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<tr>
<th>項目</th>
<th>内容</th>
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<tbody>
<tr>
<td>Title</td>
<td>The Problem of Normal Form for Unlabeled Boundary NLC Graph Languages</td>
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The Problem of Normal Form for Unlabeled Boundary NLC Graph Languages

Koichi Yamazaki * Takeo Yaku †

Abstract

In the previous paper (Rozenberg and Welzl 1986 a), it is open that whether there exists a positive integer $k_0$ such that there is a BNLC grammar $G$ with $\max r(G) \leq k_0$ and $L = \text{und}(L(G))$ for every unlabeled BNLC language $L$, where $\max r(G)$ is the maximum number of nodes productions in $G$, and $\text{und}(L(G))$ is the set of underlying unlabeled graphs which are obtained from graphs in $L(G)$ by taking off the labels. In order to negatively solve this open problem, we first show a pumping lemma for a BNLC languages. Then we will show that there is no integer $k_0$ satisfying the above conditions, using the pumping lemma.

1 Introduction

NLC graph grammars are introduced by Janssens and Rozenberg as a basic framework for the mathematical investigation of graph grammars. Since then this model has been intensively investigated (see, e.g., Janssens and Rozenberg 1980, Ehrenfeucht et al 1984 and Janssens et al 1984 ). BNLC graph grammars are introduced and investigated by Rozenberg and Welzl in (Rozenberg and Welzl 1986 a,b). BNLC graph grammars are NLC graph grammars with the property whenever -in a graph already generated- two nodes may be rewritten, then those nodes are not adjacent. BNLC languages are an attractive subfamily of the family of NLC languages (see 1986 b).

As is necessary in string grammars, in graph grammars it is important to investigate normal forms for grammars. Let $G$ be a context-free string grammar, and let $\max r(G)$ be the maximum length of the right side of the productions in $G$.

“Whether there is a positive integer $k_0$, such that for every context-free string language there is a context-free string grammar $G$ with $\max r(G) \leq k_0$ and $L = L(G)$

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The answer of the above problem for $k_0 = 2$ is well known as the Chomsky normal form. In this paper, we call the above problem the Chomsky normal form for context-free string grammars. By Ehrenfeucht et al (1984), for BNLC graph grammar the following problem was investigated: "Whether there is a positive integer $k_0$, such that for every NLC graph language $L$ there is a NLC graph grammar $G$ with $\max r(G) \leq k_0$ and $L = L(G)$?", where $\max r(G)$ is maximum of the number of nodes of axiom and graphs of right hand side of productions in $G$. In (Ehrenfeucht et al 1984), it was shown that there is no positive integer $k_0$ such that $k_0$ satisfies the above condition. This problem is similar to the Chomsky normal form for context-free string grammars. In this paper, we call the above problem the Chomsky normal form problem for NLC graph grammars.

In (Rozenberg and Welzl 1986 a), the Chomsky normal form problem for BNLC graph grammar was investigated: "Whether there is a positive integer $k_0$ such that for every BNLC graph language $L$ there is a BNLC graph grammar $G$ with $\max r(G) \leq k_0$ and $L = L(G)$?". In (Rozenberg and Welzl 1986 a), it was shown that for every BNLC graph grammar there is no positive integer $k_0$ such that $k_0$ satisfies the above condition as was previously shown in the NLC case. In (Rozenberg and Welzl 1986 a), however it is an open problem that whether there is the Chomsky normal form for unlabeled BNLC languages, i.e., "Whether there is a positive integer $k_0$, such that for every unlabeled BNLC language $L$ there is a BNLC grammar $G$ with $\max r(G) \leq k_0$ and $L = \text{und}(L(G))$?", where $\text{und}(L(G))$ is the set of underlying unlabeled graphs of $L(G)$. It turns out that there exists no positive integer $k_0$ of this problem. In this paper, by pumping lemma for NLC languages, we will show that there exists no the Chomsky normal form for unlabeled BNLC languages, i.e., there exists no positive integer $k_0$ that satisfy the following condition: For every unlabeled BNLC language $L$, there is a BNLC graph grammar $G$ such that $\max r(G) \leq k_0$ and $L = \text{und}(L(G))$.

This paper is organized as follows. In Section 2, definitions basic notions. In Section 3, pumping lemma for BNLC languages are given. In Section 4, for every $k$, we construct a u-BNLC language $L_k$ such that $L_k$ is never constructed by a BNLC graph grammar with $\max r(G) \leq k$, and give properties of $L_k$. In Section 5, it is shown, using the pumping lemma and $L_k$, that there exists no the Chomsky normal form for unlabeled BNLC languages.

2 Preliminaries

We start with basic notations concerning graphs, graph grammars and concrete derivation which we need for this paper. For details, see (1986 a). We based on the definition in (1986 a). We assume familiarity with elementary graph theory. In particular, we use the following notions as defined in Harary (1969): adjacent, neighbor, subgraphs, induced subgraphs.
2.1 Graphs

We consider finite undirected node labeled graphs without loops and without multiple edges. For a set of labels $\Sigma$, a graph $X$ (over $\Sigma$) is specified by a finite set $V_X$ of nodes, a set $E_X$ of two element subsets of $V_X$ (called the set of edges) and a function $\varphi_X$ from $V_X$ into $\Sigma$ (called the labeling function). Notice that $E_X$ is a set of the elements of $V_X$. Disregarding the labeling function one gets the underlying unlabeled graph of $X$, denoted by $\text{und}(X)$. The set of all graphs over $\Sigma$ is denoted by $G_\Sigma$. The graph $X - x$ is the subgraph of $X$ induced by $V_X - \{x\}$. The neighborhood of $x$ in $X$, $\text{nbh}_X(x)$, is the set $\{\varphi_X(y) \mid \{x, y\} \in E_X\}$. A graph $X'$ is isomorphic to $X$, if there is a bijection from $V_{X'}$ to $V_X$ which preserves labels and adjacencies. The set of all graphs isomorphic to $X$ is denoted by $[X]$. The size of $X$, $\#X$ is the number of nodes in $X$. We also denote the cardinality of $V_X$ by $\#V_X$, i.e., $\#X = \#V_X$.

Let $\Phi$ be a set of labels. A graph $X$ is called a $\Phi$ - boundary graph, if no two nodes of $X$ that are labeled by elements of $\Phi$ are adjacent.

2.2 Graph Grammars

A node label controlled (NLC) grammar is a system $G = (\Sigma, \Delta, P, \text{conn}, Z_{ax})$, where $\Sigma$ is a finite nonempty set of labels, $\Delta$ is a nonempty subset of $\Sigma$ (the set of terminals), $P$ is a finite set of pairs $(d, Y)$ where $d \in \Sigma$ and $Y \in G_\Sigma$ (the set of productions), conn is a function from $\Sigma$ into $2^\Sigma$ (the connection function), and $Z_{ax} \in G_\Sigma$ (the axiom).

By $[P]$ we denote the set $\{(d, Y') \mid Y' \in [Y] \text{ for some } (d, Y) \in P\}$. By $\text{maxr}(G)$ we denote $\text{maxr}([\#Z_{ax}] \cup \{\#Y \mid (d, Y) \in P \text{ for some } d \in \Sigma\})$. The set $\Sigma - \Delta$ is referred to as the set of nonterminals. We define the set of nonterminal nodes by $\Gamma$, i.e., $\Gamma = \Sigma - \Delta$. In the context of $G$, given a graph $X \in G_\Sigma$ we refer to nodes labeled by elements of $\Gamma$ (respectively) as nonterminal nodes (terminal nodes, respectively).

Let $X, Y$ and $Z$ be graphs over $\Sigma$ with $V_X \cap V_Y = \emptyset$ and let $x \in V_X$. Then $X$ concretely derives $Z$ (in $G$, replacing $x$ by $Y$), denoted by $X \Rightarrow_{\sigma} (x, Y) Z$ or simply by $X \Rightarrow_{(x, Y)} Z$, if $(\varphi_X(x), Y) \in [P]$, $V_Z = V_{X-x} \cup V_Y$, $E_Z = E_{X-x} \cup E_Y \cup \{(x', y) \mid x' \in \text{ nbh}_X(x), y \in V_Y, \varphi_X(x') \in \text{ conn}(\varphi_Y(y))\}$, $\varphi_Z$ equals to $\varphi_{X-x}$ on $V_{X-x}$, and $\varphi_Z$ equals to $\varphi_Y$ on $V_Y$. Intuitively speaking, we replace $x$ in $X$ by the graph $Y$ and connect a node $y$ of $Y$ to a neighbor $x'$ of $x$ if and only if $\varphi_X(x') \in \text{ conn}(\varphi_Y(y))$.

A graph $X$ directly derives a graph $Z$ (in $G$), in symbols $X \Rightarrow_{\sigma} Z$, if there is a graph $Z' \in [Z]$, such that $X$ concretely derives $Z'$ in $G$. $\Rightarrow_{\sigma}^*$ is the transitive and reflexive closure of $\Rightarrow_{\sigma}$. If $X \Rightarrow_{\sigma}^* Z$, then we say that $X$ derives $Z$ (in $G$). The language of $G$ is the set $\{X \in G_\Delta \mid Z_{ax} \Rightarrow_{\sigma}^* X\}$. A set $L$ of graphs is an NLC language if there is an NLC grammar $G$ such that $L = L(G)$. A boundary NLC (BNLC) grammar is an NLC grammar $G = (\Sigma, \Delta, P, \text{conn}, Z_{ax})$, where $Z_{ax}$ is a $\Gamma$ - boundary graph and for all $(d, Y) \in P$, $d \in \Gamma$ and $Y$ is a $\Gamma$ - boundary graph. A graph language $L$ is an NLC (BNLC) language, if there is a NLC (BNLC) grammar.
$G$ such that $L = L(G)$. A language $L$ of unlabeled graphs is $u$-$NLC$ ($u$-$BNLC$)
language, if there is an NLC (BNLC) language $L'$ such that $L = \text{und}(L')$. Let $G$
is NLC graph grammar. $G$ is chain-free, if for all $(d, Y) \in P$ with $V_Y = \{y\}$ (i.e.,
$\|Y = 1)$, $y$ is a terminal node.

2.3 Concrete Derivation

Let $G = (\Sigma, \Delta, P, \text{conn}, Z_{ax})$ be an NLC grammar. If a graph $X$ concretely
derives a graph $Z$ in $G$, replacing a node $x$ by a graph $Y$, then we refer to the
construct $X \Rightarrow_{(x,Y)} Z$ as a concrete derivation step in $G$ (from $X$ to $Z$). A sequence
of successive concrete derivation steps in $G$

$$D : X_0 \Rightarrow_{(x_0,Y_1)} X_1 \Rightarrow_{(x_1,Y_2)} \cdots \Rightarrow_{(x_{n-1},Y_n)} X_n,$$

where $n \geq 0$ and the sets $V_{X_0}, V_{Y_i}, 1 \leq i \leq n$, are pairwise disjoint, is referred to as
a concrete derivation in $G$ (from $X_0$ to $X_n$).

The node set of $D$ is $V_D = \bigcup_{n=0}^{\infty} V_{X_i}$. The edge set of $D$ is $E_D = \bigcup_{n=0}^{\infty} E_{X_i}$. The
labeling function $\varphi_D$ of $D$ is defined by $\varphi_D(x) = \varphi_{X_0}(x)$ if $x \in V_{X_0}$ and $\varphi_D(x) =
\varphi_{Y_i}(x)$ if $x \in V_{Y_i}$, for some $i$, $1 \leq i \leq n$. Note that $V_D = V_{X_0} \cup \bigcup_{i=1}^{n} V_{Y_i}$, hence $\varphi_D$
is defined on the whole set $V_D$. Moreover, if $x \in V_{X_i}$, for some $i$, $1 \leq i \leq n$, then
$\varphi_{X_i}(x) = \varphi_D(x)$. Thus every concrete derivation $D$ defines naturally a graph with
set of nodes $V_D$, set of edges $E_D$ and labeling function $\varphi_D$. Note that this graph
$D$ is a $\Gamma$-boundary graph whenever $X_0$ is a $\Gamma$-boundary graph and $G$ is a BNLC
grammar.

Let $O_D$ be a distinguished element not in $V_D$ which is called the origin of the
derivation $D$. The predecessor mapping $\text{pred}_D$ of $D$ is a function from $V_D$ into
$V_D \cup \{O_D\}$ such that for $x \in V_D$

$$\text{pred}_D(x) = \begin{cases} O_D & \text{if } x \in V_{X_0} \\
\text{x}_{i} & \text{if } x \in V_{Y_{i+1}} \text{ for an } i, 0 \leq i \leq n - 1 \end{cases}$$

Hence $\text{pred}_D$ maps every node $x$ in $V_D$ to the node from which $x$ is directly
derived (or to $O_D$ if $x$ was already present in $X_0$).

The history $\text{hist}_D(x)$ of a node $x \in V_D$ in $D$ is the sequence $(y_0, y_1, \cdots, y_m)$,
$m \geq 1$, $y_i \in V_D$ for all $i$, $1 \leq i \leq m$, such that $y_0 = O_D$, $y_m = x$, and $y_i = \text{pred}_D(y_{i+1})$
for all $i$, $0 \leq i \leq m - 1$. Let $D$ be a derivation and let $x$ and $y$ be nodes in $V_D$. A
node $x$ is an ancestor of $y$ in the derivation $D$ if $x \in \text{hist}_D(y)$. Nodes $x$ and $y$ are
independent if $x \notin \text{hist}_D(y)$ and $y \notin \text{hist}_D(x)$. Let $(y_0, y_1, \cdots, y_m)$ be a sequence
such that $\text{hist}_D(x) = (y_0, y_1, \cdots, y_m)$, and let $0 \leq i < j \leq m$. Then we denote the
sequence $(y_i, y_{i+1}, \cdots, y_j)$ by $\text{hist}_D(y_i, y_j)$. (We can define $\text{hist}_D(x, y)$ only when $x$
is an ancestor of $y$). Let

$$D : X_0 \Rightarrow_{(x_0,Y_1)} X_1 \Rightarrow_{(x_1,Y_2)} \cdots \Rightarrow_{(x_{n-1},Y_n)} X_n$$

be a derivation. We denote the set $\{x_0, x_1, \cdots, x_{n-1}\}$ of rewritten nonterminal nodes
in derivation $D$ by $C_D$. We call the graph $X_n$ the result in the derivation $D$.  

Let $D'$ be derivation

$$D': X_0' \Rightarrow_{(x_0',y_1')} X_1' \Rightarrow_{(x_1',y_2')} \cdots \Rightarrow_{(x_{n-1}',y_n')} X_n'.$$

The derivation $D'$ is isomorphic to $D$ if there is a bijection $h$ from $V_{D'}$ to $V_D$ such that $h|_{V_{X_i}'}$ is an isomorphism from $V_{X_i}'$ to $V_{X_i}$ and for $x' \in V_{X_i}' (0 \leq i \leq n - 1)$, $h(pred_{D'}(x')) = pred_D(h(x'))$. The set of all derivations isomorphic to $D$ is denoted by $[D]$.

Let $D$ be a derivation.

$$D : X_0 \Rightarrow_{(x_0,y_1)} X_1 \Rightarrow_{(x_1,y_2)} \cdots \Rightarrow_{(x_{n-1},y_n)} X_n$$

We call the graph $X_n$ the result in the derivation $D$.

3 A pumping lemma for BNLC languages

In this section, we introduce a pumping lemma for BNLC grammars. In this paper we need the pumping lemma is the proof of the main theorem.

In order to state the pumping lemma, we need to develop some concepts. Let

$$G = (\Sigma, \Delta, P, conn, Z_{ax})$$

be a BNLC grammar. Let $D$ be a following concrete derivation in $G$:

$$D : X_0 \Rightarrow_{(x_0,y_1)} X_1 \Rightarrow_{(x_1,y_2)} \cdots \Rightarrow_{(x_{n-1},y_n)} X_n,$$

such that there exist $x_p, x_q \in C_D (x_p \neq x_q)$ satisfying $x_p \in hist_D(x_q)$ and $\varphi_D(x_p) = \varphi_D(x_q)$. We call the above derivation rough derivation on $(x_p, x_q)$.

Let us consider the following derivation $D'$,

$$D' : X_0' \Rightarrow_{(x_0'y_0'+1)} X_1' \Rightarrow_{(x_1'y_1'+1)} \cdots \Rightarrow_{(x_{n-1}'y_{n-1}'+1)} X_n',$$

where the sequence $(i(0), i(1), \cdots, i(e-1))$ is a sequence ascending order of the subscripts in the set $\{x_l \in C_D | x_p \in hist_D(x_l) \text{ and } x_q \notin hist_D(x_l)\}$.

In what follows, we will show that the derivation $D'$ can be iterated in this section. Let $m$ be a positive integer. Let $D'_1, D'_2, \cdots, D'_m \in [D]$ be derivations, and $h_j$ be an isomorphism from $D'_j$ to $D$ for each $j, 1 \leq j \leq m$ of the following forms:

$$D'_j : X_0' \Rightarrow_{(x_{i(0)'},y_{i(0)'+1})} X_1' \Rightarrow_{(x_{i(1)'},y_{i(1)'+1})} X_2' \Rightarrow_{(x_{i(e-1)'},y_{i(e-1)'+1})} X_e'.$$

$$h_j : V_{D'_j} \rightarrow V_{iter(D,x_p,x_q)}, \ x_l' \mapsto x_l, \text{ where } l \in \{i(0), i(1), \cdots, i(e-1)\}, \text{ which}$$

satisfy the following four conditions:
Then we obtain the following derivation:

\[ X_0'' \Rightarrow_{(x_0,Y_1)} X_1'' \cdots \Rightarrow_{(x_{q-1},Y_q)} X_q^{n} \]

\[ \Rightarrow_{(x_{i(0)}^{1},Y_{i(0)+1}^{1})} X_{q+1}^{n} \cdots \Rightarrow_{(x_{i(e-1)}^{1},Y_{(e-1)+1}^{1})} X_{q+e}^{n} \]

\[ \Rightarrow_{(x^{2}:' Y^{2})} X_{q+e+1}'' \cdots \Rightarrow_{(x^{2},Y^{2})} X_{q+2e}'' \]

\[ \Rightarrow_{(x_{q}^{m},Y_{q+1})} X_{q+me+1}'' \cdots \Rightarrow_{(x_{q+1},Y_{q+2})} X_{q+me+2}^{n} \]

\[ \cdots \Rightarrow_{(x_{n-1},Y_{n})} X_{n+me}'' \]

We call the derivation (*) \textit{m times pumped derivation} and denote \textit{pump}(D, x_p, x_q, m).

Now we can state the pumping lemma for BNLC languages.

\begin{lemma} \textbf{(A pumping lemma for BNLC languages)} \end{lemma}

Let \(D\) be a rough derivation on \((x_p, x_q)\) and \(\text{pump}(D, x_p, x_q, m)\) be the \(m\) times pumped derivation on \((x_p, x_q)\), where \(m\) is an arbitrary positive integer. If the result of \(D\) is in \(G_\Delta\) then the result of \(\text{pump}(D, x_p, x_q, m)\) is in \(G_\Delta\).

\section{Properties of \(L_k\)}

We treat Chomsky normal form problem for u-BNLC languages, i.e., “Whether there is a positive integer \(k_0\), such that for every u-BNLC language \(L\) there is a BNLC grammar \(G\) with \(\max r(G) \leq k_0\) and \(L = \text{und}(L(G))\) ?”.

We will solve the above problem by proving the following theorem:}

\begin{theorem} \textbf{Theorem 4.1} \end{theorem}

For each positive integer \(k\), there is a u-BNLC language \(L_k\) such that the following condition hold: For all BNLC grammar \(G\) with \(\max r(G) \leq k\), \(\text{und}(L(G)) \neq L_k\) hold.

It is not difficult to see that if theorem 4.1 is true then there is no positive integer of the problem. In this section we will construct \(L_k\) of theorem 4.1, and we will deal properties on \(L_k\).
For each positive integer $k$, $L_k$ of theorem 4.1 is construct by the following method. 

A construction method of $L_k$: For each positive integer $k$, we consider a BNLC grammar $G_k = (\Sigma_k, \Delta_k, P_k, \text{conn}_k, Z_{axk})$, where $\Sigma_k = \{a_1, a_2, \ldots, a_k, s\}$, $\Delta_k = \{a_1, a_2, \ldots, a_k\}$, $\text{conn}_k(a_i) = a_i$ for all $1 \leq i \leq k$, $\text{conn}_k(s) = \Delta_k$, $Z_{axk}$ is a single node with label $s$, $P_k = \{(s, Y_{1k}), (s, Y_{2k})\}$, where $Y_{1k}$ is complete graph with set of nodes $\{u_1, u_2, \ldots, u_k, u_{k+1}\}$, where $\varphi_{Y_{1k}}(u_i) = a_i$ for all $1 \leq i \leq k$, $\varphi_{Y_{1k}}(u_{k+1}) = s$, $Y_{2k}$ is complete graph with set of nodes $\{v_1, v_2, \ldots, v_k\}$, where $\varphi_{Y_{2k}}(u_i) = a_i$ for all $1 \leq i \leq k$. We define an unlabeled graph language $L_k$ by $L_k = \text{und}(L(G_k))$.

An underlying unlabeled $L_k$ has the following characteristic properties.

**Proposition 4.2** Let $k \geq 2$ be an arbitrary integer, $e$ be a positive integer and let $H$ be a graph in $L(G_k)$ such that $|V_H| = k \cdot e$. Then
1. The graph $H$ has exactly $e$ disjoint complete subgraphs with $k$ nodes.
2. The graph $H$ has exactly $k$ disjoint complete subgraphs with $e$ nodes.

We call the complete subgraphs of (1) in proposition 4.2 different label group, and the complete subgraphs of (2) in proposition 4.2 same label group.

Same label groups and different label groups has characteristic properties. We will show properties of same label groups and different label groups.

We consider an underlying unlabeled graph $H_u$ in $L_k$. Let $H$ be a graph such that $\text{und}(H) = H_u$, and $D$ be a derivation of $H$ in $G_k$:

$$D : X_0 \Rightarrow_{(x_0, Y_1)} X_1 \Rightarrow_{(x_1, Y_2)} \cdots \Rightarrow_{(x_{n-1}, Y_n)} X_n$$

Let $x$ and $y$ be nodes in $H_u$. We say that nodes $x$ and $y$ come under an identically same label group in $H$ if $\varphi_H(x) = \varphi_H(y)$, and that nodes $x$ and $y$ come under an identically different label group in $H$ if $x \in V_{Y_i}$ and $y \in V_{Y_j}$ for some $1 \leq i \leq n - 1$.

**Proposition 4.3** Let $H_u$ be an underlying unlabeled graph with $k \cdot e$ nodes in $L_k$, $E_1, E_2, \ldots, E_k$ be different label groups in $H_u$, and let $E_1, E_2, \ldots, E_k$ be same label groups in $H_u$. For any nodes $x$ and $y$ in $V_{H_u}$, if $x$ and $y$ are adjacent then
1. Nodes $x$ and $y$ come under an identically same label group and do not come under an identically different label group. or
2. Nodes $x$ and $y$ come under an identically different label group and do not come under an identically same label group.

As a consequence of this proposition, we obtain the following Lemma 4.4.

**Lemma 4.4** Let $H_u$ be an underlying unlabeled graph with more than $k^2(k \geq 2)$ nodes in $L_k$, and let $F$ be a complete subgraph with more than $k$ nodes in $H_u$. Then for all nodes $x$ and $y$ in $V_F$, $x$ and $y$ come under an identically same label group.
5 Proof of the theorem

In this section we will show that there is no BNLC grammar $G$ such that $\text{maxr}(G) \leq k$ and $L_k = \text{und}(L(G))$ by leading contradiction.

We assume that there exists a BNLC grammar $G$ such that $\text{maxr}(G) \leq k$ and $L_k = \text{und}(L(G))$. And we consider a graph $H^1 \in L(G)$ using pumping lemma for a graph $H \in L(G)$. Then we show that $\text{und}(H^1) \notin L_k$.

From now on assume that there exists BNLC grammar $G = (\Sigma, \Delta, P, \text{conn}, Z_{ax})$ such that $L_k = \text{und}(L(G))$ and $\text{maxr}(G) \leq k$. Let $a = \downarrow \Sigma, b = \downarrow \Delta, H$ be a graph in $L(G)$ with $k(k+2+f-g)(k-1)$ nodes, and let $D$ be a derivation of $H$ in $G$. By the supposition, $\text{und}(H) \in L_k$. Hence by the (2) of proposition 4.2 $\text{und}(H)$ has $k$ complete subgraphs with $(k+2+f-g)(k-1)$ nodes. We denote the $k$ complete subgraphs in the graph $H$ by $E_1, E_2, \cdots, E_k$. (For all $1 \leq j \leq k$ complete subgraph in $\text{und}(H)$ that correspond with complete subgraph $E_j$ in $H$ is also denoted by $E_j$.)

In the derivation $D$ we denote the nonterminal node that yield $j$th created nodes in $E_i$ by $y_i^j$. For convenience we also denote the nonterminal node that yield last created nodes in $E_i$ by $y_i^{e(i)}$. Then for each $1 \leq i \leq k$, set of nonterminal nodes \{y_1^i, y_2^i, \cdots, y_{e(i)}^i\} is denoted by $N_i$.

Remarks.

(1) If a nonterminal node $y$ yield terminal nodes $u \in E_i$ and $v \in E_j$, then $y \in N_i \cap N_j$.

(2) For $1 \leq i \leq k$, $e(i) \geq k + f - g + 2$.

For each $1 \leq i \leq k$, we denote by $N_i'$ set \{y_n^i \mid y_n^i \in N_i, k + 1 \leq n \leq e(i)\}. If a graph $X$ in the derivation $D$ has a node $x \in N_i'$ then $X$ has at least $k + 1$ nodes of $E_i$. Because at least $k + 1$ nonterminal nodes in $N_i'$ must be rewritten to rewrite nonterminal node $y_{i+k+2}$. For $1 \leq i \leq k$ there exists at least a pair of nodes with identically nonterminal label, i.e., there exist a pair of nodes $y_{h(i)}^i, y_{l(i)}^i$ such that $\varphi_D(y_{h(i)}^i) = \varphi_D(y_{l(i)}^i) \in \Sigma^\ast \Delta$, where $k+1 < h(i) < t(i) \leq e(i)$. Then for all positive integer $m$ and each $1 \leq i \leq k$ we consider a derivation $\text{pump}(D, y_{h(i)}^i, y_{l(i)}^i, m)$ that is constructed by $D_{i \text{ copy}_1}, D_{i \text{ copy}_2}, \cdots, D_{i \text{ copy}_m} \in \text{orig}(D, y_{h(i)}^i, y_{l(i)}^i)$. Let $h_{i,j}$ is isomorphism from $V_{D_{i \text{ copy}_1}}$ to $V_{D_{i \text{ copy}_j}}$ for each $1 \leq j \leq m$.

We denote by $E_i(\text{orig}(D, y_{h(i)}^i, y_{l(i)}^i))$ set of nodes \{y \in V_{\text{orig}(D, y_{h(i)}^i, y_{l(i)}^i)} \mid y \in E_i\}, and by $E_i(D_{i \text{ copy}_j})$ set of nodes \{y \in V_{D_{i \text{ copy}_j}} \mid h_{i,j}(y) \in E_i(\text{orig}(D, y_{h(i)}^i, y_{l(i)}^i))\}. For convenience we denote by $x_{[i,h(i)],j}(u), x_{[i,l(i)],j}(v)$ respectively node $u, v$ such that $h_{i,j}(u) = y_{h(i)}^i, h_{i,j}(v) = y_{l(i)}^i$. For the above $y_{h(i)}^i, y_{l(i)}^i$ ($1 \leq i \leq k$), the following lemma is satisfied.

**Lemma 5.1** There exists sequence $\text{hist}_D(y_{h(i)}^i, y_{l(i)}^i)$ that hold the following condition:

Condition: For node $y \in \text{hist}_D(y_{h(i)}^i, y_{l(i)}^i)$, if in the derivation $D$ a graph $X$ has $y$, then $y$ and all terminal nodes of graph $X$ are adjacent.
We construct graph $H^{1}$ as following: Let $hist_{D}(x_{h}, x_{t})$ be sequence that be obtained by Algorithm A of Lemma 5.1, and $pump(D, x_{h}, x_{t}, 1)$ be a derivation that be constructed by the derivation $D$ and $D_{copy1} \in [iter(D, x_{h}, x_{t})]$, and let $H^{1}$ be a result of the derivation $pump(D, x_{h}, x_{t}, 1)$. Then the following lemma hold.

Lemma 5.2 $und(H^{1}) \notin L_{k}$.

References


