

# CM-TRIVIALITY AND GEOMETRIC ELIMINATION OF IMAGINARIES

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## 1. INTRODUCTION

To show CM-triviality (of generic relational structures), first of all, we showed weak elimination of imaginaries, and then, working in the real sort, we could show CM-triviality. In this note, we show that *CM-triviality in the real sort*, defined in the second section, implies geometric elimination of imaginaries and CM-triviality (in the real and imaginary sorts). To show this, we give a characterization of geometric elimination of imaginaries in simple theories.

Our notation is standard. Let  $T$  be a complete  $L$ -theory, and let  $\mathcal{M}$  be the big model of  $T$ .  $\bar{a}, \bar{b}, \dots (\subset_{\omega} \mathcal{M})$  denote finite sequences in  $\mathcal{M}$ . We work in  $\mathcal{M}^{\text{eq}}$ , which consists of  $\bar{a}_E$ , the  $E$ -class of  $\bar{a}$ , for any 0-definable equivalence relation  $E$  and  $\bar{a} \subset_{\omega} \mathcal{M}$ .  $AB$  denotes  $A \cup B$  for any  $A, B \subset \mathcal{M}^{\text{eq}}$ .

For  $a \in \mathcal{M}^{\text{eq}}, A \subset \mathcal{M}^{\text{eq}}$ , we write  $a \in \text{dcl}^{\text{eq}}(A)$ , if  $a$  is fixed by any automorphism pointwise fixing  $A$ . And we write  $a \in \text{acl}^{\text{eq}}(A)$ , if the orbit of  $a$  by automorphisms pointwise fixing  $A$ , is finite. We write  $\bar{a} \equiv_A \bar{b}$  for  $\text{tp}(\bar{a}/A) = \text{tp}(\bar{b}/A)$  in  $T$ .

We said that  $T$  geometrically eliminates imaginaries ( $T$  has GEI), if for any  $e \in \mathcal{M}^{\text{eq}}$ , there exists  $\bar{b} \subset_{\omega} \mathcal{M}$  such that  $e \in \text{acl}^{\text{eq}}(\bar{b})$  and  $\bar{b} \in \text{acl}^{\text{eq}}(e)$ .

## 2. A CHARACTERIZATION OF GEI IN SIMPLE THEORIES

Let  $T$  be a simple theory.

**Definition 2.1.** We say that  $T$  has *the independence over intersections* ( $T$  has IND/I), for any  $\bar{a}, A, B \subset \mathcal{M}$  with  $\bar{a} \perp_A B, \bar{a} \perp_B A$ , we have  $\bar{a} \perp_{\text{acl}(A) \cap \text{acl}(B)} AB$ .

**Proposition 2.2.** *IND/I implies GEI.*

*Proof.* Fix  $e = \bar{a}_E \in \mathcal{M}^{\text{eq}}$ . Take  $\bar{b}, \bar{c} \models \text{tp}(\bar{a}/e)$  such that  $\bar{b}, \bar{c}, \bar{a}$  are independent over  $e$ . Let  $A = \text{acl}(\bar{b}) \cap \text{acl}(\bar{c})$ . Then  $\bar{a} \downarrow_A \bar{b}\bar{c}$  by IND/I. By  $e \in \text{dcl}^{\text{eq}}(\bar{a}) \cap \text{dcl}^{\text{eq}}(\bar{b}\bar{c})$ ,  $e \in \text{acl}^{\text{eq}}(A)$ . On the other hand,  $A \subset \text{acl}^{\text{eq}}(e)$  follows from  $\bar{b} \downarrow_e \bar{c}$ .  $\square$

**Lemma 2.3.** *Suppose that  $T$  has GEI. Then, for any  $\text{acl}(A) = A, \text{acl}(B) = B \subset \mathcal{M}$ , we have*

$$\text{acl}^{\text{eq}}(A) \cap \text{acl}^{\text{eq}}(B) = \text{acl}^{\text{eq}}(A \cap B).$$

*Proof.* Let  $e \in \text{acl}^{\text{eq}}(A) \cap \text{acl}^{\text{eq}}(B)$ . By GEI, there exists  $\bar{a} \subset_{\omega} \mathcal{M}$  such that  $e \in \text{acl}^{\text{eq}}(\bar{a})$  and  $\bar{a} \in \text{acl}^{\text{eq}}(e)$ . As  $\bar{a} \in \text{acl}^{\text{eq}}(A)$  and  $\bar{a} \in \text{acl}^{\text{eq}}(B)$ , we see  $\bar{a} \subset A \cap B$ . Thus,  $e \in \text{acl}^{\text{eq}}(A \cap B)$ .  $\square$

From now on, we assume **elimination of hyperimaginaries (EHI)**. Then the converse of Proposition 2.2 follows.

**Proposition 2.4.**  $\text{GEI} \Leftrightarrow \text{IND/I}$

*Proof.*  $(\Leftarrow)$  by Proposition 2.2.  $(\Rightarrow)$ : Suppose that  $\bar{a} \downarrow_A B, \bar{a} \downarrow_B A$  and  $\text{acl}(A) = A, \text{acl}(B) = B$ . By the above lemma and EHI, we see  $\text{Cb}(a/AB) \subseteq \text{acl}^{\text{eq}}(A) \cap \text{acl}^{\text{eq}}(B) = \text{acl}^{\text{eq}}(A \cap B)$ .  $\square$

### 3. MAIN THEOREM

**Definition 3.1.** We say that  $T$  is *CM-trivial in the real sort*, if, for any  $\bar{a}, A = \text{acl}(A), B = \text{acl}(B) \subset \mathcal{M}$ ,  $\bar{a} \downarrow_A B$  implies  $\bar{a} \downarrow_{A \cap \text{acl}(\bar{a}, B)} B$ .

**Remark 3.2.** The original definition of CM-triviality is as follows: For any  $a, A = \text{acl}^{\text{eq}}(A), B = \text{acl}^{\text{eq}}(B) \subset \mathcal{M}^{\text{eq}}$ ,  $a \downarrow_A B$  implies  $a \downarrow_{A \cap \text{acl}^{\text{eq}}(a, B)} B$ . Clearly, under assuming GEI, CM-triviality is equivalent to CM-triviality in the real sort. In the next remark, we lay out an example which shows the difference of the definitions.

**Theorem 3.3.** *If  $T$  is CM-trivial in the real sort, then  $T$  has GEI. So CM-triviality in the real sort implies (the original) CM-triviality.*

*Proof.* By Proposition 2.2, we will show that  $T$  has IND/I, i.e. if  $\bar{a}, A = \text{acl}(A), B = \text{acl}(B) \subset \mathcal{M}$  and  $\bar{a} \downarrow_A B, \bar{a} \downarrow_B A$ , then  $\bar{a} \downarrow_{A \cap B} AB$ . By CM-triviality in the real sort, we have  $\bar{a} \downarrow_{\text{acl}(\bar{a}, B) \cap A} B$ . By  $\bar{a} \downarrow_B A$ , we see  $\text{acl}(\bar{a}, B) \cap AB = B$ . As  $A \cap B \subseteq A \cap \text{acl}(\bar{a}, B) \subseteq AB \cap \text{acl}(\bar{a}, B) = B$ , we see

$$\text{acl}(\bar{a}, B) \cap A = A \cap B.$$

By  $\bar{a} \downarrow_{\text{acl}(\bar{a}, B) \cap A} B$  and  $\bar{a} \downarrow_B A$ , we see  $\bar{a} \downarrow_{A \cap B} AB$ .  $\square$

**Remark 3.4.** (1) Let  $T$  be the theory of a simple relational structure with a closure operator  $\text{cl}(\ast)$  such that

- $\text{cl}(\text{acl}(A)) = \text{acl}(A)$  and  $\text{cl}(\text{cl}(A) \cap \text{cl}(B)) = \text{cl}(A) \cap \text{cl}(B)$ ,
- for any algebraically closed sets  $A, B \subset \mathcal{M}$ ,  $A \perp_{A \cap B} B \Leftrightarrow$   
“ $AB = \text{cl}(AB)$  and  $R^{AB} = R^A \cup R^B$  for any predicate  $R$ ”.

Then  $T$  is CM-trivial in the real sort. (Suppose that  $\bar{a} \perp_A B$ . Let  $C = \text{acl}(\bar{a}, A)$ ,  $D = \text{acl}(AB)$ . As  $C \perp_A B$  and  $C \cap B = A$ ,  $\text{cl}(CB) = CB$  and  $R^{CB} = R^C \cup R^B$  for any predicate  $R$ . Let  $E = \text{acl}(\bar{a}, B)$ . Then  $\text{cl}(CB \cap E) = CB \cap E$  and  $R^{CB \cap E} = R^{C \cap E} \cup R^{B \cap E}$  for any predicate  $R$ . So, we see  $C \cap E \perp_{A \cap E} B \cap E$ . As  $\bar{a} \subset C \cap E$ ,  $B \subset B \cap E$ ,  $\bar{a} \perp_{A \cap \text{acl}(\bar{a}, B)} B$  follows.) So, by Theorem 3.3, CM-triviality of  $T$  follows.

- (2) CM-triviality does not imply *CM-triviality in the real sort*: In [E], Evans gave an  $\omega$ -categorical CM-trivial structure  $\mathfrak{C}$ , defined below, of SU-rank one without WEI.

Here, we check that  $\mathfrak{C}$  does not have GEI.

Firstly, he constructed an  $\omega$ -categorical generic structure  $M$  (countable binary graph  $R(x, y)$  with a predimension  $\delta(A) = 2|A| - |R^A|$ ) of SU-rank two such that

- no triangles, no squares in  $M$ , and points and adjacent pairs of points are closed in  $M$
- $\text{cl}(\ast) = \text{acl}(\ast)$  in  $M$  and  $M$  is of diameter 3.

Fix  $a \in M$ . Let  $C, D$  be the sets of vertices at distance 1, 2 from  $a$ . Then we have the canonical structure  $\mathfrak{C}$  on  $C$  such that  $\text{Aut}(\mathfrak{C})$  is homeomorphic to  $\text{Aut}(M/a)$ , so  $\mathfrak{C}$  and  $(M, a)$  are biinterpretable. (See pp.136,139,348 in [H].) Then  $\mathfrak{C}$  is of SU-rank one.

We see that  $\mathfrak{C}$  does not have GEI as follows:

Let  $c, c' \in C$  and  $d, d' \in D$  be such that  $M \models R(a, c) \wedge R(a, c') \wedge R(c, d) \wedge R(c', d')$ . As no triangles and squares in  $M$ , we have  $M \models \neg R(c, c') \wedge \neg R(c, d') \wedge \neg R(c', d)$ . Note that  $c \in \text{dcl}(a, d)$  and  $acd < acdc', acdd'$ . So,  $c', d' \notin \text{cl}(a, d, c) = \text{acl}(a, d, c)$ . Therefore  $\text{cl}(a, d) = \text{acl}(a, d) = \{a, c, d\}$  follows. On the other hand,  $\text{cl}(a, c) = \{a, c\}$ . So, if  $\mathfrak{C}$  has GEI, then, as  $d \in \mathfrak{C}^{\text{eq}}$ , there exist  $\bar{c} \subset_{\omega} C$  such that  $d \in \text{acl}(a, \bar{c})$  and  $\bar{c} \in \text{acl}(a, d)$  in the sense of  $M$ . But such  $\bar{c}$  must be a singleton  $c \in C$  with  $M \models R(a, c) \wedge R(c, d)$ . Since  $\text{acl}(a, c) = \{a, c\}$  in  $M$ , so  $d \notin \text{acl}(a, c)$  in  $M$ .

**Problem 3.5.** *In stable theories, is CM-triviality equivalent to CM-triviality in the real sort?*

## REFERENCES

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