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Kyoto University
Formulas with only one atomic formula in Grzegorczyk logic and provability logic *

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Abstract. Here we discuss the set \( S(p) \) of the formulas having only one atomic formula \( p \) in Grzegorczyk logic \( \text{Grz} \) and the set \( S(\bot) \) of the formulas having only one atomic formula \( \bot \) in provability logic \( \text{GL} \). We give an inductive construction of representatives in the quotient set \( S(p)/\equiv_{\text{Grz}} \) modulo the provability of \( \text{Grz} \). On the other hand, in Boolos [1], it was shown that any formula \( A \in S(\bot) \) is equivalent to some truth-functional combination of formulas of the form \( \Box \bot \) in \( \text{GL} \). We modify it and give representatives in the quotient set \( S(\bot)/\equiv_{\text{GL}} \), which correspond to the representatives for \( \text{Grz} \). By these representatives, we clarify structures \( \langle S(p)/\equiv_{\text{Grz}}, \leq_{\text{Grz}} \rangle \) and \( \langle S(\bot)/\equiv_{\text{GL}}, \leq_{\text{GL}} \rangle \), where \( \leq_{\text{L}} \) is the derivation in \( \text{L} \in \{ \text{Grz, GL} \} \). Comparing these two structures, we also give a way to express the \( \text{GL} \)-provability of formulas in \( S(\bot) \) in \( \text{Grz} \). In spite of the simplicity of \( S(p) \) and \( S(\bot) \), it is worth considering since the quotient sets are infinite. There is few result on such structures with infinite quotient sets. One result was given in Nishimura [7] in intuitionistic propositional logic, however, the target set of formulas are also simple, with only two atomic formulas \( p \) and \( \bot \). Shehtman [11] considered more general structure for \( \text{Grz} \), however, in our simple case, our results have more information.

1 Introduction

In this section, we introduce Grzegorczyk logic \( \text{Grz} \) and provability logic (or Gödel-Löb logic) \( \text{GL} \). We use lower case Latin letters \( p, q, \ldots \), possibly with suffixes, for propositional variables. Formulas are defined inductively, as usual, from the propositional variables and \( \bot \) (contradiction) by using logical connectives \( \wedge \) (conjunction), \( \vee \) (disjunction), \( \supset \) (implication) and \( \Box \) (necessitation). We use upper case Latin letters \( A, B, \ldots \), possibly with suffixes, for formulas. The expression \( \Box^n A \) is defined inductively as \( \Box^0 A = A \) and \( \Box^{k+1} A = \Box(\Box^k A) \). For a finite \( S \) of formulas, we put \( \Box S = \{ \Box A \mid A \in S \} \). We fix the enumeration \( \text{ENU} \) of formulas. For a finite non-empty set \( S \) of formulas, the expressions \( \bigwedge S \) and \( \bigvee S \) denotes the formulas \( A_1 \wedge A_2 \wedge \cdots \wedge A_n \) and \( A_1 \vee A_2 \vee \cdots \vee A_n \), respectively.

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where \( S = \{ A_1, \ldots, A_n \} \) and \( A_i \) occurs earlier than \( A_j \) in ENU if \( i < j \). Also we put 
\( \bigwedge \emptyset = \perp \) and \( \bigvee \emptyset = \perp \). The depth \( d(A) \) of a formula \( A \) is defined inductively as 
\[ d(D) = 0 \quad \text{for any atomic formula } D, \quad d(B \land C) = d(B \lor C) = d(B \supset C) = \max\{d(B), d(C)\} \]
and \( d(\square B) = d(B) + 1 \). Let \( D \) be an atomic formula in \( \{ p, \perp \} \). By \( S(D) \), we mean the set of formulas constructed from \( D \) by using \( \land, \lor, \supset \) and \( \square \). We put \( S^n(D) = \{ B \in S(D) \mid d(B) \leq n \} \).

By \( \text{Grz} \), we mean the smallest set of formulas containing all the tautologies, and the axioms \( K : \square (A \supset B) \supset (\square A \supset \square B), T : \square A \supset A, \text{grz} : \square (\square (A \supset \square A) \supset A) \supset \square A \) (Grzegorczyk axiom), and closed under modus ponens and necessitation. By \( \text{GL} \), we mean the smallest set of formulas containing all the tautologies, and the axioms \( K \) and \( L : \square (\square A \supset A) \supset \square A \) (Löb's axiom), and closed under modus ponens and necessitation.

Let \( \text{L} \) be either \( \text{Grz} \) or \( \text{GL} \). We write \( A \equiv_{\text{L}} B \) if \( (A \supset B) \land (B \supset A) \in L \). Also for any equivalent classes \( [A] \) and \( [B] \), we put \( [A] \leq [B] \) modulo the provability of \( \text{L} \). We write \( [A] \leq [B] \) if \( A \supset B \in L \). We also use this kind of notations for other logics.

A Kripke model is a triple \( M = \langle W, R, P \rangle \), where \( W \) is a non-empty set, \( R \) is a binary relation on \( W \) and \( P \) is a mapping from the set of propositional variables to \( 2^W \). The truth valuation \( (M, \alpha) \models A \), a formula \( A \) is true at \( \alpha \in W \) in \( M \), is defined by an induction on \( A \) in the usual way. The expression \( M \models A \) denotes \( (M, \alpha) \models A \) for every \( \alpha \in W \). Since \( P(p) = \{ \alpha \mid (M, \alpha) \models p \} \), we can extend the mapping \( P \) to the set of formulas as \( P(A) = \{ \alpha \mid (M, \alpha) \models A \} \).

\begin{lemma}
(1) \( A \in \text{Grz} \) iff \( M \models A \) for any finite Kripke model \( M = \langle W, R, P \rangle \) with partial orders
(i. e. \( W \) is finite and \( R \) is a partial order).

(2) \( A \in \text{GL} \) iff \( M \models A \) for any finite Kripke model \( M = \langle W, R, P \rangle \) with strict partial orders
(i. e. \( W \) is finite and \( R \) is a strict partial order).
\end{lemma}

\begin{lemma}
For \( L \in \{ \text{Grz}, \text{GL} \} \),
(1) \( \square (A \land B) \equiv_{L} \square A \land \square B \),
(2) \( \square^k A \supset \square^{k+4} A \in L \) for \( k > 0, i > 0 \),
(3) \( A \supset B \in L \) implies \( \square A \supset \square B \in L \),
(4) \( \square A \supset B \in L \) implies \( \square A \supset \square B \in L \).
\end{lemma}

2 The structure \( \langle S(p) \rangle / \equiv_{\text{Grz}}, \leq_{\text{Grz}} \)

In this section, we construct representatives of equivalent classes in \( S^n(p) / \equiv_{\text{Grz}} \) and clarify the structure \( \langle S(p) \rangle / \equiv_{\text{Grz}}, \leq_{\text{Grz}} \). It is known, however, structures \( \langle S^n(p) / \equiv_{\text{Grz}}, \leq_{\text{Grz}} \rangle \) are boolean(cf. [3]). Also the quotient set is finite. So, we have only to construct representatives of generators of the boolean.

\[ \text{cf. K. Segerberg [10] and A. Chagrov and M. Zakharyaschev [3].} \]
Definition 2.1. A list $F_0, F_1, \ldots$ of formulas are defined inductively as $F_0 = p$ and $F_{k+1} = F_k \supset \Box F_k$.

Lemma 2.2. For $k \geq 0$,

1. $\Box F_k \supset \Box F_{k+i} \in \text{Grz}$ ($i \geq 0$),
2. $F_k \land F_{k+1} \equiv_{\text{Grz}} \Box F_k$,
3. $\Box (\Box F_{k+1} \supset \Box F_k) \equiv_{\text{Grz}} \Box F_k$,
4. $F_{k+2} \land (\Box F_{k+1} \supset \Box F_k) \equiv_{\text{Grz}} F_k$,
5. $\Box F_k \equiv_{\text{Grz}} \bigwedge \{(\Box F_{k+\ell+1} \supset \Box F_{k+\ell} | \ell = 0, 1, \ldots, i - 1) \cup \{F_{k+i}, F_{k+i+1}\}\}$ ($i \geq 0$).

Proof. By Lemma 1.2(4), we have (1). By the axiom $T$, we have (2).

We show (3). By the axiom $T$, we have $(\Box F_{k+1} \supset \Box F_k) \supset (\Box F_{k+1} \supset F_k) \in \text{Grz}$, and using Lemma 1.2(3), $\Box (\Box F_{k+1} \supset \Box F_k) \supset (\Box F_{k+1} \supset F_k) \in \text{Grz}$. By the axiom $grz$, we have $\Box (\Box F_{k+1} \supset F_k) \supset \Box F_k \in \text{Grz}$. Hence $\Box (\Box F_{k+1} \supset \Box F_k) \supset \Box F_k \in \text{Grz}$. The converse is from Lemma 1.2(4).

We show (4). Let be that $A = F_{k+2} \land (\Box F_{k+1} \supset \Box F_k)$. We note that $F_{k+1} \supset (A \supset \Box F_k)$ is a tautology. Using the axiom $T$, we have $F_{k+1} \supset (A \supset F_k) \in \text{Grz}$. We also note that $F_k \supset (A \supset F_k)$ and $F_k \lor F_{k+1}$ are tautologies. Hence we have $A \supset F_k \in \text{Grz}$. On the other hand, we note that $(F_k \land F_{k+1}) \supset \Box F_k$ is a tautology, and using (1) and the axiom $T$, we have $(F_k \land F_{k+1}) \supset \Box F_{k+1} \in \text{Grz}$ and $(F_k \land \Box F_{k+1}) \supset \Box F_k \in \text{Grz}$. So, $F_k \supset (F_{k+1} \supset \Box F_{k+1}) = F_k \supset F_{k+2} \in \text{Grz}$ and $F_k \supset (\Box F_{k+1} \supset \Box F_k) \in \text{Grz}$.

We show (5). By (1), we have $\Box F_k \equiv_{\text{Grz}} \bigwedge \{(\Box F_{k+\ell+1} \supset \Box F_{k+\ell} | \ell = 0, 1, \ldots, i - 1) \cup \{F_{k+i}, F_{k+i+1}\}\}$, and using (2), we obtain (5).

Definition 2.3. The sets $G_n$ ($n = 0, 1, 2, \ldots$) of formulas are defined as $G_0 = \{F_0\}$ and $G_{k+1} = \{F_k, F_{k+1}, \Box F_k \supset \Box F_{k-1}, \ldots, \Box F_1 \supset \Box F_0\}$.

Theorem 2.4.

1. $S^n(p)/ \equiv_{\text{Grz}} = \{[\bigwedge S] | S \subseteq G_n\}$.
2. For any subsets $S_1$ and $S_2$ of $G_n$, $S_1 \subseteq S_2$ iff $[\bigwedge S_2] \leq_{\text{Grz}} [\bigwedge S_1]$.

To prove the theorem above, we need some preparations.

By $N$, odd and even, we mean the set of integers $0, 1, 2, \ldots$, the set of odd numbers $1, 3, 5, \ldots$, and the set of even numbers $0, 2, 4, \ldots$, respectively. We define the Kripke model $M_{\text{Grz}} = \langle W_{\text{Grz}}, R_{\text{Grz}}, P_{\text{Grz}}\rangle$, where $W_{\text{Grz}} = N$, $R_{\text{Grz}} = \{(k, \ell) | k \geq \ell\}$ and $P_{\text{Grz}}(p) = \text{odd}$.

Lemma 2.5.

1. $(M_{\text{Grz}}, k) \models F_i$ iff either $i > k$ or $i + k \in \text{odd}$,
2. $(M_{\text{Grz}}, k) \models \Box F_i$ iff $i > k$.

Proof. We use an induction on $i$.

Basis ($i = 0$): We note that $0 > k$ does not hold. So, by the definition of $P_{\text{Grz}}$, we have $1$. By $kR_{\text{Grz}}0$ and $(M_{\text{Grz}}, 0) \not\models F_0$, we have $(M_{\text{Grz}}, k) \not\models \Box F_0$, and so, we obtain (2).

Induction step ($i > 0$): We show (1). By the definition of $\models$, we have $(M_{\text{Grz}}, k) \not\models F_i$ iff $(M_{\text{Grz}}, k) \models F_{i-1}$ and $(M_{\text{Grz}}, k) \not\models \Box F_{i-1}$. 

Using the induction hypothesis,

\[(M_{Grz}, k) \not\models F_i \text{ iff } \text{"}i - 1 > k \text{ or } i - 1 + k \in \text{odd} \text{"} \text{ and } i - 1 \leq k.\]

We note that \(i - 1 > k\) and \(i - 1 \leq k\) do not hold simultaneously. So,

\[(M_{Grz}, k) \not\models F_i \text{ iff } i - 1 + k \in \text{odd} \text{ and } i - 1 \leq k.\]

If \(i - 1 = k\), then \(i - 1 + k = k + 1 - 1 + k = 2k \not\in \text{odd}\). So,

\[(M_{Grz}, k) \not\models F_i \text{ iff } i - 1 + k \in \text{odd} \text{ and } i - 1 < k.\]

So,

\[(M_{Grz}, k) \not\models F_i \text{ iff } i + k \not\in \text{odd} \text{ and } i \leq k.\]

We show (2). Suppose that \(i \not< k\). Then we have \(k \geq i\), and so \(kR_{Grz}i\). On the other hand, by (1), we have \((M_{Grz}, i) \not\models F_i\). Hence we obtain \((M_{Grz}, k) \not\models \Box F_i\). Suppose that \(i > k\). If \(kR_{Grz}k'\), then we have \(i > k'\), and using (1), \(k' \models F_i\). Hence we obtain \((M_{Grz}, k) \models \Box F_i\).

**Lemma 2.6.**

1. None of the formulas in \(G_n\) is provable in \(Grz\).
2. \(S = \{F_{2n}\} \cup \{\Box F_{2k+1} \supset \Box F_{2k} \mid k = 0, 1, \cdots, n - 1\} \) is a subset of \(G_{2n}\) and a subset of \(G_{2n-1}\), and \(\land S \equiv_{Grz} p\).
3. For any \(A, B \in G_n\), \(A \not\equiv B \text{ implies } A \lor B \in Grz\).
4. For any \(A, B \in G_n\), \(A \not\equiv B \text{ implies } B \equiv_{Grz} A \lor B\).
5. \(n \not= 0 \text{ implies } \land G_n \equiv_{Grz} \Box p\).

**Proof.** For (1). By Lemma 2.5(2), we have \((M_{Grz}, i) \models \Box F_{i+1}\) and \((M_{Grz}, i) \not\models \Box F_i\), and so, \((M_{Grz}, i) \not\models \Box F_{i+1} \supset \Box F_i\). Also by Lemma 2.5(1), we have \((M_{Grz}, i) \not\models F_i\). Using Lemma 1.1(1), we obtain (1).

For (2). By an induction on \(n\) and Lemma 2.2(4).

For (3). We use an induction on \(n\). If \(n = 0\), then (3) is clear. Suppose that \(n > 0\). We note that \(F_n \lor F_{n-1} \text{ and } (\Box F_{n-1} \supset \Box F_{n-2}) \lor F_{n-1}\) are tautologies. By Lemma 2.2(1), \(F_n \lor (\Box F_{k+1} \supset \Box F_k) (k = 0, 1, \cdots, n - 2) \text{ and } (\Box F_{n-1} \supset \Box F_{n-2}) \lor (\Box F_{k+1} \supset \Box F_k) (k = 0, 1, \cdots, n - 3) \) are tautologies. The other cases can be shown by the induction hypothesis. Hence we obtain (3).

For (4). By (3).

For (5). By Lemma 2.2(5).

**Lemma 2.7.** Let \(S_1\) and \(S_2\) be subsets of \(G_n\). Then

1. \((\land S_1) \lor (\land S_2) \equiv_{Grz} \land (S_1 \cup S_2)\),
2. \((\land S_1) \lor (\land S_2) \equiv_{Grz} \land (S_1 \cap S_2)\),
3. \((\land S_1) \lor (\land S_2) \equiv_{Grz} \land (S_2 - S_1)\),
4. if \(S_1 \not= \emptyset\), then \(\Box (\land S_1) \equiv_{Grz} \Box F_k\), where \(k = \min\{i \mid F_i \in S_1\} \cup \{i \mid \Box F_{i+1} \supset \Box F_i \in S_1\}\).
Proof. (1) is from associative law and commutative law of $\land$. For (2) and (3), we use Lemma 2.6(3) and Lemma 2.6(4), respectively. We show (4). By Lemma 1.2(1) and Lemma 2.2(3),

$$
\Box(\bigwedge S_1) \equiv_{\text{Grz}} \bigwedge (\Box S_1)
$$

$$
\equiv_{\text{Grz}} \bigwedge (\{\Box F_i \mid F_i \in S_1\} \cup \{\Box(\Box F_{i+1} \supset \Box F_i) \mid \Box F_{i+1} \supset \Box F_i \in S_1\})
$$

$$
\equiv_{\text{Grz}} \bigwedge (\{\Box F_i \mid F_i \in S_1\} \cup \{\Box F_i \mid \Box F_{i+1} \supset \Box F_i \in S_1\}).
$$

Using $k = \min(\{i \mid F_i \in S_1\} \cup \{i \mid \Box F_{i+1} \supset \Box F_i \in S_1\})$, and Lemma 2.2(1), we obtain \(\Box(\bigwedge S_1) \equiv_{\text{Grz}} \Box F_k\).

Lemma 2.8. Let $A$ be a formula in $S^n(p)$. Then there exists a subset $S$ of $G_n$ such that $A \equiv_{\text{Grz}} \bigwedge S$.

Proof. We use an induction on $A$. If $A = p$, then by Lemma 2.6(2) we obtain the lemma. If $A \neq p$, then by the induction hypothesis, Lemma 2.7 and Lemma 2.2(5), we obtain the lemma.

Proof of Theorem 2.4. (1) is from Lemma 2.8. The "only if" part of (2) is clear. We show the "if part" of (2). Suppose that $[\bigwedge S_2] \leq_{\text{Grz}} [\bigwedge S_1]$ and $S_1 \not\subseteq S_2$. By $S_1 \not\subseteq S_2$, there exists a formula $A$ in $S_1 - S_2$. Using $[\bigwedge S_2] \leq_{\text{Grz}} [\bigwedge S_1]$, we have $\bigwedge S_2 \supset A \in \text{Grz}$. Since $A \not\in S_2$, using Lemma 2.6(4), we have $\bigwedge S_2 \supset A \equiv_{\text{Grz}} A$, and so, we have $A \in \text{Grz}$. This is in contradiction with Lemma 2.6(1).

Theorem 2.4 provides representatives of $S^n(p) / \equiv_{\text{Grz}}$. Next, we clarify the structure $\langle S(p)/ \equiv_{\text{Grz}}, \leq_{\text{Grz}} \rangle$. We first introduce an exact model, which is useful to clarify this kind of structures if the quotient set is finite. Let $S$ be a set of formulas closed under $\supset$. A Kripke model $\langle W, R, P \rangle$ is said to be an exact model for $S$ in a logic $L$ if the following two conditions hold:

1. $P$ maps $S$ onto $2^W$,
2. for any $A \in S$, $A \in L$ iff $P(A) = W$.

The condition (2) above is equivalent to

3. for any $A \in S$, $A \supset B \in L$ iff $P(A) \subseteq P(B)$.

So, $P$ is a homomorphism from $\langle S, R \rangle$ to $\langle 2^W, \subseteq \rangle$, where $R = \{(A, B) \mid A \supset B \in L\}$. If an exact model for a set $S$ in a logic $L$ is given, then we can construct a structure isomorphic to $\langle S/ \equiv_L, \leq_L \rangle$ as follows. By (3), for any $B \in [A]_{\equiv_L}$, we have $P(A) = P(B)$. Hence we can define a one-to-one mapping $f$ from $S/ \equiv_L$ to $2^W$ as $f([A]) = P(A)$. By (1), $f$ is onto, and so, an isomorphism. Hence $\langle S/ \equiv_L, \leq_L \rangle$ is isomorphic to $\langle 2^W, \subseteq \rangle$.

So, giving a concrete exact model for $S$ in $L$ is an effective way to clarify a structure $\langle S/ \equiv_L, \leq_L \rangle$. Bruijn [2], Hendriks [6] used this model for the set of disjunction free formulas with finite number of atomic formulas in intuitionistic propositional logic, and gave precise description on the structure, while Diego [4], Urquhart [12] and [8] treated the same structure. [9] also used this model for the set of disjunction free formulas with finite number of atomic formulas in a normal modal logic, called propositional lax logic, and gave precise description on the structure. Also exact models are useful to clarify such kind of structures if quotient sets are finite.
However, if $S \not\equiv_{L}$ is not finite, then there is no exact model for $S$ in $L$. Suppose that there is an exact model. Then $(S \not\equiv_{L}, \leq_{L})$ is isomorphic to $(2^{W}, \subseteq)$. So, $2^{W}$ is not finite, and neither is $W$. So, $2^{W}$ is not countable. On the other hand, since $S$ is countable, so are the quotient set and $2^{W}$. This is a contradiction. Although Nishimura [7] clarified the structure, with infinite quotient set, for the set of formulas with only two atomic formulas $p$ and $\bot$ in intuitionist logic, results on such structures with infinite quotient sets are few.

The quotient set $S(p)/\equiv_{Grz}$ is not finite. This makes our problem difficult. There is no exact model for $S(p)$ in $Grz$. So, we modify exact model in order to treat such infinite case. The idea has already used in general frame described in [3]. The structure $(W, R, Q)$ in the definition below is called a general frame in [3].

**Definition 2.9.** A structure $(W, R, P, Q)$ is said to be a general exact model for a set $S$ of formulas in a logic $L$ if the following four conditions hold:

1. $(W, R, P)$ is a Kripke model,
2. $Q$ is a subset of $2^{W}$
3. $Im(P)(= \{P(A) \mid A \in S\}) = Q$,
4. for any $A \in S$, $A \in L$ iff $P(A) = W$.

Similarly to the description of exact models, we have

**Lemma 2.10.** Let $(W, R, P, Q)$ be a general exact model for a set $S$ of formulas in a logic $L$. Then

1. we can define a one-to-one mapping $f$ from $S/\equiv_{L}$ to $Q$ as $f([A]) = P(A)$, and $f$ is an isomorphism from $(S/\equiv_{L}, \leq_{L})$ to $(Q, \subseteq)$,
2. $(S/\equiv_{L}, \leq_{L})$ is isomorphic to $(Q, \subseteq)$.

[11] constructed the structure $(W, R, P)$ satisfying the first two conditions in Definition 2.9 and

5. $\{\{w\} \mid w \in W\} \subseteq Im(P),$

for the set $S$ of formulas having only $n$ propositional variables $p_{1}, \cdots, p_{n}$. However, he did not clarify $Im(P)(= Q)$, in a sense, and so, we have not known the structure $(S/\equiv_{L}, \leq_{L})$. After the proof of Theorem 2.13, we will give an example as this reason.

We will clarify $(S(p)/\equiv_{Grz}, \leq_{Grz})$ by giving a concrete general exact model for $S(p)$ in $Grz$. Theorem 2.4 and Lemma 2.5 are useful for it. We put

$Q_{fn}(S) = \{S_{1} \mid S_{1}$ is a finite subset of $S\},$

$\text{odd}_{2k+1} = \{i \in \text{odd} \mid i \geq 2k + 1\},$

$\text{even}_{2k} = \{i \in \text{odd} \mid i \geq 2k\},$

$N_{k} = \{i \mid i \geq k\},$

$Q_{inf} = \{\text{odd}_{2k+1} \mid k \in N\} \cup \{\text{even}_{2k} \mid k \in N\} \cup \{N_{k} \mid k \in N\},$

$Q_{Grz} = \{S_{1} \cup S_{2} \mid S_{1} \in Q_{fn}(N), S_{2} \in Q_{inf} \cup \{\emptyset\}\}.$

**Lemma 2.11.** $S, T \in Q_{inf}$ implies $S \cap T \in Q_{Grz}$.

**Proof.** Below, we only show the case that $S \in \{\text{odd}_{2k+1} \mid k \in N\}$. The other cases can be shown similarly.
odd_{2k+1} \cap \text{odd}_{2\ell+1} = \text{odd}_{\max\{2k+1,2\ell+1\}}
odd_{2k+1} \cap \text{even}_{2\ell} = \emptyset,
odd_{2k+1} \cap N_{2\ell+1} = \text{odd}_{\max\{2k+1,2\ell+1\}}.
odd_{2k+1} \cap N_{2\ell} = \begin{cases} odd_{2k+1} & \text{if } 2k + 1 > 2\ell \\
2\ell \} \cup \text{odd}_{2k+1} & \text{if } 2k + 1 < 2\ell. \end{cases}

Lemma 2.12. \mathcal{Q}_{Grz} is closed under \cap.
Proof. Let S and T be any sets in \mathcal{Q}_{Grz}. Then there exist S_1, T_1 \in Q_{fin}(N) and S_2, T_2 \in Q_{inf} \cup \{\emptyset\} such that S = S_1 \cup S_2 and T = T_1 \cup T_2. So,

S \cap T = (S_1 \cap S_2) \cap (T_1 \cup T_2) = (S_1 \cap (T_1 \cup T_2)) \cup (S_2 \cap T_1) \cup (S_2 \cap T_2).

We put U = (S_1 \cap (T_1 \cup T_2)) \cup (S_2 \cap T_1), and note that U \in Q_{fin}(N). If either S_2 or T_2 is \emptyset, then we have S \cap T = U \cup \emptyset \in \mathcal{Q}_{Grz}. So, we assume that S_2, T_2 \in Q_{inf}. By Lemma 2.10, there exist U_1 \in Q_{fin}(N) and U_2 \in Q_{inf} such that S_2 \cap T_2 = U_1 \cup U_2. Clearly, (U \cup U_1) \in Q_{fin}(N), and so, S \cap T = U \cup (S_2 \cap T_2) = (U \cup U_1) \cup U_2 \in \mathcal{Q}_{Grz}. \square

Theorem 2.13. The structure \langle W_{Grz}, R_{Grz}, P_{Grz}, Q_{Grz}\rangle is a general exact model for S(p) in Grz.

Proof. The conditions (1) and (2) in Definition 2.9 are clear. So, it is sufficient to show the following:

(3) Im(P_{Grz}) = \mathcal{Q}_{Grz},
(4) A \in Grz iff P_{Grz}(A) = W_{Grz}.

We show (3). Suppose that S \in Im(P_{Grz}). Then there exists a formula A such that P_{Grz}(A) = S. Let be that n = max\{d(A), 1\}. Then using Lemma 2.8, there exists a subset S_1 of G_n such that A \equiv_{Grz} \bigwedge S_1. So,

P_{Grz}(A) = \{k | (M_{Grz}, k) \models \bigwedge S_1\} = \{k | (M_{Grz}, k) \models B, \text{ for any } B \in S_1\}

= \bigcap_{B \in S_1} \{k | (M_{Grz}, k) \models B\} = \bigcap_{B \in S_1} P_{Grz}(B).

By Lemma 2.5, we have

(5) P_{Grz}(\Box F_{k+1} \supset \Box F_k) = N - \{k\} = \{i | 0 \leq i < k\} \cup N_{k+1} \in \mathcal{Q}_{Grz},
(6) P_{Grz}(F_{2k}) = \{i | i < 2k\} \cup \text{odd}_1 \in \mathcal{Q}_{Grz},
(7) P_{Grz}(F_{2k+1}) = \{i | i < 2k + 1\} \cup \text{even}_0 \in \mathcal{Q}_{Grz}.

So, for any formula B \in G_n, we have P_{Grz}(B) \in \mathcal{Q}_{Grz}, and using Lemma 2.11, we obtain S = P_{Grz}(A) \in \mathcal{Q}_{Grz}.

Suppose that S \in \mathcal{Q}_{Grz}. Then there exist S_1 \in Q_{fin}(N) and S_2 \in Q_{inf} \cup \{\emptyset\} such that S = S_1 \cup S_2. On the other hand, by Lemma 2.5, (5), (6) and (7), we have

P_{Grz}(\Box p) = \emptyset,
P_{Grz}(\Box F_{k+1} \supset \Box F_k) \supset \Box p) = \{k\},
P_{Grz}(F_{2k} \supset \Box p) = N - \{i | i < 2k\} \cup \text{odd}_1 = \text{even}_{2k},
P_{Grz}(F_{2k+1} \supset \Box p) = N - \{i | i < 2k + 1\} \cup \text{even}_0 = \text{odd}_{2k+1}.

Also,
P_{Grz}(F_{2k} \supset \Box p) \lor (F_{2k+1} \supset \Box p)) = N_{2k},
Here we note that $S_2 = P_{Grz}(B)$ for some $B \in S(p)$. We also note for any formulas $C$ and $D$,

$$P_{Grz}(C) \cup P_{Grz}(D) = \{ i \mid (M_{Grz}, i) \models C \} \cup \{ i \mid (M_{Grz}, i) \models D \}$$

$$= \{ i \mid (M_{Grz}, i) \models C \lor D \} = P_{Grz}(C \lor D).$$

Hence

$$S = S_1 \cup S_2 = \left( \bigcup_{k \in S_1} \{ k \} \right) \cup P_{Grz}(B) = \left( \bigcup_{k \in S_1} P_{Grz}(\square F_{k+1} \supset \square F_k \supset \square p)) \cup P_{Grz}(B) \right)$$

$$= P_{Grz}(\bigvee \{ \square F_{k+1} \supset \square F_k \supset \square p \mid k \in S_1 \} \lor B).$$

If $S_1 = \emptyset$, then the set above is $P_{Grz}(B)$, and hence $S \in Im(P_{Grz})$; if not, we have

$$\bigvee \{ \square F_{k+1} \supset \square F_k \supset \square p \mid k \in S_1 \} \in S(p),$$

and hence we can see that the set $\{3n \mid n \in \mathbb{N}\}$ does not belongs to $Q(= Im(P))$, while from [11] we can’t see it directly.

**Corollary 2.14.** $\langle S(p)/ \equiv_{Grz}, \leq_{Grz} \rangle$ is isomorphic to $\langle Q_{Grz}, \subseteq \rangle$.

Also by sketching the proof of Theorem 2.13, we also have

**Theorem 2.15.** Let be that $M^n_{Grz} = \langle W^n_{Grz}, R^n_{Grz}, P^n_{Grz} \rangle$, where $W^n_{Grz} = \{ i \mid 0 \leq i \leq n \}$, $R^n_{Grz} = R_{Grz} \cap (W^n_{Grz} \times W^n_{Grz})$ and $P^n_{Grz}(p) = P_{Grz}(p) \cap W^n_{Grz}$. Then

1. $M^n_{Grz}$ is an exact model for $S^n(p)$ in $Grz$.
2. $\langle S^n(p)/ \equiv_{Grz}, \leq_{Grz} \rangle$ is isomorphic to $\langle 2^{W^n_{Grz}}, \subseteq \rangle$.
3. for any $A \in S^n(p)$, $A \equiv_{Grz} \bigwedge \{ f(k) \mid (M^n_{Grz}, k) \not\equiv A \}$, where

$$f(k) = \begin{cases} \square F_{k+1} \supset \square F_k & \text{if } 0 \leq k \leq n - 2 \\ F_k & \text{if } n - 1 \leq k \leq n. \end{cases}$$

### 3 The structure $\langle S(\perp)/ \equiv_{GL}, \leq_{GL} \rangle$

In this section, we treat the structure $\langle S(\perp)/ \equiv_{GL}, \leq_{GL} \rangle$ as in the previous section. In [1], we can see many useful results for our study, and most of the result here can be given by considering carefully the correspondence between proofs in [1] and the structure. However, since [1] does not aim at the structure, there are notions that we do not need.
Also we would like to compare the structures here with the one in the previous section. So, we treat the structure in a similar way to the previous sections. Some lemmas below can be proved using results in [1] but we will give their proof directly in the case that we have to define new notions in order to use the result in [1] and the case that it seems to be better in order to compare the structure with the one in the previous section.

First, we construct representatives in the quotient set \( S^n(\bot)/\equiv_{GL} \).

**Definition 3.1.** The sets \( G_n^* \) (\( n = 0, 1, 2, \cdots \)) of formulas are defined as \( G_n^* = \{ \Box^n \bot, \Box^n \bot \supset \Box^{n-1} \bot, \cdots, \Box \bot \supset \bot \} \).

**Lemma 3.2.** For \( k \geq 0 \), \( \Box(\Box^{k+1} \bot \supset \Box^k \bot) \equiv_{Grz} \Box^{k+1} \bot \).

**Proof.** By Lemma 1.2(3) and the axiom \( L \).

**Theorem 3.3.**
(1) \( S^n(\bot)/\equiv_{GL} = \{ [\bigwedge S] \mid S \subseteq G_n^* \} \).
(2) For any subsets \( S_1 \) and \( S_2 \) of \( G_n^* \), \( S_1 \subseteq S_2 \) iff \( [\bigwedge S_2] \leq_{GL} [\bigwedge S_1] \).

Theorem 3.3 can be proved using the lemma in [1] below, but we have to check depth of formulas and independence of the elements in \( G_n^* \).

**Lemma 3.4([1]).** If \( A \) is a formula in \( S(\bot) \), then there exists a truth-functional combination \( B \) of formulas of the form \( \Box^k \bot \) such that \( A \equiv_{GL} B \).

Here we prove Theorem 3.3 in a similar way to section 2. Some lemmas we will show are useful to the investigation in section 4.

We define the Kripke model \( M_{GL} = (W_{GL}, R_{GL}, P_{GL}) \), where \( W_{GL} = N \), \( R_{GL} = \{(k, \ell) \mid k > \ell \} \) and \( P_{GL}(A) = \emptyset \) for any propositional variable \( A \).

**Lemma 3.5.** \((M_{Grz}, k) \models \Box^i \bot \iff i > k \).

**Proof.** By an induction \( i \).

[1] introduced two notions rank and trace, and using the result in [1], we can show that the rank of \( k \in W_{GL} \) is \( k \) and that the trace of a formula \( A \) is \( P_{GL}(A) \). As the result, the lemma above is just the lemma in [1](Lemma 5 in Chapter 7).

**Lemma 3.6.**
(1) None of the formulas in \( G_n^* \) is provable in GL.
(2) For \( k \leq n \), \( \wedge \{ \Box^n \bot, \Box^n \bot \supset \Box^{n-1} \bot, \cdots, \Box^{k+1} \bot \supset \Box^k \bot \} \equiv_{GL} \Box^k \bot \).
(3) For any \( A, B \in G_n^* \), \( A \neq B \) implies \( A \lor B \in GL \),
(4) For any \( A, B \in G_n^* \), \( A \neq B \) implies \( B \equiv_{GL} A \lor B \).

**Proof.** By Lemma 3.5 and Lemma 1.1(2), we have (1). By Lemma 1.2(2), we have (2) and (3). By (3), we have (4).
Lemma 3.7. Let $S_1$ and $S_2$ be subsets of $G^*_n$. Then

1. $(\bigvee S_1) \land (\bigvee S_2) \equiv GL \bigvee (S_1 \cup S_2),$
2. $(\bigvee S_1) \lor (\bigvee S_2) \equiv GL \bigvee (S_1 \cap S_2),$
3. $(\bigvee S_1) \supset (\bigvee S_2) \equiv GL \bigvee (S_2 - S_1),$
4. If $S_1 \neq \emptyset$, then $\Box (\bigvee S_1) \equiv GL \Box^k \bot$, where $k = \min \{n + 1 \mid \Box^n \bot \in S_1\} \cup \{i + 1 \mid \Box^{i+1} \bot \supset \Box^i \bot \in S_1\}$.

Proof. (1),(2) and (3) can be shown similarly to Lemma 2.7. (4) can also be shown, but we use Lemma 3.2 and Lemma 1.2(2) instead of Lemma 2.2(3) and Lemma 2.2(1). \(\square\)

Lemma 3.8. Let $A$ be a formula in $S^n(\bot)$. Then there exists a subset $S$ of $G^*_n$ such that $A \equiv_{GL} \bigwedge S$.

Proof. We use an induction on $A$. If $A = \bot$, then by Lemma 3.6(2),

$$\bigwedge G^*_n = \bigwedge \{\Box^n \bot, \Box^n \bot \supset \Box^{n-1} \bot, \ldots, \Box \bot \supset \bot\} \equiv_{GL} \bot = A.$$ If $A \neq \bot$, then by the induction hypothesis, Lemma 3.7 and Lemma 3.6(2), we obtain the lemma. \(\square\)

Similarly to Theorem 2.4, Theorem 3.3 is proved by Lemma 3.8, Lemma 3.6(1) and Lemma 3.6(4).

Next, we clarify the structure $\langle S(\bot)/ \equiv_{GL}, \leq_{GL}\rangle$ by giving a concrete general exact model for $S(\bot)$ in $GL$. We put

$$Q_{GL} = \{S_1 \cup S_2 \mid S_1 \in Q_{fin}(N), S_2 \in \{N_k \mid k \in N\} \cup \{\emptyset\}\}.$$ 

Lemma 3.9. $Q_{GL}$ is closed under $\cap$.

Proof. Let $S$ and $T$ be any sets in $Q_{GL}$. Then there exist $S_1, T_1 \in Q_{fin}(N)$ and $S_2, T_2 \in \{N_k \mid k \in N\} \cup \{\emptyset\}$ such that $S = S_1 \cup S_2$ and $T = T_1 \cup T_2$. Similarly to the proof of Lemma 2.11, $S \cap T = U \cup (S_2 \cap T_2)$, where $U = (S_1 \cap (T_1 \cup T_2)) \cup (S_2 \cap T_1) \in Q_{fin}(N)$. On the other hand, we have $S_2 \cap T_2 \in \{N_k \mid k \in N\} \cup \{\emptyset\}$ since $N_i \cap N_j = N_{\max(i,j)}$ and $N_i \cap \emptyset = \emptyset$. Hence $S \cap T \in Q_{GL}$. \(\square\)

Theorem 3.10. The structure $\langle W_{GL}, R_{GL}, P_{GL}, Q_{GL}\rangle$ is a general exact model for $S(\bot)$ in $GL$.

Proof. The conditions (1) and (2) in Definition 2.9 are clear. Also (4) in Definition 2.9 can be shown similarly to the proof of Theorem 2.13 using Lemma 3.6(1) and Lemma 3.8. So, it is sufficient to show

3. $\text{Im}(P_{GL}) = Q_{GL}$.

Suppose that $S \in \text{Im}(P_{Grz})$. Then there exists a formula $A$ such that $P_{GL}(A) = S$. Let be that $n = d(A)$. Then using Lemma 3.8, there exists a subset $S_1$ of $G^*_n$ such that $A \equiv_{GL} \bigwedge S_1$. So, similarly to Theorem 2.13, we have $P_{GL}(A) = \bigcap_{B \in S_1} P_{GL}(B)$. By Lemma 3.5, we have

5. $P_{GL}(\Box^{k+1} \bot \supset \Box^k \bot) = N - \{k\} = \{i \mid 0 \leq i < k\} \cup N_{k+1} \in Q_{GL}$,
\( P_{GL}(\square^k \perp) = \{ i | i < k \} \in Q_{GL} \).

So, for any formula \( B \in G_n \), we have \( P_{GL}(B) \in Q_{GL} \), and using Lemma 3.9, we obtain
\[ S = P_{Grz}(A) \in Q_{Grz}. \]

Suppose that \( S \in Q_{GL} \). Then there exist \( S_1 \in Q_{fin}(N) \) and \( S_2 \in \{ N_k | k \in N \} \cup \{ \emptyset \} \) such that \( S = S_1 \cup S_2 \). On the other hand, by Lemma 3.5, (5) and (6), we have
\[ (7) \quad P_{GL}(\perp) = \emptyset, \]
\[ (8) \quad P_{GL}(\bigcup_{k \in S_1} \{ k \}) \in Q_{Grz}, \]
\[ (9) \quad P_{GR}(\square^k \perp \supset \perp) = N_k \in Q_{Grz}. \]

Here we note that \( S_2 = P_{GL}(B) \) for some \( B \in S(\perp) \).

Also similarly to Theorem 2.14, for any formulas \( C \) and \( D \), \( P_{GL}(C) \cup P_{GL}(D) = P_{GL}(C \lor D) \).

Hence we obtain \( S \in Im(P_{GL}). \)

Considering the relations that the rank of \( k \in W_{GL} \) is \( k \) and that the trace of a formula \( A \) is \( P_{GL}(A) \), the conditions (4), (8) and \( Im(P_{GL}) \subseteq Q_{GL} \), in the proof above, have been shown in [1].

Corollary 3.11. \( \langle S(\perp)/\equiv_{GL}, \leq_{GL} \rangle \) is isomorphic to \( \langle Q_{GL}, \subseteq \rangle \).

Also by sketching the proof of Theorem 3.10, we also have

**Theorem 3.12.** Let be that \( M^0_{GL} = \langle W^0_{GL}, R^0_{GL}, P^0_{GL} \rangle \), where \( W^0_{GL} = \{ i | 0 \leq i \leq n \} \), \( R^0_{GL} = R_{GL} \cap (W^0_{GL} \times W^0_{GL}) \) and \( P^0_{GL}(q) = P_{GL}(q) \) for any \( q \). Then
(1) \( M^0_{GL} \) is an exact model for \( S_n(\perp) \) in \( GL \),
(2) \( \langle S^n(\perp)/\equiv_{GL}, \leq_{GL} \rangle \) is isomorphic to \( \langle 2^{W^0_{GL}}, \subseteq \rangle \),
(3) for any \( A \in S^n(\perp) \), \( A \equiv_{GL} \{ f(k) | (M^0_{GL}, k) \not\in A \} \), where
\[ f(k) = \begin{cases} 
\square^{k+1} \supset \square^k \perp & \text{if } 0 \leq k \leq n - 1 \\
\square^k \perp & \text{if } k = n.
\end{cases} \]

### 4 GL-provability of formulas in \( S(\perp) \) in \( Grz \)

Comparing \( \langle S(p)/\equiv_{Grz}, \leq_{Grz} \rangle \) with \( \langle S(\perp)/\equiv_{GL}, \leq_{GL} \rangle \), we can see some kinds of correspondence between them. We show one kind of correspondence by giving a way to express the GL-provability of formulas in \( S(\perp) \) in \( Grz \). More precisely, we give a mapping \( g \) from \( S(\perp) \) to \( S(p) \) satisfying for any formula \( A \in S(\perp) \),
\[ A \in GL \text{ iff } g(A) \in Grz. \]
On the other hand, Goldblatt [5] gave a mapping $*$, satisfying for any formula $A$,

$$A^* \in \text{GL} \iff A \in \text{Grz}.$$ 

From this, we can know the Grz-provability of any formulas in $\text{GL}$. Our result is a kind of the converse of his result with restriction.

**Definition 4.1.** A list $g_0, g_1, \ldots$ of mappings from $S(\bot)$ to $S(p)$ are defined inductively as follows:

1. $g_i(\bot) = \Box F_i$,
2. $g_i(B \otimes C) = g_i(B) \otimes g_i(C)$, for $\otimes \in \{\land, \lor, \supset\}$,
3. $g_i(\Box B) = \Box g_{i+1}(B)$.

Also we put $g_i(S) = \{g_i(A) \mid A \in S\}$. The mapping $g_0$ transforms the formula $\Box(\Box \Box \bot \supset \bot) \supset \Box \bot$, an instance of the axiom $L$, into a formula

$$g_0(\Box(\Box \Box \bot \supset \bot) \supset \Box \bot) = \Box(\Box \Box (F_1 \supset \Box F_1) \supset \Box F_1).$$

Here we note that the image is similar to $\Box(\Box (F_1 \supset \Box F_1) \supset F_1)$, an instance of the axiom $grz$, and that the image and instance are equivalent in $\text{Grz}$. The mappings $g_i$ also have the same property since $g_i(\Box(\Box \Box \bot \supset \bot) \supset \Box \bot) = \Box(\Box \Box (F_{i+1} \supset \Box F_{i+1}) \supset \Box F_{i+1}).$
have either $[\bigwedge h_i(S)] \not\leq_{\text{Grz}} [\bigwedge h_i(S)]$ or $[\bigwedge h_i(S)] \not\leq_{\text{Grz}} [\bigwedge h_i(S)]$. So, either $(\bigwedge h_i(S)) \supset (\bigwedge h_i(S_i)) \not\subseteq \text{Grz}$ or $(\bigwedge h_i(S_i)) \supset (\bigwedge h_i(S)) \not\subseteq \text{Grz}$. Hence $\bigwedge h_i(S) \not\subseteq \text{Grz} \bigwedge h_i(S_i)$.

**Lemma 4.6.** Let $S_1$ and $S_2$ be subsets of $G_\infty^*$. Then for any $i$,

1. $(\bigwedge h_i(S_1)) \land (\bigwedge h_i(S_2)) \equiv_{\text{Grz}} h_i(S_1 \cup S_2),$
2. $(\bigwedge h_i(S_1)) \lor (\bigwedge h_i(S_2)) \equiv_{\text{Grz}} h_i(S_1 \cap S_2),$
3. $(\bigwedge h_i(S_1)) \supset (\bigwedge h_i(S_2)) \equiv_{\text{Grz}} h_i(S_2 - S_1),$
4. if $S_1 \neq \emptyset$, then $\Box(\bigwedge h_i(S_1)) \equiv_{\text{Grz}} \Box F_k$, where $k = \min\{n+i \mid \Box^n \perp \in S_1\} \cup \{j+i \mid \Box^j \perp \in S_1\}$.

**Proof.** We note that $h_i(S_1) \cup h_i(S_2) = h_i(S_1 \cup S_2)$, $h_i(S_1) \cap h_i(S_2) = h_i(S_1 \cap S_2)$ and $h_i(S_2) - h_i(S_1) = h_i(S_2 - S_1)$. So, using Lemma 2.7 and Lemma 4.5(1) we obtain (1), (2) and (3).

We show (4). By Lemma 2.7(4) and Lemma 4.5(1), $\Box(\bigwedge h_i(S_1)) \equiv_{\text{Grz}} \Box F_k$, where $k = \min\{n+i \mid \Box^n \perp \in S_1\} \cup \{j+i \mid \Box^j \perp \in S_1\}$.

**Lemma 4.7.** Let $A$ be a formula in $S^n(\perp)$ and let $S$ be a subset of $G_\infty^*$. Then for any $i$,

$$A \equiv_{\text{GL}} \bigwedge S \iff g_i(A) \equiv_{\text{Grz}} \bigwedge g_i(S).$$

**Proof.** We use an induction on $A$.

**Basis** ($A = \perp$): By Lemma 3.6(2), we have $A = \perp \equiv_{\text{GL}} \bigwedge G_\infty^*$. By Lemma 4.5(2), Lemma 4.6(4),

$$\bigwedge g_i(G_\infty^*) \equiv_{\text{Grz}} \bigwedge h_i(G_\infty^*) \equiv_{\text{Grz}} \Box F_i = g_i(\perp) = g_i(A).$$

So, if $S = G_\infty^*$, then we have both of $A \equiv_{\text{GL}} \bigwedge S$ and $g_i(A) \equiv_{\text{Grz}} \bigwedge g_i(S)$. If not, then by Theorem 3.3(2), $A \not\equiv_{\text{GL}} \bigwedge S$, and by Lemma 4.5(2) and Lemma 4.5(3),

$$\bigwedge g_i(S) \equiv_{\text{Grz}} \bigwedge h_i(S) \not\equiv_{\text{Grz}} \bigwedge h_i(G_\infty^*) \equiv_{\text{Grz}} \bigwedge g_i(G_\infty^*) \equiv_{\text{Grz}} g_i(A),$$

and so, $g_i(A) \not\equiv_{\text{Grz}} \bigwedge g_i(S)$.

**Induction step** ($A \neq \perp$): We divide the cases.

The case that $A = A_1 \land A_2$: We note $A_1, A_2 \in S^n(\perp)$. So, by Lemma 3.8, there exist subsets $S_1, S_2$ of $G_\infty^*$ such that $A_1 \equiv_{\text{GL}} \bigwedge S_1$ and $A_2 \equiv_{\text{GL}} \bigwedge S_2$. Using the induction hypothesis, $g_i(A_1) \equiv_{\text{Grz}} \bigwedge g_i(S_1)$ and $g_i(A_2) \equiv_{\text{Grz}} \bigwedge g_i(S_2)$. By Lemma 3.7(1),

$$A = A_1 \land A_2 \equiv_{\text{GL}} (\bigwedge S_1) \land (\bigwedge S_2) \equiv_{\text{GL}} (\bigwedge S_1 \cup S_2).$$

Also by Lemma 4.5(2) and Lemma 4.6(1),

$$g_i(A) = g_i(A_1) \land g_i(A_2) \equiv_{\text{Grz}} (\bigwedge g_i(S_1)) \land (\bigwedge g_i(S_2)) \equiv_{\text{Grz}} (\bigwedge h_i(S_1)) \land (\bigwedge h_i(S_2))$$
\[ \equiv_{\text{Grz}} \bigwedge h_i(S_1 \cup S_2) \equiv_{\text{Grz}} \bigwedge g_i(S_1 \cup S_2). \]

So, considering the cases that \( S = S_1 \cup S_2 \) and that \( S \neq S_1 \cup S_2 \), we obtain the lemma similarly to Basis. The case that \( A = A_1 \lor A_2 \) can be shown similarly, but we use Lemma 3.7(2) and Lemma 4.6(2) instead of Lemma 3.7(1) and Lemma 4.6(1). Also the case that \( A = A_1 \supset A_2 \) we use Lemma 3.7(3) and Lemma 4.6(3) instead of them.

The case that \( A = \square A_1 \): Similarly to the above cases, there exists a subset \( S_1 \) of \( G_{n-1}^* \) such that for any \( k, A_1 \equiv_{\text{GL}} \bigwedge S_1 \) and \( g_k(A_1) \equiv_{\text{Grz}} \bigwedge g_k(S_1) \).

If \( S_1 = \emptyset \), then we have
\[
A = \square A_1 \equiv_{\text{GL}} \square \bigwedge \emptyset \equiv_{\text{GL}} \square (\perp \supset \perp) \equiv_{\text{GL}} \perp \supset \perp \equiv_{\text{Grz}} \perp \supset \perp \equiv_{\text{GRz}} \bigwedge \emptyset.
\]
and
\[
g_i(A) = \square g_{i+1}(A_1) \equiv_{\text{Grz}} \square \bigwedge g_{i+1}(\emptyset) = \square \bigwedge g_i(\emptyset) \equiv_{\text{Grz}} \square \bigwedge g_i(0).
\]
So, considering the cases that \( S = \emptyset \) and that \( S \neq \emptyset \), we obtain the lemma similarly to Basis.

If \( S_1 = \{\square^n \perp\} \), then we have
\[
A = \square A_1 \equiv_{\text{GL}} \square \bigwedge \{\square^n \perp\} \equiv_{\text{GL}} \bigwedge \{\square^n \perp\}
\]
and
\[
g_i(A) = g_i(\square A_1) = \square g_{i+1}(A_1) \equiv_{\text{Grz}} \square \bigwedge g_{i+1}(\{\square^n \perp\})
\]
\[
\equiv_{\text{Grz}} \square g_{i+1}(\bigwedge \{\square^n \perp\}) \equiv_{\text{Grz}} \bigwedge g_i(\{\square^n \perp\}) \equiv_{\text{Grz}} \square \bigwedge g_i(\bigwedge \{\square^n \perp\}).
\]
So, considering the cases that \( S = \{\square^n \perp\} \) and that \( S \neq \{\square^n \perp\} \), we obtain the lemma similarly to Basis.

Suppose that \( S_1 \not\in \{\emptyset, \{\square^n \perp\}\} = P(\{\square^n \perp\}) \). Then we have \( S_1 \not\in \{\square^n \perp\}\). Since \( S_1 \subseteq G_{n-1}^* \), we have \( \emptyset \neq S_1 - \{\square^n \perp\} \subseteq \{\square^n \perp \supset \square^n \perp, \ldots, \square \perp \supset \perp\} \). So, there exists the minimum \( k \) of \( \{\ell | \square^\ell \perp \supset \square^{\ell-1} \in S_1\} \). By Lemma 3.7(4) and Lemma 3.6(2),
\[
A = \square A_1 \equiv_{\text{GL}} \square \bigwedge S_1 \equiv_{\text{GL}} \square^{k}\perp
\]
\[
\equiv_{\text{GL}} \bigwedge \{\square^n \perp, \square^n \perp \supset \square^{n-1} \perp, \ldots, \square^{k+1} \perp \supset \square^k \perp\}.
\]
Also by Lemma 4.5(2) and Lemma 4.6(4),
\[
g_i(A) = g_i(\square A_1) = \square g_{i+1}(A_1) \equiv_{\text{Grz}} \square \bigwedge g_{i+1}(S_1) \equiv_{\text{Grz}} \square \bigwedge h_{i+1}(S_1) \equiv_{\text{Grz}} \square F_{k+i}
\]
\[
\equiv_{\text{Grz}} \bigwedge \{F_{n+i+1}, F_{n+i}, \square F_{n+i+1} \supset \square F_{n+i-1}, \ldots, \square F_{k+i+1} \supset \square F_{k+i}\}
\]
\[
\equiv_{\text{Grz}} \bigwedge h_i(\{\square^n \perp, \square^n \perp \supset \square^{n-1} \perp, \ldots, \square^{k+1} \perp \supset \square^k \perp\})
\]
\[
\equiv_{\text{Grz}} \bigwedge g_i(\{\square^n \perp, \square^n \perp \supset \square^{n-1} \perp, \ldots, \square^{k+1} \perp \supset \square^k \perp\}).
\]
So, considering the case that $S = \{\Box^n \bot, \Box^n \bot \supset \Box^{n-1} \bot, \ldots, \Box^{k+1} \bot \supset \Box^k \bot\}$ and the other case, we obtain the lemma similarly to Basis.

Considering the case that $S = \emptyset$ in Lemma 4.7, we obtain Theorem 4.3.

References


