On Confluent PCE grammars

Kunio Aizawa¹

and

Akira Nakamura²

¹Department of Applied Mathematics, Hiroshima University, Higashi-Hiroshima, 724 Japan
²Department of Information Science, Meiji University, Kawasaki, Kanagawa, 214 Japan

Abstract: In this paper, a path controlled embedding graph grammar (PCE graph grammar) having the confluent property is proposed. Then the relationships between confluent PCE grammars and array languages are investigated.

1. Introduction

In the recent years, many models for context-free graph grammars have been proposed (see, e.g., Ehrig, et al. [2]). Some grammars are node rewriting and others are edge (or hyper-edge) rewriting. These grammars are context-free in the sense that one node is replaced without considering any other part of the rewrited graph. However, they may still be context-sensitive in the sense that generated graph depends on the order in which the production rules are applied. A graph grammar that does not suffer from this context-sensitivity is said to be confluent (see, e.g., Engelfriet [3]).

In general, graph grammars are less powerful to describe structures in contrast of their generative power. In Aizawa and Nakamura [1], we introduced a graph grammar called node-replacement graph grammar with path controlled embedding (nPCE grammars) which use a sequence of edges instead of the single edge to embedding a newly replaced graph into the host graph. It has been shown that there exists a subclass of nPCE grammars generating context-free array languages. In this paper, a path controlled embedding graph grammar (PCE graph grammar) having the confluent property is proposed. Then the relationships between confluent PCE grammars and array languages are investigated.

We assume for readers to be familiar with the theories of two-dimensional grammars and graph grammars (see e.g., Nagl [5] and Rosenfeld [6]). The following notations will be used in the rest of this paper.

(1) Let X be set. By \(2^X\) we denote the set of subsets of X and if X is finite, then \#X denotes
the cardinality of X.

(2) Let \( \pi=\{c_1c_2...c_i\} \) be a string. By \( \pi^R \) we denote the reverse string of \( \pi \), i.e., 
\[
\pi^R=\{c_i...c_2c_1\}.
\]
and \( \pi \) denotes the length of \( \pi \), i.e., \( i \).

(3) A graph is a system \( H=(V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E) \), where \( V \) is a finite nonempty set of natural numbers called the set of nodes, \( E \) is a set of pairs of two elements from \( V \) called the set of edges, \( \Sigma_V \) is a finite nonempty set called the set of node labels, \( \Sigma_E \) is a finite nonempty set called the set of edge labels, \( \phi_V \) is a mapping from \( V \) into \( \Sigma_V \) called the nodes labelling function, and \( \phi_E \) is a mapping from \( E \) into \( \Sigma_E \) called the edges labelling function. \( H \) is called a graph over \( (\Sigma_V, \Sigma_E) \). Throughout this paper, \( V(H) \) and \( E(H) \) denote the set of nodes and the set of edges of \( H \), respectively.

(4) Let \( H=(V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E) \) be a graph and let \((x,y)\) is an edge of \( H \). We say that the edge \((x,y)\) is incident with the nodes \( x \) and \( y \), and the nodes \( x, y \) are neighbors.

(5) Let \( H=(V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E) \) be a graph and let \( x \) be a node of \( H \). Then the degree of \( x \), denoted as \( \deg(x) \), is the number of edges incident with \( x \).

(6) Let \( A=(V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E) \) and \( B=(V', E', \Sigma_{V'}, \Sigma_{E'}, \phi_{V'}, \phi_{E'}) \) be graphs. \( A \) is a subgraph of \( B \) if \( V'\supseteq V \), \( E'\cap\{(x, y)\mid x, y\in V\}\supseteq E \), \( \Sigma_{V'}\supseteq \Sigma_{V} \), \( \Sigma_{E'}\supseteq \Sigma_{E} \), \( \phi_{V'}(x)=\phi_{V}(x) \) for \( \forall x\in V \), and \( \phi_{E'}((x, y))=\phi_{E}((x, y)) \) for \( \forall (x, y)\in E \). In this case, we call \( A \) the subgraph spanned by \( V \) in \( B \). By \( B-A \) we denote the subgraph spanned by \( V\setminus V \) in \( B \).

(7) Let \( A=(V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E) \) and \( B=(V', E', \Sigma_{V'}, \Sigma_{E'}, \phi_{V'}, \phi_{E'}) \) over \( (\Sigma_V, \Sigma_E) \). An isomorphism from \( A \) into \( B \) is a bijective mapping \( h \) from \( V \) into \( V' \) such that \( \phi_{V'} h=\phi_{V} \) and \( E'=\{(h(x),h(y))\mid (x, y)\in E\} \). We say that \( A \) is isomorphic to \( B \).

(8) A graph \( A=(V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E) \) is connected if for every \( x, y \) in \( V \), there exists a sequence \( x_1, x_2, ..., x_n \) of nodes in \( V \) such that \( x_1=x, x_n=y \) and for \( 1\leq i\leq n-1, x_i \) is a neighbor of \( x_{i+1} \).

2. Path controlled embedding

In this section, we define a kind of graph grammars called U-nPCE graph grammars and languages. They are restricted version of nPCE graph grammars proposed in Aizawa and Nakamura [1].

**Definition 2.1.** For any given graph \( H \), its node \( P \), and a string \( \pi=\{c_1c_2...c_i\} \) of its edge labels, \( P\pi \) is realizable on \( H \) if and only if there exists a set of nodes \( \{P_0, P_1, ..., P_i\} \) such that \( P_0=P \) and \( P_j \) is a neighbor of \( P_{j-1} \) joined by an edge labelled with \( c_j \) (\( 1\leq j\leq i \)).

**Definition 2.2.** A node-replacement graph grammar with path controlled embedding,
denoted as nPCE grammar, is a construction
\[ G = \langle \Sigma_N, \Sigma_E, P, Z, \Delta_N, \Delta_E \rangle, \]
where
- \( \Sigma_N \) is a finite nonempty set of node labels,
- \( \Sigma_E \) is a finite nonempty set of edge labels,
- \( \Delta_N \) is a finite nonempty subset of \( \Sigma_N \), called terminal node labels,
- \( \Delta_E \) is a finite nonempty subset of \( \Sigma_E \), called terminal edge labels,
- \( P \) is a finite set of productions of form \( (a, \beta, \psi) \), where \( a \) is a node, \( \beta = (V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E) \) is a connected graph, \( \psi \) is a mapping from \( \Sigma_E^+ \) into \( V(\beta) \times \Sigma_E \). \( \psi \) is called embedding function.
- \( Z \) is a connected graph over \( (\Sigma_N, \Sigma_E) \) called the axiom.

A direct derivation step in a nPCE grammar is performed as follows:
Let \( H = (V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E) \) be a graph. Let \( p = (a, \beta) \) be a production in \( P \) and \( \psi \) be the embedding function. Let \( \beta' \) be isomorphic to \( \beta \) (with \( h \) an isomorphism from \( \beta' \) into \( \beta \)), where \( \beta' \) and \( H-a \) have no common nodes. Then the result of the application of \( p \) to \( H \) (by using \( h \)) is obtained by first removing \( a \) from \( H \), then replacing \( a \) with \( \beta' \) and finally adding edges \((u,v)\) between every nodes \( u \) in \( \beta' \) and every \( v \) in \( H-a \) such that there exists a path \( \pi \in \text{domain(}\psi\text{)} \) with \( v = a\pi \) in \( H \). The production \( p \) is applicable to the graph \( H \) if following conditions hold:

1. If \( a \) has at least one neighbor in \( H \), there exist at least one realizable path \( \pi \), and
2. if a node \( v \) of \( H-a \) is adjacent to the node labelled with \( a \) in \( H \), \( v \) must be adjacent to the node 1 of \( \beta' \).

Note here that the embedding function \( \psi \) bring no significant context into the generation procedure of an nPCE grammar since no node labels are referred in any place of embedding steps except the one of the newly replaced graph.

Formally the notion of a direct derivation step is defined as follows:

**Definition 2.3.** Let \( G = \langle \Sigma_N, \Sigma_E, P, Z, \Delta_N, \Delta_E \rangle \) be a nPCE grammar and let \( H, H' \) be graphs over \( (\Sigma_V, \Sigma_E) \).

1. \( H \) directly derives \( H' \) in \( G \), denoted as \( H \Rightarrow_G H' \), if there exists a production \( p = (a, \beta, \psi) \) in \( P \), a graph \( \beta' \) with \( V(\beta') \cap V(\beta-a) = \emptyset \) and an isomorphism \( h \) from \( \beta' \) into \( \beta \) such that \( H' \) is isomorphic to the graph \( X \) constructed as follows:

\[
X = (V, E, \Sigma_V, \Sigma_E, \phi_V, \phi_E),
\]
where
- \( V = V(H-a) \cup V(\beta') \),
- \( E = \{ (x,y) \mid x,y \in V(H-a) \text{ and } (x,y) \in E(H) \} \cup \{ (x,y) \mid x,y \in V(\beta') \text{ and } (h(x), h(y)) \in E(\beta) \} \)
∪ \{ (x,y) \mid x \in V(\beta'), y \in V(H-a), \text{ there exists a path } \pi \in \text{domain}(\psi) \\
\text{with } y = a\pi \text{ and } <h(x), e> \in \psi(\pi) \text{ for some } e \in \Sigma_E \} ,

\phi_{\psi} \text{ is equal to the node labelling function of } H \text{ for nodes in } V(H-a), \text{ equal to the node labelling function of } \beta' \text{ for nodes in } V(\beta'),

\phi_{\beta} \text{ is equal to the edge labelling function of } H \text{ for edges between the nodes of } H, \text{ is equal to the edge labelling function of } \beta' \text{ for nodes in } V(\beta'),

We also say that \( H' \) is derived from \( H \) by replacing \( a \) using the production \( p \).

(2) We will denote the reflexive and the transitive closure of \( \Rightarrow_G \) by \( \Rightarrow_G^* \) and the transitive closure of \( \Rightarrow_G \) by \( \Rightarrow_G^+ \).

(3) The language of \( G \), denoted as \( L(G) \), is defined by \( L(G) = \{ H \mid H \text{ is a graph over } (\Sigma_V, \Sigma_E) \text{ and } Z \Rightarrow_G^* H \} \).

We here present an example of the derivations of \( \text{nPCE} \) grammars.

**Example 2.1.** The application of the production in Fig. 1a to the graph in Fig. 1b results the graph in Fig. 1c.

By making use of the path controlled embedding mechanism defined above, it is possible to construct \( \text{nPCE} \) grammars generating array languages under some way to regard a graph as a two-dimensional array. In Aizawa and Nakamura [1], it is shown that, for any given context-free array grammar \( G \), there exists an \( \text{nPCE} \) grammar \( G' \) such that \( L(G') \) is regarded as a set of two-dimensional array, \( L(G) \). However, describing the \( \text{nPCE} \) grammars generating array languages needs somewhat clumsy production rules especially in path descriptions. So, \( \text{nPCE} \) grammars using partial path group to describe paths in the embedding functions seems to be more suitable for describing patterns having some geometrical structures. We review here the definitions of the partial path groups. For more precise definitions of the partial path groups, see Rosenfeld [7].

**Definition 2.4.** Let \( H \) be a connected graph of degree \( d \), i.e., no more than \( d \) edges emanate from any node. By an edge coloring of \( H \) we mean an assignment of colors to the edges of \( H \) such that the edges emanating from any given node all have different colors.

**Definition 2.5.** Let \( p \) be a node of a graph \( H \) and let \( \pi \) be a string of colors. Then \( p\pi \) is defined as the terminal node of the path defined by \( \pi \) starting from \( p \), provided this path is realizable. For convenience, let us define a fictitious "blank" color representing "no move". It is obvious that the structure defined above resembles a group structure called "partial group". From now on we shall refer to this partial group as the partial path group of \( H \) defined by given coloring, and denote it by \( \Pi(H) \).
Note that the partial path groups can be defined on a graph generated by a graph grammar by regarding the edge labels of the generated graph to the colors mentioned in above definitions.

**Definition 2.6.** Let $\alpha = c_{i_1} \ldots c_{i_k}$ be a string of colors. The string $\alpha'$ is called an elementary reduction of $\alpha$ if one of the following statements is true:

a) $k>1$; for some $1 \leq j \leq k$ we have $c_{i_{j}} = c_{0}$; and $\alpha' = c_{i_1} \ldots c_{i_{j-1}} c_{i_{j+1}} \ldots c_{i_k}$.

b) $k>2$; for some $1 \leq j < k$ we have $c_{i_{j}} = c_{i_{j+1}}$; and $\alpha' = c_{i_1} \ldots c_{i_{j-1}} c_{i_{j+2}} \ldots c_{i_k}$.

c) $k=2$; $c_{i_1} = c_{i_2}$; and $\alpha' = c_{0}$.

The string $\alpha'$ is called a reduction of $\alpha$ if there exists strings $\alpha = \alpha_0$, $\alpha_1$, ..., $\alpha_m = \alpha'$ such that $\alpha_i$ is an elementary reduction of $\alpha_{i-1}$, $1 \leq i \leq m$. $\alpha$ is called fully reducible if $\alpha = c_{0}$, or if $c_{0}$ is a reduction of $\alpha$. $\Pi(G)$ is called free if $P\alpha = P$ implies that $\alpha$ is fully reducible.

**Definition 2.7.** A partial path group $\Pi(H)$ is called near-abelian if each color except blank color commutes with all but one of the other colors, and does not commute with the remaining one.

It is shown in Rosenfeld [7] that the free near-abelian partial path groups with four colors correspond to a subgraph of a two-dimensional array.

We define nPCE grammars using embedding functions with near-abelian partial path groups to describe paths.

**Definition 2.8.** A nPCE grammar with 4 colors free near-abelian partial path groups, denoted as $\text{nPCE}_{\Pi 4}$ grammar, is a construction $G = <\Sigma_N, \Sigma_E, P, Z, \Delta_N, \Delta_E>$, where $\Sigma_N$, $\Sigma_E$, $Z$, $\Delta_N$, $\Delta_E$ are same as in the definition of nPCE grammars provided that $\Sigma_E = \Delta_E$ and $|\Delta_E| = 4$. $P$ is a finite set of production rules of form $(a, \beta, \psi_{\Pi 4})$, where $\psi_{\Pi 4}$ is a mapping from $\Sigma_E^{+}$ into $(V(\beta) \times \Sigma_E)$ provided that if $\psi_{\Pi 4}$ maps $\pi$ into $(i, c)$ for some $c \in \Delta_E$, then $c$ is the reduction of $\sigma \pi$, where $\sigma$ is the path between node 1 to $i$ in $\beta$; $\psi_{\Pi 4}$ is called embedding function with 4 colors near-abelian partial path groups. For a path $\pi$, $\psi_{\Pi 4}(\pi)$ implies the shortest realizable path equivalent to $\pi$ under near-abelian partial path groups instead of $\pi$ itself.

In the $\text{nPCE}_{\Pi 4}$ grammars, we use $\psi_{\Pi 4}$ instead of the embedding function $\psi$ of nPCE grammars. In this case, arbitrarily long paths can be used to embed the newly replaced graphs. Since the derivation steps of our $\text{nPCE}_{\Pi 4}$ grammars are proceeded in the context-free node-replacement style, there are no insurance that the degree of each node is always at most 4 and
the edges emanating from any given node have always different colors. In such cases, even if 
α is fully reducible, Pα may not be P. However, the embedding mechanism of nPCE 
grahamars still works correctly for such cases as far as using the shortest realizable paths.

One of the advantages to introduce the partial path groups into the path controlled 
embedding mechanism is flexibility of describing various structures. As mentioned above, 
two-dimensional square arrays correspond to free near-abelian groups with 4 colors, and other 
types of arrays (triangular, hexagonal) correspond to other types of abelian groups. The partial 
path groups allow us to treat other classes of graphs (such as trees and hypercubes etc.) in 
addition to arrays, and they have very simple characterizations. For detail discussions of those 
structures, see Rosenfeld [7]. The partial path groups for more complicated structures are 
discussed in Melter [4].

**Definition 2.9.** An nPCE$_{4}$ grammar is called having uniform embedding function, denoted 
as U-nPCE$_{4}$ grammar, if the following conditions for embedding functions hold.

1. For each pair of production rules $p_i = (a_i, \beta_i, \psi_{4_i})$, $p_j = (a_j, \beta_j, \psi_{4_j})$ and a path $\pi \in \Sigma_{E}^{+}$, 
   if both of $\psi_{4_i}(\pi)$ and $\psi_{4_j}(\pi)$ are defined, then $\psi_{4_i}(\pi) = \psi_{4_j}(\pi)$.

2. For each production rule $p = (a, \beta, \psi_{4})$, if $\psi_{4}(\pi) = (v, c)$ is defined for some $\pi \in \Sigma_{E}^{+}$ and 
v = 1, then $\psi_{4}(\pi^{R})$ is also defined and $\psi_{4}(\pi) = \psi_{4}(\pi^{R})$.

3. For each production rule $p = (a, \beta, \psi_{4})$, if $\psi_{4}(\pi) = (v, c)$ and $v \neq 1$, then there exist $\rho$ and 
   $\sigma$ such that $\pi = \rho \sigma R$, $\psi_{4}(\rho) = \psi_{4}(\rho^{R}) = (1, c)$ and $\psi_{4}(\rho^{R} \sigma) = \psi_{4}(\pi)$ are defined, and $\sigma$ 
is a shortest realizable path between node 1 and $v$ in $\beta$.

So we can describe a U-nPCE$_{4}$ grammar $G$ as a construction $G = \langle \Sigma_{N}, \Sigma_{E}, P, \psi_{4}, Z, \Delta_{N}, \Delta_{E} \rangle$, where $\psi_{4} = \bigcup_{p \in P} \psi_{4_i}$ and each production rule has no exclusive embedding 
function.

**3. Confluent derivations on PCE embedding**

As defined in the last section, nPCE grammars are context-free in the sense that one node is 
replaced without considering any other part of the rewrited graph. However, the grammar may 
still be context-sensitive in the sense that generated graph depends on the order in which the 
production rules are applied. A graph grammar is said to be confluent, if derivation steps on 
distinct vertices can be done in any order.

**Definition 3.1.** Let $G$ be an nPCE grammar. $G$ is confluent, denoted as C-nPCE grammar, 
if the following condition holds for every sentential form $H$ of $G$: Let $v_1$ and $v_2$ be distinct
nodes of H labelled with nonterminal labels, and let \( p_1 \) and \( p_2 \) be production rules applicable to \( v_1 \) and \( v_2 \), respectively. If \( H_{12} \) is derived from H by applying \( p_1 \) at first then \( p_2 \) and \( H_{21} \) is derived by applying \( p_2 \) at first then \( p_1 \), then \( H_{12}=H_{21} \).

The class of all graph languages generated by U-nPCE grammars are included in the class of all graph languages generated by C-nPCE grammars.

**Theorem 3.1.** For any given nPCE\( \Pi 4 \) grammar G, G is confluent if G has uniform embedding function.

**Proof:** Let \( G=<\Sigma_N, \Sigma_E, P, \psi_{\Pi 4}, Z, \Delta_N, \Delta_E> \) be U-nPCE\( \Pi 4 \) grammar. For any given sentential form H of G, let \( v_1 \) and \( v_2 \) be distinct nodes of H labelled with nonterminal labels, and let \( p_1=(a_{v_1}, \beta_1) \) and \( p_2=(a_{v_2}, \beta_2) \) be production rules applicable to \( v_1 \) and \( v_2 \), respectively. If \( v_1 \) and \( v_2 \) are not neighbors of H and any pair of nodes from \( \beta_1 \) and \( \beta_2 \) do not become neighbors by applying both rules, then obviously \( H_{12}=H_{21} \). If \( v_1 \) and \( v_2 \) are neighbors of H and the node 1s of \( \beta_1 \) and \( \beta_2 \) are not neighbors, then at least one of the production rules is not applicable. This is a contradiction. Assume here that \( v_1 \) and \( v_2 \) are not neighbors in H and become neighbors after applying \( p_1 \) and \( p_2 \). There exist four possible cases.

**Case 1:** Node 1s of \( \beta_1 \) and \( \beta_2 \) become neighbors.

From the definition of uniform embedding, there exists a path \( \pi \) from \( v_1 \) to \( v_2 \) and \( \psi_{\Pi 4}(\pi)=\psi_{\Pi 4}(\pi^R)=(1, c) \). Thus obviously \( H_{12}=H_{21} \).

**Case 2:** Node \( i \neq 1 \) of \( \beta_1 \) and node 1 of \( \beta_2 \) become neighbors.

From the definition of uniform embedding, there exists a path \( \pi \) from \( v_1 \) to \( v_2 \) such that \( \psi_{\Pi 4}(\pi)=(i, c) \) and \( \pi=\rho \sigma \) where \( \sigma \) is the path between node 1 to \( i \) in \( \beta_1 \) and c is the reduction of \( \rho \). \( \psi_{\Pi 4}(\rho^R)=(1, c) \) is also defined. Since c is an element of \( \Delta_E \), c is also the reduction of \( \rho^R \). Then \( H_{12}=H_{21} \).

**Case 3:** Node 1 of \( \beta_1 \) and node \( i \neq 1 \) of \( \beta_2 \) become neighbors.

Same as in the Case 2.

**Case 4:** Node \( i \neq 1 \) of \( \beta_1 \) and node \( j \neq 1 \) of \( \beta_2 \) become neighbors.

In this case, the path \( \sigma_i \) from node 1 to \( i \) in \( \beta_1 \) is equal to the path \( \sigma_j \) from 1 to \( j \) in \( \beta_2 \) under \( \Pi 4 \) since c of \( \psi_{\Pi 4}(\pi \sigma_j)=(i, c) \) is the reduction of \( \pi \), and also the reduction of \( \sigma_i^R \pi \sigma_j \). Then from the fact \( \psi_{\Pi 4}(\pi \sigma_j)=\psi_{\Pi 4}(\pi^R \sigma_j^R) \), \( H_{12}=H_{21} \).

The case in which \( v_1 \) and \( v_2 \) are neighbors in H and are also neighbors after applying \( p_1 \) and \( p_2 \) is proved in the almost same way as in the above case.
4. Array languages defined by U-nPCE grammars

In this section, we introduce two mappings which map graph languages generated by U-nPCE\(_{\Pi 4}\) grammar into the set of array languages. Then, we investigate the generative powers of U-nPCE\(_{\Pi 4}\) grammars as the array patterns generators.

**Definition 4.1.** A mapping \(k\) from a graph generated by U-nPCE\(_{\Pi 4}\) grammars into two-dimensional arrays is such that

1. The horizontal neighborhood on array is defined as the set of edges labelled with two colors, say \(h\) and \(h'\), which do not commute each other.
2. The vertical neighborhood on array is defined as the set of edges labelled with remaining two colors, say \(v\) and \(v'\), which also do not commute each other.
3. The label of each node is the symbol in the corresponding position of array.
4. For a graph \(H\) which has more than one node mapped to a position, \(k(H)\) is undefined.

The mapping \(k\) can be extended to the set of graphs in two different ways.

**Definition 4.2.** For any given graph language \(L(G)\) generated by a U-nPCE\(_{\Pi 4}\) grammar,

1. \(K(L(G)) = \begin{cases} \{k(g) \mid g \in L(G)\} & \text{if } k(g) \text{ is defined for all elements of } L(G) \\ \text{undefined} & \text{otherwise} \end{cases}\)
2. \(K'(L(G)) = \{k(g) \mid g \in L(G) \text{ and } k(g) \text{ is defined}\}\).

Both \(K\) and \(K'\) can be extended to the families of languages in the natural way, i.e., \(K(X) = \{K(L(G)) \mid G \text{ is a grammar in } X\}\) and \(K' = \{K'(L(G)) \mid G \text{ is a grammar in } X\}\).

From the definitions of nPCE\(_{\Pi 4}\) grammars and uniform embedding, it is not so difficult to see the following lemma holds:

**Lemma 4.1.** \(\mathcal{F}(K'(U-nPCE_{\Pi 4})) = \mathcal{F}(CFAG)\).

The same result is obtained in Aizawa and Nakamura [1] but the concept of uniform embedding is not used in it. If all production rules of a U-nPCE\(_{\Pi 4}\) grammar is restricted to be *strongly linear*, i.e., the right hand side of each rule has at most one nonterminals and there exists a single-stroke path covering the whole of the right hand side (see Yamamoto, et al. [8] for more detail definition), denoted as U-nPCE\(_{\Pi 4}\)-SLAG grammar, then following corollary is obtained:

**Corollary 4.1.** \(\mathcal{F}(K'(U-nPCE_{\Pi 4}-SLAG)) = \mathcal{F}(RAG)\).
Once we use the mapping $K$ instead of $K'$, the situation is entirely changed. In fact, $\mathcal{F}(K(U-nPCE_{\Pi 4}))$ is no longer equal to $\mathcal{F}(CFAG)$.

Lemma 4.2. $\mathcal{F}(K(U-nPCE_{\Pi 4})) \subset \mathcal{F}(CFAG)$. $\mathcal{F}(K(U-nPCE_{\Pi 4}))$ is incomparable with $\mathcal{F}(RAG)$.

Proof: It is easy to see that $\mathcal{F}(K(U-nPCE_{\Pi 4}) \subset \mathcal{F}(CFAG)$. Thus, to prove $\mathcal{F}(RAG)$ is not included in $\mathcal{F}(K(U-nPCE_{\Pi 4}))$, assume that there exists a $U-nPCE_{\Pi 4}$ grammar $G$ such that $K(G)$ is the set $R$ of all rectangles. As shown in Yamamoto, et al. [8], $R$ is in $\mathcal{F}(RAG)$. If such grammar $G$ exist, at least one production rule $P$ is applied to a nonterminal node which is generated from a nonterminal node of the right hand side of $P$ itself. Since otherwise arbitrary large pattern cannot be generated by application of rules whose right hand side have constant size. If so, we can remove the array pattern generated from the first application of $P$ and then connect the array pattern generated from the second application of $P$ without any shearing effects since $K$ is defined for $G$. Such a removed array pattern is finitely large unless there exists another repeated rule $P'$. So $G$ is not the set of all rectangles. If such $P'$ exists, then repeated application process of these rules proceed independently. Again $G$ is not the set of all rectangles. This is a contradiction.

The results of this section are summarized in the following theorem:

Theorem 4.1. The diagram in Fig. 2 holds.

References

$O_{1} \rightarrow B \begin{array}{c} h \\ 1 \end{array} C \begin{array}{c} 2 \end{array}

\psi(h) = (1, h)
\psi(v) = (1, v)
\psi(vh) = (2, v)

(a)

$H \begin{array}{c} h \\ 1 \end{array} J
I \begin{array}{c} h' \\ 1 \end{array} A

(b)

c)

Fig. 1. An example of derivations of nPCE grammar.

$\mathcal{F}(K'(U\text{-nPCE}_{\Pi 4})) = \mathcal{F}(\text{CFAG})$

$\mathcal{F}(K'(U\text{-nPCE}_{\Pi 4} - \text{SLAG})) = \mathcal{F}(\text{RAG})$

Fig. 2. Hierarchical results.