

## Spectrum and connectivity of graphs

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**Problem** Let  $\Gamma$  be a nice graph. Show that  $\Gamma$  is very connected.

In this talk I would like to give three examples of results about the connectivity of a graph that follow by considering its spectrum.

Three measures of connectivity play a rôle here:

- (i) is  $\Gamma$  connected or not?
- (ii)  $\kappa(\Gamma)$ , the vertex connectivity of  $\Gamma$ , that is, the minimum number of vertices that one has to remove in order to disconnect  $\Gamma$ .
- (iii)  $t(\Gamma)$ , the toughness of  $\Gamma$ , is defined as

$$\min_S \frac{|S|}{c(\Gamma \setminus S)}$$

where  $S$  runs over all sets of vertices such that  $\Gamma \setminus S$  is disconnected, and  $c(\Gamma \setminus S)$  is its number of connected components.

The graph  $K_0$  without vertices is not connected (we have  $c(K_0) = 0$ , while  $c(\Gamma) = 1$  for connected graphs  $\Gamma$ ) but I shall leave undefined whether it is disconnected, and hence do not define  $\kappa(\Gamma)$  and  $t(\Gamma)$  when  $\Gamma$  is complete.

For example, for the Petersen graph we find  $\kappa(\Gamma) = 3$  and  $t(\Gamma) = \frac{4}{3}$ . More generally, we clearly have  $\kappa(\Gamma) \leq k(\Gamma)$  if  $k(\Gamma)$  is the (minimal) valency of  $\Gamma$ . One may also ask about the size of 'nonlocal' cut sets. For example,

(1) ('unimodality') Is it true that if  $S$  is a cut set of  $\Gamma$ , with separation  $\Gamma \setminus S = A + B$ , then  $\min(|\Gamma(S) \cap A|, |\Gamma(S) \cap B|) \leq |S|$ ? (Here  $\Gamma(S)$  denotes the set of all vertices adjacent to some vertex of  $S$ .)

[Jack Koolen remarks that some condition is necessary: for each  $i$ ,  $0 \leq i \leq 17$ , the Biggs-Smith graph has a cut set  $S$  of size 17 such that  $|\Gamma(S) \cap A| = 17 + i$ ,  $|\Gamma(S) \cap B| = 34 - i$ .]

(2) Show that  $|S|$  is substantially larger than  $k$  when  $S$  is nonlocal (say, given a lower bound on the size or the minimum valency of each component of  $\Gamma \setminus S$ ).

## 1 The connectivity of strongly regular graphs

**Theorem 1.1** (Brouwer & Mesner [4]) *Let  $\Gamma$  be strongly regular of valency  $k$ . Then  $\kappa(\Gamma) = k$ , and the only cut sets of size  $k$  are the point neighbourhoods.*

Open problems are for example:

(3) Prove the above result for distance-regular graphs.

(4) Let  $\Gamma$  be strongly regular with parameters  $(v, k, \lambda, \mu)$ , and let  $S$  be a disconnecting set not containing any point neighbourhood  $\Gamma(x)$ . Show that  $|S| \geq 2k - 2 - \lambda$ .

(5) Let  $S$  be a disconnecting set such that  $|S \cap \Gamma(x)| \leq \alpha k$  for some fixed  $\alpha$ ,  $0 < \alpha < 1$ , and all vertices  $x$  of  $\Gamma$ . Prove a superlinear (in  $k$ ) lower bound for  $|S|$ .

**Note** (added July '94): Brouwer & Mulder [5] showed  $\kappa(\Gamma) = k$  for graphs with the property that any two distinct vertices have either 0 or 2 common neighbours. This settles (3) in the case  $\lambda \in \{0, 2\}$ ,  $\mu = 2$ .

## 2 The connectedness of generic pieces of generalized polygons

**Theorem 2.1** (Brouwer [2]) *Let  $\Gamma$  be the point graph or the flag graph of a finite generalized polygon. Then the subgraph  $\Delta$  of  $\Gamma$  induced on the set of all vertices far away from ('in general position w.r.t.') a point or flag is connected, except in the cases  $G_2(2)$ ,  ${}^2F_4(2)$  and (for the flag graph)  $B_2(2)$ ,  $G_2(3)$ . A similar result holds more generally for the complement of a geometric hyperplane.*

Open problems:

- (6) Generalize this to near polygons.
- (7) Generalize this to distance-regular graphs.

It is very easy to see that in a strongly regular graph the subgraph on the vertices far away from a point is connected (except when the graph is complete multipartite).

### 3 The toughness of a regular graph

**Theorem 3.1** (Alon-Brouwer, cf. [1, 3]) *Let  $\Gamma$  be a graph on  $v$  vertices, regular of valency  $k$ , and with eigenvalues  $k = \theta_1 \geq \theta_2 \geq \dots \geq \theta_v$ . Put*

$$\lambda := \max_{2 \leq j \leq v} |\theta_j|.$$

Then

$$t(\Gamma) > \frac{k}{\lambda} - 2.$$

Open problems:

- (8) Prove  $t(\Gamma) \geq \frac{k}{\lambda} - 1$ . (I conjecture that this is the right bound.)
- (9) Prove  $t(\Gamma) = \frac{k}{\lambda}$  in many cases.

**Examples** We have bipartite graphs of small toughness, so the ‘ $-1$ ’ would be best possible. The Delsarte-Hoffman bound for cliques  $C$  in strongly regular graphs states

$$|C| \leq \frac{v}{1 + k/(-\theta_v)}.$$

If equality holds, and  $\lambda = -\theta_v$  (as is often the case), then we find with  $S = \Gamma \setminus C$ :  $t(\Gamma) \leq (v - |C|)/|C| = \frac{k}{\lambda}$ .

### 4 Tools

How are these results proved? Essentially, only interlacing (cf. Haemers [6]) is used. Interlacing comes in two main forms:

- (i) If  $\Delta$  is an induced subgraph of a graph  $\Gamma$ , then the eigenvalues  $\eta_j$  ( $1 \leq j \leq u$ ) of  $\Delta$  interlace the eigenvalues  $\theta_i$  ( $1 \leq i \leq v$ ) of  $\Gamma$ : we have  $\theta_i \geq \eta_i$  ( $1 \leq i \leq u$ ) and  $\eta_{u-j} \geq \theta_{v-j}$  ( $0 \leq j \leq u - 1$ ).

(ii) Given a partition  $\Pi$  of the index set of a symmetric matrix  $A$ , let  $B = (B_{R,S})_{R,S \in \Pi}$  be the matrix of average row sums of the corresponding submatrices of  $A$ . Then the eigenvalues of  $B$  interlace those of  $A$ .

### Examples

**Lemma 4.1** *The average valency of a graph is not more than its largest eigenvalue.*

**Proof:** Use a partition with 1 part. □

**Lemma 4.2** *Let  $\Gamma$  be regular of valency  $k$  on  $v$  vertices, and let the graph induced on the  $r$ -set  $R$  have average valency  $k_R$ . Then*

$$\theta_2 \geq (vk_R - rk)/(v - r) \geq \theta_v,$$

(and hence

$$r \leq v(k_R - \theta_v)/(k - \theta_v).$$

For  $k_R = 0$  we find the Delsarte-Hoffman bound).

**Proof:** Use a partition with 2 parts. □

**Lemma 4.3** *Let  $\Gamma$  and  $R$  be as before. Put  $\lambda = \max(|\theta_2|, |\theta_v|)$ . Then*

$$\sum_x (|\Gamma(x) \cap R| - \frac{rk}{v})^2 \leq \lambda^2 r(v - r)/v.$$

**Proof:** Use a partition with 2 parts, and apply to  $A^2$ , the square of the adjacency matrix of  $\Gamma$ . □

## 5 Proofs of the results in sections 1,2,3

Let  $\Gamma$  have eigenvalues  $k = \theta_1 \gg \theta_2 \geq \dots$ . If  $\Delta$  is a disconnected subgraph, then its spectrum is the union of the spectra of its components. Each component has a largest eigenvalue at least as large as its average degree, and by interlacing it follows that all components except perhaps one have average degree at most  $\theta_2$ , but this is much too small (except when  $\Gamma$  is very small).

This proves the results of Sections 1 and 2. For those of Section 3, use the above three Lemmata and compute.

## References

- [1] Noga Alon, *Tough Ramsey graphs without short cycles*, preprint, 1993.
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- [4] A.E. Brouwer & D.M. Mesner, *The connectivity of strongly regular graphs*, *Europ. J. Combin.* **6** (1985) 215-216.
- [5] A.E. Brouwer & H.M. Mulder, *The vertex connectivity of a  $\{0, 2\}$ -graph equals its valency*, preprint, 1994.
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