数理解析研究所講究録
516巻 1984年 79-91

Countable J_a^S-admissible ordinals

Juichi SHINODA (篠田 壽一)

Department of Mathematics
College of General Education
Nagoya University

§0. Introduction.

In [3], Platek constructs a hierarchy of jumps J_a^S indexed by elements a of a set 0^S of ordinal notations. He asserts that a real $X \subseteq \omega$ is recursive in the superjump S if and only if it is recursive in some J_a^S. Unfortunately, his assertion is not correct as is shown in [1]. In [1], it also has been shown that an ordinal $>\omega$ is J_a^S-admissible if it is $|a|_S$-recursively inaccessible, where $|a|_S$ is the ordinal denoted by a.

Let A be an arbitrary set. We say that an ordinal α is A-admissible if the structure $< L_\alpha[A], \in, A \cap L_\alpha[A] >$, which we will denote by $L_\alpha[A]$ for simplicity, is admissible, a model of the Kripke-Platek set theory KP, where $L_\alpha[A]$ is the sets constructible relative to A in fewer than α steps. We use ω_1^A or $\omega_1(A)$ to denote the first A-admissible ordinal $>\omega$, and use $\omega_1(A_1,\cdots,A_n)$ for $\omega_1(<A_1,\cdots,A_n>)$.

The purpose of this paper is to prove the following theorem.

<u>Theorem 1</u>. Suppose $a \in 0^S$ and $\alpha > \omega$ is a countable $|a|_S$-recursively inaccessible ordinal. Then, there exists a real $X \subseteq \omega$ such that $\alpha = \omega_1(J_a^S, X)$.

In the case $|a|_S = 0$, $J_a^S = {}^2E$, the Kleene object of type 2, and $\omega_1({}^2E, X) = \omega_1^X$ for all reals $X \subseteq \omega$. α is an admissible ordinal if and only if it is 0-recursively inaccessible. Therefore, Theorem 1 is an extension of the following theorem of Sacks.

Theorem 2. (Sacks [4]). If $\alpha > \omega$ is a countable admissible ordinal, then there exists a real X such that $\alpha = \omega_1^X$.

Sacks also showed that the real X mentioned in Theorem 2 can be taken to have the minimality property:

$$\omega_1^Y < \alpha \quad \text{for every } Y \text{ of lower hyperdegree than } X.$$

Likewise, we can show that for every countable $|a|_S$-recursively inaccessible ordinal $\alpha > \omega$ there is a real X such that:

$$\alpha = \omega_1(J_a^S, X) ;$$

and

$$\omega_1(J_a^S, Y) < \alpha \quad \text{for every } Y \text{ of lower } J_a^S\text{-degree than } X.^{1)}$$

Theorem 1 will be proved by the forcing with J_a^S-pointed perfect trees. Let $\alpha > \omega$ be a countable $|a|_S$-recursively inaccessible ordinal and X be a generic real with respect to this forcing relation. Then $L_\alpha[X]$ is admissible and $\alpha \leqq \omega_1(J_a^S, X)$. To see $\omega_1(J_a^S, X) \leqq \alpha$, we must show that X preserves sufficiently many admissible ordinals below α to make α to be $<J_a^S, X>$-admissible.

§1. $|a|_S$-recursively inaccessible ordinals.

A normal type 2 object is a total function F from ω^ω to ω such that the Kleene object ${}^2E$ of type 2:

$$^2E(f) = \begin{cases} 0 & \text{if } (\exists n) f(n) = 0, \\ 1 & \text{otherwise,} \end{cases}$$

is recursive in F. The superjump $S(F)$ of F is a type 2 object

1) For J_a^S-degrees, the reader may refer to [5].

defined by:

$$S(F)(\langle n, f\rangle) = \begin{cases} 0 & \text{if } \{n\}^F(f) \text{ is defined,} \\ 1 & \text{otherwise.} \end{cases}$$

Platek [3] defines a hierarchy J_a^S of type 2 objects along with a set 0^S of ordinal notations, starting from ${}^2E$ and iterating the superjump operation transfinitely.

An ordinal α is 0-recursively inaccessible if it is admissible. α is $(\sigma+1)$-recursively inaccessible if it is σ-recursively inaccessible and a limit of σ-recursively inaccessible ordinals. For limit λ, α is said to be λ-recursively inaccessible if it is σ-recursively inaccessible for all $\sigma < \lambda$. Let X be an arbitrary set. σ-recursively-in-X inaccessible ordinals are defined in the same way starting from X-admissible ordinals. By $RI(\sigma, X)$, we denote the class of all σ-recursively-in-X inaccessible ordinals. In the case $X = \phi$, $RI(\sigma, \phi)$ is the class of all σ-recursively inaccessible ordinals.

The following lemma, due to Aczel and Hinman, gives a characterization of $\omega_1(J_a^S, X)$ for $X \subseteq \omega$.

Lemma 3. (Aczel and Hinman [1]). Suppose $a \in 0^S$ and $\sigma = |a|_S$, the ordinal denoted by a. Then $\sigma < \omega_1(J_a^S)$, and for any ordinal $\alpha > \omega$ and $X \subseteq \omega$:

$$\alpha \in RI(\sigma, X) \longrightarrow \alpha \text{ is } \langle J_a^S, X\rangle\text{-admissible.}$$

$\omega_1(J_a^S, X)$ is the least ordinal in $RI(\sigma, X)$.

Let λ_0 be the least ordinal λ such that λ is λ-recursively inaccessible. Lemma 3 shows that $|0^S| = \sup\{|a|_S : a \in 0^S\} \leq \lambda_0$. In [1], it has shown that $|0^S| = \lambda_0$.

Let $\alpha > \omega$ be a countable admissible ordinal. Using the unbounded

Levy forcing over L_α, we can add to L_α a generic function $K : (\alpha - \omega) \times \omega \to \alpha$ such that if $\omega \leq \beta < \alpha$ then the function $\lambda n K(\beta, n)$ is a bijection from ω onto β. Therefore, in $L_\alpha[K]$ all sets are countable. It has been shown in [4] that $< L_\alpha[K], \in, K >$ is an admissible structure in which Σ_1 - DC (Σ_1-Dependent Choice) holds.

Suppose $a \in O^S$. For any $X, Y \subseteq \omega$, $X \leq_{J^S_a} Y$ means X is recursive in $< J^S_a, Y >$, which is equivalent to that $X \in L_\rho[Y]$, where $\rho = \omega_1(J^S_a, Y)$. X and Y have the same J^S_a - degree, $X \equiv_{J^S_a} Y$, if $X \leq_{J^S_a} Y$ and $Y \leq_{J^S_a} X$. $X <_{J^S_a} Y$ if $X \leq_{J^S_a} Y$ but $X \not\equiv_{J^S_a} Y$.

Lemma 4. Suppose $\alpha > \omega$ is a countable $|a|_S$-recursively inaccessible ordinal and K is a generic function with respect to the unbounded Levy forcing over L_α. Then for any $X, Y \subseteq \omega$:

$$X \leq_{J^S_a} Y \quad \& \quad Y \in L_\alpha[K] \longrightarrow X \in L_\alpha[K].$$

Proof. The unbounded Levy forcing preserves admissible ordinals. That is, if $\beta < \alpha$ is an admissible ordinal then β is K-admissible. This is because for admissible β, $K \upharpoonright (\beta - \omega) \times \omega$ is generic with respect to the unbounded Levy forcing over L_β. Therefore, if $Y \in L_\alpha[K]$ then α is $|a|_S$-recursively-in-Y inaccessible, so $L_\rho[Y] \subseteq L_\alpha[K]$, where $\rho = \omega_1(J^S_a, Y)$. Thus we have the lemma. □

§2. J^S_a-pointed perfect trees.

Let a be an element of O^S such that $|a|_S > 0$. We put $J = J^S_a$ for simplicity.

A perfect tree is a set P of finite sequences of 0's and 1's such that:

(1) $p \in P$ & $q \subseteq p \longrightarrow q \in P$;

and

(2) $(\forall p \in P)(\exists q, r \in P)$ (q and r are incompatible extensions of p),

where $q \subseteq p$ denotes that p is an extension of q. For a perfect tree P, $[P]$ denotes the set of all infinite paths through P:

$$[P] = \{f \in 2^{\omega} : (\forall n)\bar{f}(n) \in P\}.$$

We say that P is J-pointed if:

(3) $(\forall f \in [P])(\omega_1(J, P) \leq \omega_1(J, f)$ & $P \in L_{\omega_1(J,P)}[f])$.

Note that if P is J-pointed then it is $\leq_J$-pointed in the sense of Sacks [4: 2.1], but not vice versa.

Lemma 5. Suppose P is J-pointed. If $X \subseteq \omega$ and $P \leq_J X$, then there exists a J-pointed $Q \subseteq P$ such that $Q \equiv_J X$.

Proof. In [4: 2.3], Sacks constructed a perfect subtree Q of P such that:

(4) Q is recursive in P and f for every $f \in [Q]$;

and

(5) $Q \equiv_J X$.

To see Q is J-pointed in our sense, fix $f \in [Q]$. Since P is J-pointed and $f \in [P]$, by (3), we have:

(6) $P \in L_{\omega_1(J,P)}[f]$.

Clearly:

(7) $f \in L_{\omega_1(J,P)}[f]$.

From (4) (6) and (7), we obtain:

(8) $Q \in L_{\omega_1(J,P)}[f]$.

From (5) and the assumption $P \leq_J X$, we see:

(9) $\omega_1(J, P) \leqq \omega_1(J, Q)$.

From (8) and (9), we obtain $Q \in L_{\omega_1(J,Q)}[f]$. □

For any ordinal δ, $\{\delta\}^f$ denotes the δ-th element of $L[f]$ in the canonical wellordering on $L[f]$. A perfect tree P is said to be uniformly J-pointed if there exists an ordinal δ such that:

(10) $(\forall f \in [P])(P = \{\delta\}^f \;\&\; \delta < \omega_1(J, f))$.

Obviously, uniformly J-pointed perfect trees are J-pointed. Let $\alpha > \omega$ be a countable $|a|_S$-recursively inaccessible ordinal and K a generic function over L_α in the sense of the unbounded Levy forcing. Observe that if P is uniformly J-pointed and $P \in L_\alpha[K]$ then there exists a $\delta < \alpha$ which satisfies (10) since the leftmost path f_P through P is recursive in P and so $\omega_1(J, f_P) \leqq \omega_1(J, P) < \alpha$.

Let M be a countable admissible set and P be a perfect tree in M. Then P becomes a partially ordered set as usual. The forcing with P as the set of conditions is called the local Cohen forcing over M and denoted by $\Vdash^P_M$, or simply by $\Vdash^P$. If $f \in [P]$ is generic with respect to $\Vdash^P$, then $M[f]$ is an admissible set, and so is $L_\mu[f]$, where $\mu = M \cap On$.

Lemma 6. For any $\xi < \alpha$ and any J-pointed perfect tree P in $L_\alpha[K]$, there exists a uniformly J-pointed perfect tree $Q \subseteq P$ such that $\xi < \omega_1(J, Q)$ and $Q \in L_\alpha[K]$.

Proof. Since ξ is countable in $L_\alpha[K]$, there is a real $X \in L_\alpha[K]$ such that ξ is recursive in X. By Lemma 5, there is a J-pointed perfect subtree P_1 of P such that $P_1 \equiv_J X$. Then we see $\xi < \omega_1(J, P_1)$, and $P_1 \in L_\alpha[K]$ by Lemma 4. Thus, we may assume $\xi < \omega_1(J, P)$ from the beginning. Put $M = L_{\omega_1(J,P)}[P]$. Consider the local Cohen forcing relation $\Vdash^P_M$ over M. Since P is J-pointed, we have:

(11) $(\forall f \in [P])\omega_1(J, P) \leqq \omega_1(J, f)$;

and

(12) $(\forall f \in [P])(\exists\gamma < \omega_1(J, P))\ \{\gamma\}^f = P$.

By (12), there exists a $p_0 \in P$ and $\gamma < \omega_1(J, P)$ such that:

(13) $p_0 \Vdash^P_M \{\check{\gamma}\}^{\mathcal{J}} = \check{P}$,

where $\mathcal{J}$ is the canonical name which denotes the generic reals. As in [4: 2.10], we can construct a perfect tree $Q \subseteq P$ such that:

(14) $Q \in L_{\omega_1(J,P)}[P]$;

and

(15) $(\forall f \in [Q])\ \{\gamma\}^f = P$.

From (14), we can find a $\delta < \omega_1(J, P)$ such that $\{\delta\}^P = Q$. So, by (15), there is an $\varepsilon < \omega_1(J, P)$ such that:

(16) $(\forall f \in [Q])\ \{\varepsilon\}^f = Q$.

Let f_Q be the leftmost branch of Q. Then, by (11):

(17) $\omega_1(J, P) \leqq \omega_1(J, f_Q) \leqq \omega_1(J, Q)$.

Hence, from (16), we see that Q is uniformly J-pointed. By (17), we also see $\xi < \omega_1(J, Q)$. Since $P \in L_\alpha[K]$, we have $\omega_1(J, P) \leqq \alpha$, and so $Q \in L_{\omega_1(J,P)}[P] \subseteq L_\alpha[K]$. □

Let $\mathcal{L}$ be a first-order language. A Π^1_1 formula in $\mathcal{L}$ is a second-order formula of the form:

$$(\forall S_1) \cdots (\forall S_m) \psi,$$

where $S_1, \cdots, S_m$ are predicate variables and ψ is first-order formula in the expanded language $\mathcal{L} \cup \{S_1, \cdots, S_m\}$.

Lemma 7. Suppose A is a countable admissible set such that $\omega \in A$ and $\mathcal{L} \in A$ is a first-order language. Let $\theta(x_1, \cdots, x_n)$ be a Π^1_1 formula in $\mathcal{L}$. Then there exists a Σ_1 formula $\Phi(x_1, \cdots, x_n, y)$ such that for any structure $\mathfrak{M} = \langle M, \cdots \rangle \in A$ for $\mathcal{L}$ and any $a_1, \cdots, a_n \in M$:

$$A \models \Phi(a_1, \cdots, a_n, \mathfrak{M}) \longleftrightarrow \mathfrak{M} \models \theta(a_1, \cdots, a_n).$$

Proof. This is well-known. See, e.g., Barwise [2: IV. 3.1]. □

Using this lemma, we obtain the following lemma.

Lemma 8. The set of all uniformly J-pointed perfect trees in $L_\alpha[K]$ is Σ_1 over $L_\alpha[K]$.

Proof. Put $\sigma = |a|_S$, (recall that $J = J^S_a$). Let P be a perfect tree in $L_\alpha[K]$ and $\delta < \alpha$. Let $\beta(P, \delta, \sigma)$ denote the least admissible ordinal $\beta < \alpha$ such that $\max(\delta, \sigma, \omega) < \beta$ and $P \in L_\beta[K]$. The function β is Σ_1 over $L_\alpha[K]$. We can easily find a Π^1_1 formula θ in the language of set theory such that for any perfect tree $P \in L_\alpha[K]$:

$$P \text{ is uniformly J-pointed} \longleftrightarrow (\exists \delta < \alpha) L_{\beta(P,\delta,\sigma)}[K] \models \theta(P, \delta, \sigma).$$

Thus the lemma follows from Lemma 7. □

§3. Forcing with uniform J^S_a - pointed perfect trees.

Suppose $|a|_S > 0$ and put $J = J^S_a$. Let $\alpha > \omega$ be a countable $|a|_S$-recursively inaccessible ordinal and K a generic function with respect to the unbounded Levy forcing over L_α, which we fix throughout this section.

Let $\mathcal{L}(\alpha, \mathcal{J})$ be a ramified language containing names for all members of $L_\alpha[f]$. $\mathcal{L}(\alpha, \mathcal{J})$ includes: a numeral $\bar{n}$ for each $n \in \omega$, unranked variables $x, y, z, \cdots$; ranked variables $x^\beta, y^\beta, z^\beta, \cdots$ for each $\beta < \alpha$; and abstraction operator $\hat{}$. It is intended that $\mathcal{J}$ denotes $\{n \in \omega: f(n) = 1\}$, that x ranges over $L_\alpha[f]$, that x^β ranges over $L_\beta[f]$, and that $\hat{x}^\beta \phi(x^\beta)$ denotes the set:

$$\{x \in L_\beta[f] : L_\beta[f] \models \phi(x)\}.$$

$C(\beta)$ is the set of names for elements of $L_\beta[f]$ and $C = \bigcup_{\beta<\alpha} C(\beta)$.

Let $\mathcal{P}$ denote the set of all uniformly J-pointed perfect trees in $L_\alpha[K]$. $P, Q, R, \cdots$ denote the members of $\mathcal{P}$. For a ranked sentence ϕ of $\mathcal{L}(\alpha, \mathcal{J})$ and $P \in \mathcal{P}$, let $\rho(P, \phi)$ be the least admissible ordinal $\rho < \alpha$ such that $P \in L_\rho[K]$ and $\text{rank}(\phi) < \rho$. The function ρ is Σ_1 over $L_\alpha[K]$. The forcing relation $P \Vdash \phi$, where ϕ is a sentence of $\mathcal{L}(\alpha, \mathcal{J})$, is defined inductively:

(1) ϕ is ranked. $P \Vdash \phi$ iff $(\forall f \in [P])\, L_{\rho(P,\phi)}[f] \models \phi$;

(2) $\phi \vee \psi$ is not ranked. $P \Vdash \phi \vee \psi$ iff $P \Vdash \phi$ or $P \Vdash \psi$;

(3) $(\exists x^\beta)\phi(x^\beta)$ is not ranked. $P \Vdash (\exists x^\beta)\phi(x^\beta)$ if $P \Vdash \phi(c)$ for some $c \in C(\beta)$;

(4) $P \Vdash (\exists x)\phi(x)$ iff $P \Vdash \phi(c)$ for some $c \in C$;

(5) ϕ is not ranked. $P \Vdash \neg\phi$ iff $(\forall Q \subseteq P)\neg(Q \Vdash \phi)$.

Using Lemma 7 and 8, it is easy to see that the forcing relation $P \Vdash \phi$, restricted Σ_1 sentences ϕ, is Σ_1 over $L_\alpha[K]$.

Lemma 9. For each P and ϕ, there exists a $Q \subseteq P$ such that $Q \Vdash \phi$ or $Q \Vdash \neg \phi$.

Proof. In view of (5), we may assume that ϕ is ranked. By Lemma 6, we may also assume that $\phi \in L_\delta[P]$ for some P-admissible δ such that $\delta < \omega_1(J, P)$. Then, in $L_\delta[P]$, all sets are countable. Thus, in $L_\delta[P]$, we can enumerate all ranked sentences of rank $\leq$ rank(ϕ) :

$$\phi = \phi_0, \phi_1, \cdots, \phi_n, \cdots \quad (n \in \omega).$$

Let $\Vdash^P$ be the local Cohen forcing relation over $L_\delta[P]$. In $L_\delta[P]$, we can construct a family $< q_s : s \in Seq(2) >$ of elements of P such that:

(6) $q_s \Vdash^P \phi_n$ or $q_s \Vdash^P \neg \phi_n$, where $n = lh(s)$;

and

(7) $q_{s\frown<0>}$ and $q_{s\frown<1>}$ are incompatible extensions of q_s,

where Seq(2) is the set of all finite sequences of 0's and 1's. Let $Q = \{q \in P : (\exists s)\, q \subseteq q_s\}$. Then by (7) Q is a perfect subtree of P. By (6), it is easy to see that $Q \Vdash \phi$ or $Q \Vdash \neg \phi$. Since $Q \in L_\delta[P]$, $Q = \{\gamma\}^P$ for some $\gamma < \delta$. Therefore Q is uniformly J-pointed because P is. □

A real $f \in 2^\omega$ is said to be generic if for every dense subset $\mathcal{D}$ of $\mathcal{P}$ which is definable over $L_\alpha[K]$ there is a $P \in \mathcal{D}$ such that $f \in [P]$. For every $P \in \mathcal{P}$, there is a generic f such that $f \in [P]$. From Lemma 9, it follows that for every generic f and sentence ϕ:

$$L_\alpha[f] \models \phi \text{ iff } (\exists P)(f \in [P] \;\&\; P \Vdash \phi).$$

<u>Lemma 10</u>. If f is generic, then $L_\alpha[f]$ is admissible.

<u>Proof</u>. We need to show that $L_\alpha[f]$ satisfies the Δ_0 Collection. Let $\phi(x, y)$ be a formula of $\mathcal{L}(\alpha, \mathcal{T})$ with no unranked quantifiers. We claim that if $P \Vdash (\forall n)(\exists y)\phi(n, y)$ then there exists a $Q \subseteq P$ and a $\beta < \alpha$ such that $Q \Vdash (\forall n)(\exists y^\beta)\phi(n, y^\beta)$. The proof of this claim is almost the same as that of [4: 3.12] with some notational changes. So, we omit the proof here. From the claim, it follows that $L_\alpha[f]$ satisfies the Δ_0 Collection. □

<u>Proof of Theorem 1</u>. Let $\alpha > \omega$ be a countable $|a|_S$-recursively inaccessible ordinal and K be as before. Put $\sigma = |a|_S$ and $J = J_a^S$. In the case $\sigma = 0$, Theorem 1 is exactly Theorem 2, which has already been established by Sacks [4]. So we may assume $\sigma > 0$. Let $f_0 \in 2^\omega$ be a generic real over $L_\alpha[K]$ with respect to the forcing with uniform J-pointed perfect trees. By Lemma 6, for each $\xi < \alpha$, the set $\{P \in \wp : \xi < \omega_1(J, P)\}$ is dense in $\wp$. It is obviously definable over $L_\alpha[K]$. Therefore there is a $P \in \wp$ such that $f_0 \in [P]$ and $\xi < \omega_1(J, P)$. Since P is J-pointed, we have:

$$\xi < \omega_1(J, P) \leqq \omega_1(J, f_0).$$

Thus, we have $\alpha \leqq \omega_1(J, f_0)$. To see $\alpha = \omega_1(J, f_0)$, we must show that $\alpha \in RI(\sigma, f_0)$. At first we consider the case where $\sigma = \tau + 1$ for some τ. It is sufficient to prove that α is a limit of ordinals in $RI(\tau, f_0)$, since then by induction on τ we can show that $\alpha \in RI(\tau, f_0)$, (note that $\alpha \in RI(0, f_0)$ by Lemma 10). Suppose $\xi < \alpha$. We shall show that the following set $\mathcal{D}_\xi$ is dense in $\mathcal{P}$:

$$\mathcal{D}_\xi = \{P \in \mathcal{P} : (\exists \delta < \alpha)(\xi < \delta \ \& \ (\forall f \in [P])\delta \in RI(\tau, f))\}.$$

Assume this can be done. Using Lemma 7, it is easy to see that $\mathcal{D}_\xi$ is Σ_1 over $L_\alpha[K]$. Therefore, for every $\xi < \alpha$, there exists a $\delta < \alpha$ such that $\xi < \delta$ and $\delta \in RI(\tau, f_0)$.

To show that $\mathcal{D}_\xi$ is dense in $\mathcal{P}$, take an arbitrary $P \in \mathcal{P}$. By Lemma 6, we may assume $\xi < \omega_1(J, P)$. Take a $\delta \in RI(\tau, P)$ so that $\xi < \delta < \omega_1(J, P)$. Such a δ exists because $\omega_1(J, P)$ is a limit of ordinals in $RI(\tau, P)$. Consider the local Cohen forcing relation $\Vdash^P$ over $L_\delta[P]$. Let δ^+ be the next P-admissible ordinal of δ. Then, $L_\delta[P]$ is countable in $L_{\delta+}[P]$. So we can enumerate inside $L_{\delta+}[P]$ all sentences of the appropriate forcing language:

$$\phi_0, \phi_1, \cdots, \phi_n, \cdots \qquad (n \in \omega).$$

As in the proof of Lemma 9, we can construct a perfect subtree $Q \in L_{\delta+}[P]$ of P such that:

$$(\forall f \in [Q])\ f \text{ is generic with respect to } \Vdash^P.$$

Q is uniformly J-pointed since $Q \in L_{\delta+}[P]$, $\delta^+ < \omega_1(J, P)$ and P is uniformly J-pointed. To show that $\delta \in RI(\tau, f)$ for all $f \in [Q]$, take $f \in [Q]$. Let $\beta \leqq \delta$ be an arbitrary P-admissible ordinal $> \omega$, and $\Vdash^P_\beta$ be the local Cohen forcing relation over $L_\beta[P]$. It is easy to see that f is generic with respect to $\Vdash^P_\beta$, and so β is f-admissible. From this, by induction on τ, we see that $\delta \in RI(\tau, f)$.

Now we consider the case whre σ is a limit ordinal. The proof is carried out in the same way. For any $\xi < \alpha$ and any $\tau < \sigma$, let $\mathcal{D}_{\xi\tau}$ be the set:

$$\{P \in \mathcal{P} : (\exists \delta < \alpha)(\xi < \delta \ \ \& \ \ (\forall f \in [P])\ \delta \in RI(\tau, f)\}.$$

Then $\mathcal{D}_{\xi\tau}$ is dense in $\mathcal{P}$ and definable over $L_\alpha[K]$. Therefore, we have that $\alpha = \omega_1(J, f_0)$ for any generic f_0 with respect to $\Vdash$. □

REFERENCES

[1] P. Aczel and P.G. Hinman, Recursion in the superjump, in: Generalized Recursion Theory, edited by J.E. Fenstad and P.G. Hinman (North-Holland, Amsterdam, 1974), 3-41.

[2] J. Barwise, Admissible Sets and Structures, Springer, Berlin, 1975.

[3] R. Platek, A countable hierarchy for the superjump, in: Logic Colloquium '69, edited by R.O. Gandy and C.E.M. Yates (North-Holland, Amsterdam, 1971), 257-271.

[4] G.E. Sacks, Countable admissible ordinals and hyperdegrees, Advances in Math., 19(1976), 213-262.

[5] J. Shinoda, On the upper semi-lattice of J_a^S-degrees, Nagoya Math. J., 80(1980), 75-106.

Department of Mathematics
College of General Education
Nagoya University
Chikusa-ku, Nagoya 464, Japan