employ multiple channels to broadcast continuous media data. However, clients for most broadcast systems such as wireless LAN, DVB and ISDB-T cannot receive data from multiple channels concurrently. In this paper, we propose and evaluate a scheduling scheme to reduce the waiting time of clients with a single channel. Preliminary results suggest that this scheme outperforms existing schemes.

Index Terms—Broadcasting, continuous media data, multimedia communication, streaming, video on demand.

I. INTRODUCTION

RECENTLY, broadcasting continuous media data such as audio and video has become increasingly popular. In a typical broadcast system, a server repetitively broadcasts data. Although the server can simultaneously deliver data to many clients, each client has to wait until its desired data is broadcast. On the other hand, in continuous media data broadcasting, it is important for clients to be able to play data without interruption from beginning to end. Accordingly, various schemes are proposed to reduce clients’ waiting time under the condition that clients can play data continuously from start to finish ([1]–[3], [5]–[9], [16]). These schemes reduce the waiting time by dividing data into several segments and broadcasting segments via multiple channels. However, with multiple channels, clients are forced to receive data from multiple channels concurrently, and the server has to control all these multiple channels; hence, their mechanisms are more complicated than broadcasting data via a single channel. In practical systems, unfortunately, many clients cannot receive data from multiple channels for such hardware reasons as complexity or mounting space. For example, most of the commercial DVB (Digital Video Broadcasting) systems cannot receive data from multiple channels concurrently. Moreover, since almost all set-top-boxes, i.e., clients of DVB systems, have only one tuner, they cannot receive data from multiple channels concurrently. Accordingly, it is more realistic to assume that clients can receive data only from a single channel.

In this paper, we propose a scheduling scheme to reduce the waiting time of clients under the condition that clients can play data without interruption from beginning to end. Our scheme divides data into several segments and broadcasts them via a single channel. By broadcasting precedent segments frequently, the waiting time is reduced. In our scheme, the number of segments and the broadcast schedule affect the waiting time. We present a schedule that provides the minimum average waiting time in schedules where data is divided into two segments of equal size. In the case where it is necessary to divide data into more than two segments, we propose a heuristic scheduling method.

The remainder of this paper is organized as follows. Section II explains our basic idea. We show how to reduce the waiting time by broadcasting precedent segments frequently with a single channel. Our assumed system environment and proposed scheme are explained in Section III. In Section IV, we analyze schedules where data is divided into two segments of equal size and present a schedule that provides the minimum average waiting time. In the case of dividing data into more than two segments, we propose a heuristic scheduling method in Section V. Our scheme is discussed in Section VI. In Section VII, we make remarks on related work. Finally, we conclude the paper in Section VIII.

II. BASIC IDEA

In this section, first we calculate the average waiting time in the case of broadcasting data repetitively without division. Next, we explain how broadcasting precedent segments frequently with a single channel reduces the waiting time.

Suppose that it takes 180 seconds to broadcast a continuous media data for 30 minutes. Note that this case is suitable for broadcasting an MPEG2 (5 Mbps)-encoded video for 30 minutes on 50 Mbps satellite broadcast systems. In the case of broadcasting the data repetitively without division as shown in Fig. 1 (a simple repetition scheme), a client which demands the data should wait for 90 seconds on average to start playing the data (max. 180 seconds, min. 0 seconds). Note that the data is streaming media data, i.e., the client can start playing the data as soon as it starts receiving it.

Next, we explain our basic idea. Suppose the case of dividing the data into two segments of equal size. The playing time of each segment is 15 minutes and it takes 90 seconds to broadcast it. Let \( S_1 \) be the first segment and \( S_2 \) the last segment. The server
broadcasts these segments in the order of $S_1S_2S_3$ as shown in Fig. 2, and in this case, the broadcast cycle is 270 seconds. Let $T_i (i = 1, 2, \cdots)$ be the time slot. The broadcast cycle starts from $T_1$. That is, $S_1$ is broadcast from $T_1$ and $T_2$, and $S_2$ is broadcast from $T_3$. The average waiting time of clients is calculated as follows.

- In the case of demanding the data in $T_1$:

  Since $S_1$ is broadcast from $T_2$, the client can start playing the data from the beginning of $T_2$. The time to finish playing $S_1$ is after $T_3$, because the playing time of $S_1$ is 15 minutes. Hence, the client can store $S_2$ broadcast from $T_3$ while it plays $S_1$. After the client finishes playing $S_1$, it can play the data without interruption by playing stored $S_2$ immediately. In this case, the average waiting time of clients is $90/2 = 45$ seconds.

- In the case of demanding the data in $T_2$:

  Since $S_2$ is broadcast from $T_3$, the client can store $S_2$ before it starts playing the data. In this case, the average waiting time is 135 seconds.

- In the case of demanding the data in $T_3$:

  In this case, the average waiting time is 45 seconds.

As a result, the average waiting time becomes $(45 + 135 + 45)/3 = 75$ seconds. This is 17% shorter than that of the simple repetition scheme.

Moreover, when the server broadcasts segments in the order of $S_1S_1S_1S_1S_1S_1S_1S_1S_1S_2$, it is similarly shown that the average waiting time falls to 53 seconds, 42% shorter than that of the simple repetition method. This is the minimum average waiting time in this example.

In this way, by broadcasting precedent segments frequently with a single channel, clients’ waiting time can be reduced.

### III. OUR APPROACH

Let $D$ be the playing time of a continuous media data and $D'$ the time required to broadcast the data. We define the playback ratio $a$ by the following equation:

$$a = \frac{D}{D'}.$$  \hspace{1cm} (1)

For example, an MPEG2 (5 Mbps)-encoded video lasting 30 minutes is 1,125 Mbytes. If the data is broadcast via a 50 Mbps satellite broadcast system, it takes 180 seconds to broadcast it. In this case, the playback ratio is $1800/180 = 10$. Moreover, MPEG1-encoded video is usually assumed to be 1 Mbps. If the data is broadcast via a 20 Mbps DVB system, the playback ratio is 20 [10]. Not only video, but also audio can be applied to this discussion. An MP3 (Moving Picture Experts Group Audio Layer 3)-encoded audio is usually 128 Kbps. If the data is broadcast via a 1 Mbps wireless LAN system, the playback ratio is 7.8. Note that the playback ratio is derived while encoding the data.

In this paper, we assume that $a$ is constant and known. However, if $a$ varies, our proposed scheme can be applied by using the maximum value of the playback ratio as $a$.

### A. Assumed System Environment

- The client can receive data only from a single channel.
- The client can play data without interruption from beginning to end.
- The client can start playing data as soon as it starts receiving it.
- The client has its own buffer. The client can receive and store new data while playing data. The buffer has enough capacity to store the data.
- The client starts receiving broadcast data after it demands playing the data.
- Since a segment includes additional information such as the segment number in front of it, the client cannot receive segments midstream.
- The server does not update data.
- $a$ is larger than 1.

We assume guarantee-type networks that guarantee steady bandwidth and reliable data receiving. The Internet is assumed to be a guarantee-type network by applying a bandwidth reservation protocol such as RSVP (Resource ReSerVation Protocol).

In the case where the server broadcasts some contents, it broadcasts them one by one. Since clients do not receive a data broadcast prior to demanding the data, they do not always have to observe broadcast data.

### B. Proposed Scheme

We propose a scheduling scheme called the division-based broadcasting scheme. In this scheme, a continuous media data is divided into $N$ segments $S_i (i = 1, \cdots, N)$ of equal size. The necessary time to broadcast a segment is $D' / N$ (Fig. 3). The server broadcasts $C$ segments in a broadcast cycle.

The problem is, how we produce a broadcast schedule that effectively reduces the waiting time? For example, suppose the
case where it takes 180 seconds to broadcast a video for 30 minutes. Table I shows broadcast schedules that include four or five ($c = 4, 5$) schedule segments in one broadcast cycle when the data is divided into two segments. Since $S_1, S_2$ are regarded as $c = 2$, this broadcast schedule is not shown in the table. The minimum average waiting time in Table I is 63 minutes for the broadcast schedule $S_1, S_2, S_1, S_2$. Table II shows broadcast schedules that include four schedule segments in a broadcast cycle when the data is divided into three segments. In this case, the minimum average waiting time is 60 seconds for the broadcast schedule $S_1, S_2, S_1, S_2$. Thus, the average waiting time is calculated from given $a, N$, and $c$. However, since these values are given voluntarily and the number of broadcast schedules is infinity, finding a broadcast schedule that provides the minimum average waiting time in all schedules is not realistic. Hence, in Section IV, we analyze schedules in the case of dividing data into two segments of equal size and present a schedule that provides the minimum average waiting time. In the case of dividing data into more than two segments, since the large number of broadcast schedules makes analysis difficult, we propose a heuristic scheduling method in Section V.

### IV. Dividing Data into Two Segments

In this section, we analyze schedules in the case of dividing data into two segments and present a schedule that provides the minimum average waiting time.

#### A. Classification of the Average Waiting Time

To find a broadcast schedule that provides the minimum average waiting time, we classify equations. Let $d$ be the necessary time to broadcast a segment. Since we assume that the additional information data size of segments is small, $d = D'/2$. We call the time in which a segment is broadcast a time slot $T_1 (i = 1, 2, \cdots)$. The start time of $T_1$ is denoted by $t_1$. Let $T_1$ be the time slot in which the client demands the data.

The timing for a client to receive $S_1$ or $S_2$ after the demand is essential for calculating the average waiting time. Therefore, we term the time slot in which the first $S_1$ is broadcast after $T_1$ as $T_m$, and the time slot in which the first $S_2$ is broadcast as $T_n$. By comparing $m$ with $n$, the situation can be divided into two cases, as shown in Figs. 4 and 5. Since clients cannot receive the segment broadcast from $T_1$, the segment broadcast from $T_1$ is not represented in Fig. 4. $m$ and $n$ must be different numbers. When the value of either $m$ or $n$ is 2, the other value must be more than 2.

Suppose that the client starts playing the data as soon as it receives $S_1$. Let $\tau$ be the time at which the client finishes playing $S_1$. $\tau$ is expressed by $t_m + ad$. In the case of $n \leq a + m$, since $t_n \leq \tau$, the client can receive $S_2$ while it plays $S_1$. In the case of $n > a + m$, since $t_n > \tau$, the client cannot receive $S_2$ while it plays $S_1$. Below, we show the average waiting time of each case.

- **Where $n \leq a + m$:**

  After the client finishes playing $S_1$, it can play $S_2$ continuously. The average waiting time of this case $W^A(m)$ is given by the following equation:

  \[
  W^A(m) = \frac{d}{2} + (m-2)d. \quad (2)
  \]

#### Table I

<table>
<thead>
<tr>
<th>Broadcast Schedule</th>
<th>Average Waiting Time (seconds)</th>
<th>Broadcast Schedule</th>
<th>Average Waiting Time (seconds)</th>
</tr>
</thead>
<tbody>
<tr>
<td>$S_1, S_1, S_2, S_2$</td>
<td>67.5</td>
<td>$S_1, S_1, S_1, S_2$</td>
<td>63</td>
</tr>
<tr>
<td>$S_1, S_2, S_2, S_2$</td>
<td>112.5</td>
<td>$S_1, S_1, S_2, S_2$</td>
<td>99</td>
</tr>
</tbody>
</table>

#### Table II

<table>
<thead>
<tr>
<th>Broadcast Schedule</th>
<th>the average waiting time (seconds)</th>
</tr>
</thead>
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<td>$S_1, S_1, S_2, S_2$</td>
<td>75</td>
</tr>
<tr>
<td>$S_1, S_2, S_2, S_2$</td>
<td>75</td>
</tr>
<tr>
<td>$S_1, S_2, S_2, S_2$</td>
<td>60</td>
</tr>
<tr>
<td>$S_1, S_2, S_2, S_2$</td>
<td>120</td>
</tr>
<tr>
<td>$S_1, S_2, S_2, S_2$</td>
<td>120</td>
</tr>
<tr>
<td>$S_1, S_2, S_2, S_2$</td>
<td>120</td>
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<tr>
<td>$S_1, S_2, S_2, S_2$</td>
<td>120</td>
</tr>
<tr>
<td>$S_1, S_2, S_2, S_2$</td>
<td>120</td>
</tr>
</tbody>
</table>

Fig. 3. A division of a continuous media data.

Fig. 4. An example of the value of $m$ and the value of $n \leq a + m$.

Fig. 5. An example of the value of $m$ and the value of $n > a + m$. 

![Broadcast Schedules](image.png)

![Table I](table.png)

![Table II](table.png)
The meanings of terms are shown in Fig. 4. Since \( W^I(M) \) are the average waiting time, we showed the case which a client demands the data at the center of a time slot in the figure.

- Where \( n > m + \alpha \)

In this case, if the client starts playing the data from \( t_m \), it cannot play \( S_2 \) continuously after \( S_1 \). Hence, the client has to wait before it starts playing \( S_1 \) to play the data continuously. To calculate the average waiting time, we suppose that the client starts playing the data from time \( t_w \). Let \( \tau_w \) be the time at which the client finishes playing \( S_1 \). That is, \( \tau_w = t_w + ad \). Therefore, \( S_2 \) is broadcast until or from \( \tau_w \), and the client can play \( S_2 \) after \( S_1 \). Hence, \( t_w \leq \tau_w \) must be established. That is, \( w \geq n - a \). In this case, when \( w = n - a \), the average waiting time is the minimum. The average waiting time of this case \( W^{II}(n) \) is given by the following equation:

\[
W^{II}(n) = \frac{d}{2} + (n-a-2)d. \tag{3}
\]

The meanings of terms are shown in Fig. 5.

Here, we consider the order of the average waiting times. Since \( \alpha \leq a < \alpha + 1 \) (\( \alpha \) denotes the integral part of \( a \), i.e., \( \alpha = \lfloor a \rfloor \)), the following inequality is derived.

\[
W^I(2) < W^{II}(\alpha + 3) \leq W^{II}(3) < W^{II}(\alpha + 4) \leq W^{II}(4) < W^{II}(\alpha + 5) \leq W^I(5) < \cdots \tag{4}
\]

By using the order of \( W^I(M) \) and \( W^{II}(N) \) in (4), we will find the optimal schedule in Section IV-C.

### B. Segment Sets

In the case of dividing the data into two segments, the broadcast schedule consists of \( S_1 \) and \( S_2 \). Hence, we divide the broadcast schedule into segment sets at the point between \( S_2 \) and \( S_1 \), i.e., let \( U_i = S_1 \cdots S_iS_2 \cdots S_2 \) \((i = 1, \cdots, I)\) be segment sets. The broadcast schedule is represented by \( U_1 \cdots U_I \). \( T_{pi} \) indicates the time slot in which the first \( S_1 \) of \( U_i \) is broadcast, \( T_{qi} \) indicates the time slot in which the first \( S_2 \) of \( U_i \) is broadcast, and \( T_{ri} \) indicates the time slot in which the last \( S_2 \) of \( U_i \) is broadcast. Then, \( \tau_i + 1 = p_{i+1} \) holds for all \( i = 1, \cdots, I-1 \). Examples of segment sets are shown in Table III. As the table demonstrates, all broadcast schedules can be divided into segment sets. Since the first segment of a segment set is \( S_1 \), and the last segment of the segment set is \( S_2 \), each combination dividing a broadcast schedule into segment sets is unique.

As for segment sets, the next theorem is established.

**Theorem 1**: A segment set \( U_i \) defines the waiting time of clients that demand the data in \( T_{pi-1} \cdots, T_{ri-1} \).

For example, a client that demands the data in \( T_{pi-1} \) can receive \( S_1 \) from \( T_{pi} \), and receive \( S_2 \) from \( T_{qi} \). Hence, the client’s waiting time is decided by the value of \( p_i \) and \( q_i \) of \( U_i \). Moreover, a client that demands the data in \( T_{ri-1} \) can receive \( S_2 \) from \( T_{ri} \) while it waits for \( S_1 \). In this case, since the first segment of the next segment set is \( S_1 \), the client can receive \( S_1 \) from \( T_{pi+1} \). From the reasons described above, segment sets included in a broadcast schedule are independent of the waiting time. Hence, the average waiting time is calculated by the average of the average waiting time of each segment set.

That is, we can determine the waiting time of a request by looking ahead and seeing which \( U_i \) is present. This is simple and can be implemented for actual systems. For example, by calculating the waiting time for each time slot beforehand, the server can get the waiting time needed for clients to play the data without interruption. The server adds the information about the necessary waiting time in front of the segment included in each time slot and broadcasts them. Since clients can get the necessary waiting time as soon as they start receiving the segment, they do not have to look ahead and see which \( U_i \) is there. This can be implemented for actual systems and is realistic.

Fig. 6 shows how the average waiting time is calculated. In this case, the broadcast schedule is divided into 3 segment sets. Since the second \( S_2 \) is broadcast after 3 time slots from the first time slot, the value of \( m \) of the first time slot is 3. Also, since the first \( S_2 \) is broadcast after 2 time slots from the first time slot, the value of \( n \) of the first time slot is 2. When \( a = 2 \), since \( n \leq a + m \) is established, the type is II (Section IV-A). Therefore, the average waiting time of a client which demands the data in the first time slot is \( 3d/2 \). Moreover, that of a client which demands the data in the last time slot is \( d/2 \). As a result, in the case of broadcasting \( U_1 \) as a broadcast schedule, the average waiting time is \( d \). In the case of the second time slot, since \( m = 2 \) and \( n = 5 \), \( n > a + m \) is established. Accordingly, the type is II and the average waiting time is \( 3d/2 \). In the same way as \( U_1 \), the average waiting time of \( U_2 \) is \( d \). The average waiting time of \( U_3 \) is \( 3d/2 \). To calculate the average waiting time of the schedule, we have to consider the number of segments included in each segment set. Since the number of segments included in \( U_1, U_2, \)

<table>
<thead>
<tr>
<th>The broadcast schedule</th>
<th>Segment sets</th>
</tr>
</thead>
<tbody>
<tr>
<td>(S_1S_2)</td>
<td>(U_1 = {S_1S_2})</td>
</tr>
<tr>
<td>(S_1S_2)</td>
<td>(U_2 = {S_1})</td>
</tr>
<tr>
<td>(S_1S_1S_2)</td>
<td>(U_3 = {S_1S_2})</td>
</tr>
<tr>
<td>(S_1S_1S_2S_2)</td>
<td>(U_4 = {S_1S_2S_2})</td>
</tr>
<tr>
<td>(S_1S_2S_2S_2)</td>
<td>(U_5 = {S_1})</td>
</tr>
<tr>
<td>(S_1S_1S_2S_2S_2)</td>
<td>(U_6 = {S_1S_2S_2})</td>
</tr>
<tr>
<td>(S_1S_2S_2S_2S_2)</td>
<td>(U_7 = {S_1})</td>
</tr>
<tr>
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<td>(U_8 = {S_1S_2S_2S_2})</td>
</tr>
<tr>
<td>(S_1S_2S_2S_2S_2S_2)</td>
<td>(U_9 = {S_1})</td>
</tr>
<tr>
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<td>(U_{10} = {S_1S_2S_2S_2})</td>
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<td>(U_{11} = {S_1})</td>
</tr>
<tr>
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<td>(U_{12} = {S_1S_2S_2S_2})</td>
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<tr>
<td>(S_1S_2S_2S_2S_2S_2)</td>
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<table>
<thead>
<tr>
<th>Type</th>
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<th>(U_2)</th>
<th>(U_3)</th>
<th>(U_4)</th>
<th>(U_5)</th>
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<th>(U_7)</th>
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<th>(U_{10})</th>
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<th>(U_{12})</th>
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<tr>
<td>Time</td>
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<td>II</td>
<td>III</td>
<td>I</td>
<td>II</td>
<td>III</td>
<td>I</td>
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<td>III</td>
<td>I</td>
<td>II</td>
<td>III</td>
<td>I</td>
</tr>
<tr>
<td>Waiting time</td>
<td>determined by (U_1)</td>
<td>determined by (U_2)</td>
<td>determined by (U_3)</td>
<td>Waiting time</td>
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</tbody>
</table>

![Fig. 6](image_url) 
An example of segment sets.
The average waiting time for a segment set

<table>
<thead>
<tr>
<th>$U$</th>
<th>$W^1(2)$</th>
<th>$W^1(\alpha + 3)$</th>
<th>$W^1(3)$</th>
<th>$W^1(\alpha + 4)$</th>
<th>$W^1(4)$</th>
<th>\ldots</th>
</tr>
</thead>
<tbody>
<tr>
<td>$S_1 S_2$</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td></td>
</tr>
<tr>
<td>$S_1 S_2$</td>
<td>1</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>$S_1 S_2$</td>
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<td>0</td>
<td>1</td>
<td>0</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>$S_1 S_1 S_2$</td>
<td>3</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td></td>
</tr>
</tbody>
</table>

$U_3$ is 2, 4, 3, respectively, the average waiting time becomes $(2 \times d + 4 \times d + 3 \times 3d/2)/9 = 7d/6$.

Let $U_{\min}$ be the segment set that gives the minimum average waiting time of all segment sets. Here, the following theorem is established.

**Theorem 2**: A broadcast schedule that repeats $U_{\min}$ provides the minimum average waiting time.

To prove the theorem, suppose that a broadcast schedule consists of the segment sets $U_1 \cdots U_i$. When all of $U_i$ are not the same as $U_{\min}$, the average waiting time of $U_i$ is longer than $U_{\min}$. Therefore, if one of these segment sets is changed to $U_{\min}$, the average waiting time decreases. By doing this over the same arguments, we can see that a broadcast schedule that provides the minimum average waiting time is the schedule that consists of repetitions of $U_{\min}$. Hence, by finding $U_{\min}$, we can find a broadcast schedule that provides the minimum average waiting time in the case where the data is divided into two segments of equal size.

C. Detection of $t_{\min}$

To detect $U_{\min}$, we can easily see the number of $W^1(n)$ and $W^1(\alpha)$ in a segment set $U$. The result is shown in Table IV. Since $U$ is a segment set, the form of $U$ is $S_1 \cdots S_1 S_2 \cdots S_2$. In Table IV, equations of average waiting times are lined up in the first line. According to (4), $W^1(2)$ is the shortest of all. A equation gives longer average waiting time along with right. The number of equation which gives the average waiting time of each time slot is counted in the right side of schedules shown in the table. Suppose that $U$ is $S_1 S_2$ as shown in the first line of Table IV. This schedule is same as the simple method. In this case, the number of $W^1(2)$ is one and $W^1(3)$ is also one. First, by adding $S_2$ to $U$, $U$ becomes the second line, and the number of $W^1(4)$ becomes one. Since $W^1(3) < W^1(4)$, the average waiting time of $U$ increases. By increasing $S_2$, the number of $W^1(5)$ becomes one. Hence, the more the number of $S_2$ increases, the longer the waiting time becomes. Accordingly, when $U$ includes only one $S_2$, the average waiting time is shorter. Next, by adding $S_1$ to $U$, $U$ becomes the third line. The number of $W^1(2)$ becomes two. Since the average waiting time by $W^1(2)$ is the minimum, the average waiting time decreases. As $S_1$ increases, the average waiting time is reduced while the number of $S_1$ is less than $\alpha$. In the case where the number of $S_1$ equals $\alpha$, $U$ becomes $(a)$ in Table IV. By adding $S_1$ to $U$, the number of $W^1(\alpha + 3)$ becomes one, and $U$ becomes $(b)$. If the average waiting time of $(a)$ is less than $W^1(\alpha + 3)$, the average waiting time of $(a)$ is less than the average waiting time of $(b)$. Otherwise, the average waiting time of $(a)$ is greater than the average waiting time of $(b)$. By adding $S_1$ to $(b)$, the number of $W^1(\alpha + 4)$ becomes one. Since $W^1(\alpha + 3) < W^1(\alpha + 4)$, the average waiting time increases. Hence, a schedule that provides the minimum waiting time (i.e., $U_{\min}$) can be $(a)$ or $(b)$ in Table IV. In the case of $(a)$, the average waiting time $W_a$ is given by the following equation:

$$W_a = \frac{\alpha + 3}{2\alpha + 2}d.$$  

(5)

In the case of $(b)$, the average waiting time $W_b$ is given by the following equation:

$$W_b = \frac{3\alpha + 6 - 2\alpha}{2\alpha + 4}d.$$  

(6)

From (5) and (6), we compare the average waiting time of $(a)$ with that of $(b)$.

$$W_a - W_b = \frac{(\alpha + 3)d - (3\alpha + 6 - 2\alpha)}{2\alpha + 4}d = \frac{(\alpha + 1)(\alpha - \alpha - \alpha)}{(\alpha + 1)(\alpha + 2)}d.$$  

(7)

Here, let $A$ be the numerator of (7).

$$A = \left(\text{the minimum integer larger than } \alpha \right) \times \left(\text{the decimal part of } \alpha \right) - \left(\text{the integer part of } \alpha \right)$$

A broadcast schedule that provides the minimum average waiting time is given as follows.

- In the case of $A < 0$:

$$\frac{S_1 \cdots S_1 S_2}{\alpha}$$  

(8)

and the average waiting time is obtained by (5).

- In the case of $A > 0$:

$$\frac{S_1 S_1 \cdots S_1 S_2}{\alpha + 1}$$  

(9)

and the average waiting time is obtained by (6).

- In the case of $A = 0$:

the above two schedules have the same average waiting time.
A broadcast schedule that provides the minimum average waiting time in the case of dividing data into two segments in equal size is found above.

V. DIVIDING DATA INTO MORE THAN TWO SEGMENTS

Dividing data into more than two segments makes it difficult to analyze schedules. A solution is to find the optimal broadcast schedule that provides the minimum average waiting time in all broadcast schedules through computer programming. Hence, we produced a computer program which finds the optimal schedule that includes \( c \) segments in a broadcast cycle. This program enables us to find the optimal broadcast schedule when \( a = 10 \) and \( c = 7, \ldots, 11 \). The computation time is shown in Table V (CPU: Pentium III 1.2 GHz, Memory: 512 MBytes). Since the computation time is \( O(N^c) \), the computation time for finding the optimal schedule of \( c = 24 \) is estimated at 3,000 years. Clearly then, finding the optimal broadcast schedule with this program is not realistic. Moreover, since we can choose arbitrary \( N \) and \( c \), the combinations of segments are infinity. Accordingly, we cannot find the optimal solution for all \( N \) and \( c \). For above reasons, finding the optimal solution is impossible. Therefore, in this paper, we propose a heuristic scheduling method.

A. Alternative Broadcasting Method

Our proposed method, the alternative broadcasting method for single data, is based on the division-based broadcasting scheme. As a result of finding the schedule which gives the minimum average waiting time for values that require relative low computation power, we got the tendency that the waiting time is effectively reduced when the broadcast schedule includes precedent segments frequently. Hence, in the alternative scheme, the server divides a continuous media data into several segments of equal size and schedules segments in order to broadcast precedent segments frequently. segments of equal size, and by scheduling segments in order to broadcast precedent segments frequently.

In the alternative broadcasting method, the server broadcasts \( S_1 \) in the following time slot. \( T_1 \) is a time slot from which a broadcast cycle starts:

- In the case of \( n = 1 \): \( S_1 \) is broadcast in \( T_2j-2 \) \((j = 1, \ldots, N - 1)\).
- In the case of \( n \neq 1 \): \( S_n \) is broadcast in \( T_2n-2 \).

Examples of schedules produced by the alternative broadcasting method are shown in Fig. 7. As the figure demonstrates, the first segment and the other segments are broadcast alternatively in the alternative broadcasting method. The segment number next to \( S_1 \) becomes larger, and the number of segments broadcast in one broadcast cycle is \( 2(N - 1) \).

### Table V

<table>
<thead>
<tr>
<th>( N )</th>
<th>( c = 7 )</th>
<th>( c = 8 )</th>
<th>( c = 9 )</th>
<th>( c = 10 )</th>
<th>( c = 11 )</th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>1</td>
<td>2</td>
<td>8</td>
<td>32</td>
<td>114</td>
</tr>
<tr>
<td>4</td>
<td>2</td>
<td>17</td>
<td>93</td>
<td>531</td>
<td>2817</td>
</tr>
<tr>
<td>5</td>
<td>6</td>
<td>56</td>
<td>503</td>
<td>3691</td>
<td>25753</td>
</tr>
</tbody>
</table>

Fig. 7. Examples of schedules produced by the alternative broadcasting method.

**Broadcast Data**

\[ a = 2.0 \]

\[ \begin{array}{ccccccc}
T_1 & T_2 & T_3 & T_4 & T_5 & T_6 \\
\cdots & S_1 & S_2 & S_1 & S_3 & S_1 & S_2 & \cdots \\
\end{array} \]

**Time**

\[ \begin{array}{ccccccc}
t_1 & t_2 & t_3 & t_4 & t_5 & t_6 \\
\end{array} \]

**Demand**

\[ \begin{array}{ccccccc}
(\text{for } t_{d1}) & S_1 & S_2 & S_3 \\
\end{array} \]

Fig. 8. An example of implementation of the alternative broadcasting method \((a = 2.0, N = 3)\).

### B. Scenario Description

After a client demands playing a continuous media data, it starts receiving the broadcast data. The client ceases to start playing the data until it can play the data without interruption from beginning to end. The client stores the data into its buffer while it plays the data. After the client finishes playing \( S_n (n = 1, \ldots, N - 1) \), it immediately starts playing \( S_{n+1} \), which is in the buffer or is being broadcast. For example, suppose the case of \( a = 2.0 \) and \( N = 3 \) (Fig. 8). When a client demands the data at \( t_{d1} \), it receives the \( S_1 \) broadcast from \( T_3 \). Hence, the client can start playing the data from \( t_3 \). The client finishes playing \( S_1 \) at \( t_5 \). Since the client can store \( S_2 \) broadcast in \( T_2 \) until \( t_5 \), it can play \( S_2 \) continuously after it finishes playing \( S_1 \). Also, the client can store \( S_3 \) until the time to start playing \( S_3 \). In this way, the client can play the data continuously. In another case, when a client demands the data at \( t_{d2} \), it receives the \( S_1 \) broadcast from \( T_3 \). However, the client has to start playing \( S_1 \) from \( t_4 \), because if the client starts playing \( S_1 \) from \( t_3 \), the time to finish playing \( S_1 \) is \( t_5 \). Since the client cannot receive \( S_2 \) until \( t_5 \), it has to wait before playing \( S_2 \). Hence, the client waits until \( t_4 \) to be able to play the data continuously. By starting to play the data from \( t_4 \), the client can store \( S_2 \) and \( S_3 \) until the time to start playing them. Hence, the client can play the data without interruption from beginning to end.

### C. The Number of Segments and the Waiting Time

Fig. 9 shows the average waiting time by a computer simulation. Since the average waiting time is proportional to the playing time, the value of the average waiting time divided by
the playing time is shown on the vertical axis. The horizontal axis shows playback ratio $a$. “Simple” is the average waiting time under the simple repetition scheme, which I explained in Section II. “AB ($N = i$)” ($i = 3, \cdots, 7$) is the average waiting time under the alternative broadcasting method.

Although the average waiting time under the alternative broadcasting method of $N = 3$ is the shortest at $1.5 \leq a \leq 3.0$, that of $N = 4$ is the shortest at $3.0 \leq a \leq 5.0$. Thus, the $N$ that gives the shortest average waiting time depends on the playback ratio. Hence, in this paper, we propose some methods for determining $N$.

### D. Methods for Determining $N$

We propose two methods for determining $N$.

- **The mechanism dominance (MD) method:**

  In this method, even if a client starts playing the data as soon as it receives $S_1$, it can play the data without interruption from beginning to end. Hence, the advantage of the MD method is that clients need not get information on the broadcast schedule and that the client can start playing the data as soon as it receives $S_1$. That is, cases such as demanding the data at $t_{DJ}$ in Fig. 8 do not occur. To satisfy the above condition, in the MD method, $N$ is given by

  $$N = \left\lceil \frac{a - 1}{2} + 2 \right\rceil.$$  

  (10)

  This equation is discussed in Section VI-C. The average waiting time under the MD method is sometimes longer than that under the other method below (Section VI-B).

- **The waiting-time dominance (WD) method:**

  The average waiting time under this method is sometimes shorter than that under the MD method. However, when a client starts playing $S_1$ as soon as it receives $S_1$, it cannot always satisfy the continuity condition. Therefore, the client has to wait for the data to start playing after it receives $S_1$ to play the data without interruption from beginning to end. Hence, the disadvantage of the WD method is that clients need to obtain information on this waiting time. Clients can get this information by adding it to the header of segments. In the WD method, $N$ is given by

  $$N = \left\lceil \frac{a + 3}{2} \right\rceil.$$  

  (11)

  This is discussed in Section VI-D.

### VI. DISCUSSION

#### A. Comparison

To demonstrate the effectiveness of our proposed scheme, we compare the average waiting time under the division-based broadcasting scheme with that under the simple repetition scheme. Fig. 10 shows the average waiting time. To make the figure clearly understandable, the ordinate range is different from that in Fig. 8. “Minimum ($N = 2$)” is the minimum average waiting time under the division-based broadcasting scheme of $N = 2$. These minimum average waiting times are discussed in Section IV. “MD-AB” is the average waiting time under the mechanism dominance determination method of the alternative broadcasting method. “WD-AB” is the average waiting time under the waiting-time dominance determination method of the alternative broadcasting method. “Optimum ($c = N, \cdots, 11$)” is the minimum average waiting time under schedules when $c$ is changed from $N$ to 11. $N$ is the same value as WD-AB selects. This waiting time is calculated by the computer program explained in Section V.

From these figures, we can see that the alternative broadcasting scheme gives close waiting time to the optimum for values that require relative low computation.

For example, suppose that a 5 Mbps video for 30 minutes is broadcast via a 50 Mbps satellite broadcast system ($a = 10$). In this case, the average waiting time under the simple repetition scheme is $0.05 \times 30 = 1.5$ minutes. When dividing the data into two segments, the minimum average waiting time becomes $0.03 \times 30 = 0.9$ minutes. This is 42% shorter than that under the simple repetition scheme. Moreover, the average waiting time under the MD-AB method is $0.017 \times 30 = 0.51$ minutes, and that under the WD-AB method is $0.015 \times 30 = 0.45$ minutes. In the case of broadcasting using the WD-AB method, the average waiting time is 69% shorter than that under the simple repetition scheme.

#### B. The Mechanism or the Waiting-Time Dominance Determination Method

The WD-AB method determines $N$ in order to further reduce the average waiting time. For example, the average waiting time under the WD-AB method in Fig. 10 is the same as the minimum average waiting time in Fig. 9. However, since a client has to wait after receiving $S_1$ to play the data continuously, it must obtain information on how long this waiting time will be.
In the MD-AB method, a client that starts playing $S_1$ as soon as it receives $S_1$ can always play the data without interruption from beginning to end. However, the average waiting time can be longer than that under the WD method. For example, in Fig. 10, the average waiting time is longer at $a \leq 2.5, 3.5 \leq a \leq 4.5$, and so on. The MD method is used to simplify the mechanism, and the WD method is used to further reduce the waiting time.

We show the mathematical argument of each determination method below.

C. Argument for the Mechanism Dominance Method

The following theorem is established in the alternative broadcasting method.

**Theorem 3:** The condition for clients to play the data without interruption between $S_n$ and $S_{n+1}$ is equivalent to the following inequities.

- In the case of $n = 1$:
  \[ D'(2N - 3) \leq D. \]  \hspace{1cm} (12)

- In the case of $n = 2, \cdots, N - 1$:
  \[ D' \leq D. \]  \hspace{1cm} (13)

Proof is shown in Appendix I.

From Theorem 3, $N \leq (a - 1)/2 + 2$ must be established to play the data without interruption from beginning to end. Considering that a larger $N$ further reduces the average waiting time, the equation becomes $N = \lceil (a - 1)/2 + 2 \rceil$. This is same as (10). Hence, in the MD method, a client that starts playing the data as soon as it receives $S_1$ can play the data continuously. In this case, the average waiting time is equal to the time for clients to receive $S_1$. Since the interval of $S_1$ is $2D'/N$, the average waiting time $W$ is given by the following equation:

\[ W = \frac{2D'}{2N} = \frac{D'}{N}. \]  \hspace{1cm} (14)

D. Argument for the Waiting-Time Dominance Determination Method

It is supposed that $D'(2N - 3) < D_m \leq D'(2N' - 3)$ is established for all $N$ when the average waiting time under $N$ becomes longer than that under $N' = N + 1$. This is supposed from Fig. 9. Namely, when the number of segments is $N$, a client that starts playing $S_1$ as soon as it receives $S_1$ can play the data without interruption, but when the number of segments is $N'$, it cannot. For example, in Fig. 9, the playback ratio at which the average waiting time under $N = 4$ becomes shorter than that of $N = 3$ is $9.6$. This satisfies $7 < \alpha < 13$. Accordingly, we can obtain the average waiting time in this bound. When the number of segments is $N$, $D'(2N - 3) \leq D$ is established. Hence, by Theorem 3, the average waiting time $W_N$ is

\[ W_N = \frac{D'}{N}. \]  \hspace{1cm} (15)

Next, we show the average waiting time under $N'$. Here, the following theorem is established.

**Theorem 4:** In the case of $D'(2N - 2i + 1) \leq D \leq D'(2N - 2i + 3)(i = 2, \cdots, N - 1)$, the average waiting time $W$ is

\[ W = \frac{1}{2N(N - 1)} (D(-2i + 5) - D'(2i^2 - 4i(N + 1) + 8N - 1)), \]  \hspace{1cm} (16)

Proof is shown in Appendix 2. By Theorem 4, in the case of $D' + 2D'(N - 2) < D \leq D' + 2D'(N' - 2)$, the average waiting time $W_N'$ is:

\[ W_N' = \frac{1}{2N'(N'-1)} (-D + D'(4N' - 5)) \]  \hspace{1cm} (17)

Compared with $W_N$, when the following equation is established, the average waiting time under $N + 1$ is shorter than that under $N$:

\[ N \leq \frac{a + 3}{2}, \]  \hspace{1cm} (18)

Considering that $N$ is incremented when (18) is established, $N$ is given by $N = \lceil (a + 3)/2 \rceil$ to reduce the average waiting time effectively. This is same as (11). Hence, a shorter average waiting time than the MD method is given by (11).

VII. RELATED WORK

Various strategies have been proposed to reduce the waiting time under the condition that clients can play data without interruption from beginning to end. In pyramid broadcasting (PB) [16], a continuous media data is divided into $N$ segments, and $N$ channels are used. Let $D_l$ be the time to broadcast the $i^{th}$ segment. The data is divided so that $D_l = D_1 \times \alpha(i - 1)$, where, $\alpha$ is 1 or more. The server broadcasts the $i^{th}$ segment via the $i^{th}$ channel. Since the interval of broadcasting precedent segments is shorter than the other, the first segment is broadcasted frequently. Hence, clients’ waiting time is reduced, because they have many opportunities to receive the first segment. In this scheme, however, the greater the value of $N$ becomes, the larger the data size of the $i^{th}$ segment is. Therefore, a larger buffer capacity is required.

Skyscraper broadcasting (SB) [6] considers this problem. Segments with a longer playing time than $D_1 \times W$ are divided so that its playing time becomes $D_1 \times W$. By modifying $W$, the buffer requirement and the average waiting time are voluntarily given.

Dynamic skyscraper broadcasting [3] dynamically changes broadcast data based on the SB scheme. This improves the performance of SB when clients infrequently request data.

The permutation-based pyramid broadcasting scheme [1] divides $N$ channels into $p$ sub-channels. A server broadcasts the $i^{th}$ segment using $p$ sub channels. The beginning of the broadcast cycle of each sub-channel is shifted. As a result, the storage requirements, the disk transfer rate, and the waiting time are improved compared to PB. This scheme proposes time-division multiplexing, which employs a single channel. However, the paper does not detail a technique for reducing the waiting time of clients with a single channel.

To further reduce the waiting time of clients, harmonic broadcasting (HB) [8] has been proposed. In this scheme, the $i^{th}$ segment is divided into $i$ sub-segments. A server broadcasts $i$ sub-segments of the $i^{th}$ segment of each data via the $i^{th}$ channel. The bandwidth of the $i^{th}$ channel is $1/i$ of the first channel, while the bit rate of the first channel is the same as the consumption rate. Hence, we can say that the playback ratio of the $i^{th}$ channel is $1/i$. However this work did not cover a method of reducing the waiting time for clients with a single channel. This scheme always requires a buffer on the client side.

In fast broadcasting (FB) [9], clients that do not have a buffer can play data by waiting longer. Regarding FB, methods for channel transition have been studied [15].
These strategies reduce the waiting time by broadcasting continuous media data with multiple channels. Because of this, clients are forced to receive data from multiple channels concurrently, and the server has to control multiple channels. Hence, their mechanisms are more complicated than broadcasting data via a single channel. Regarding constructing an actual system, neither wireless LAN nor satellite broadcast systems can receive data from multiple channels concurrently.

VIII. CONCLUSION AND FUTURE WORK

In this paper, we proposed the division-based broadcasting scheme to reduce the waiting time of clients with a single channel. In our approach, the number of segments \( N \) and the broadcast schedule affect the waiting time.

In the case of dividing data into two segments of equal size, we presented a schedule that provides the minimum average waiting time. The number of the first segment, \( S_1 \), included in the broadcast schedule is decided by the playback ratio (the data playing time divided by the time required to broadcast the data), and the last segment \( S_2 \) is included only once. In the case of broadcasting a 5 Mbps video for 30 minutes via a 50 Mbps satellite broadcast system, the minimum average waiting time under \( N = 2 \) is 53 seconds. This is 42% shorter than that in the case of broadcasting the data without division. Hence, dividing the data into only two segments is sufficiently realistic.

For the case where it is necessary to divide data into more than two segments, we propose the alternative broadcasting method. In the alternative broadcasting method, the \( N \) that gives the shortest average waiting time depends on the playback ratio. Hence, we proposed two methods for determining \( N \). For example, in the aforesaid case, the average waiting time under the mechanism dominance determination method of alternative broadcasting method (MD-AB) is 30 seconds, and that under the waiting-time dominance determination method (WD-AB) is 28 seconds. In the case of broadcasting using the WD-AB method, the average waiting time is 69% shorter than that under the simple repetition scheme.

In the future, we plan to study other heuristics and delivering multiple data with a single channel.

APPENDIX

A. Proof of Theorem 3

Let \( I_n \) be the time from starting broadcasting \( S_n(n = 1, \cdots, N - 1) \) to starting broadcasting \( S_{n+1} \) (Fig. 11). \( I_n \) are given by the following equations:

- In the case of \( n = 1 \):
  
  \[ I_1(i) = \begin{cases} \frac{D}{N} & (i = 1) \\ \frac{D}{N} + \frac{2D(N - i)}{N} & (i = 2, \cdots, N - 1) \end{cases} \]  

- In the case of \( n = 2, \cdots, N - 1 \):
  
  \[ I_n = \frac{D}{N} \]  

B. Proof of Theorem 4

In the case of \( D'/(2N - 2i + 1) \leq D \leq D'/(2N - 2i + 3)(i = 2, \cdots, N - 1) \), (13) is always established. Hence, by Theorem 3, if the client can play \( S_2 \) after \( S_1 \) continuously, it can play the data without interruption from beginning to end. In this case, the average waiting time is calculated by the following three cases:

- In the case of demanding data in \( T_1 \):
  
  Since a client can store \( S_2 \) broadcast from \( T_2 \), the client that starts playing \( S_1 \) as soon as it receives \( S_1 \) from \( T_3 \) can play the data without interruption from beginning to end. The average waiting time is \( D'/2N \).

- In the case of demanding data in from \( T_2 \) to \( T_{2N - 1} \):
  
  A client that starts playing \( S_1 \) as soon as it receives \( S_1 \) cannot play \( S_2 \) continuously. Hence the client has to wait. In this case, the average waiting time is \( -D/N + (4N - 2i + 1)D'/2N \).

- In the case of demanding the data in from \( T_{2N - 3} \) to \( T_{2N - 2} \):
  
  A client that starts playing \( S_1 \) as soon as it receives \( S_1 \) can play the data continuously. The average waiting time is \( D'/N \).

Accordingly, the average waiting time \( W \) is:

\[ W = \frac{1}{2N(N - 1)} (D(-2i + 5) - D' \times (2i^2 - 4i(N + 1) + 8N - 1)) \]  

REFERENCES


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